C Model: Memory & Control Flow

Optional Textbook Readings: CP:AMA 6.1–6.4, 7.1–7.3, 7.6, Appendix E

- the ordering of topics is different in the text
- some portions of the above sections have not been covered yet

The primary goal of this section is to be able to model how C programs execute.
Models of computation

In CS 135, we modelled the computational behaviour of Racket with substitutions (the “stepping rules”).

- all arguments are evaluated to values **before** a function can be called (“applied”)

- to call (“apply”) a function, we substitute the *body* of the function, replacing the parameters with the argument values.

```
(define (my-sqr x) (* x x))

(+ 2 (my-sqr (+ 3 1)))
=> (+ 2 (my-sqr 4))
=> (+ 2 (* 4 4))
=> (+ 2 16)
=> 18
```
In this course, we model the behaviour of C with two complimentary mechanisms:

- control flow
- memory
Control flow

We use control flow to model how programs are executed.

During execution, we keep track of the program location, which is “where” in the code the execution is currently occurring.

When a program is “run”, the program location starts at the beginning of the main function.

In hardware, the location is known as the program counter, which contains the address within the machine code of the current instruction (more on this in CS 241).
Types of control flow

In this course, we explore four types of control flow:

- compound statements (blocks)
- function calls
- conditionals (i.e., if statements)
- iteration (i.e., loops)
Compound statements (blocks)

We have already seen compound statements (blocks) where a sequence of statements (and definitions) are executed in order.

```c
int main(void) {
    trace_int(1 + 1);       // first
    assert(3 > 2);         // second
    int i = 7;              // third (i is now in scope)
    printf("%d\n", i);      // fourth
    return 0;              // fifth
}
```
Function Calls

When a function is called, the program location “jumps” from the current location to the start of the function.

The return control flow statement changes the program location to go back to the most recent calling function (where it “jumped from”). Obviously, C needs to “keep track” of where to go. We revisit this when we introduce memory later in this section.
example: function call flow

```c
void blue(void) {
    printf("three\n");
    return;
}

void green(void) {
    blue();
    printf("four\n");
    return;
}

void red(void) {
    printf("two\n");
    green();
    printf("five\n");
    return;
}

int main(void) {
    printf("one\n");
    red();
    printf("six\n");
}
```

There is a supplemental video for this slide online.
Conditionals (if)

We introduced the if control flow statement in Section 02. We now discuss if in more detail.

The syntax of if is

    if (expression) statement

where the statement is only executed if the expression is true (non-zero).

    if (n < 0) printf("n is less than zero\n");

Remember: the if statement does not produce a value. It only controls the flow of execution.
The if statement only affects whether the next statement is executed. To conditionally execute more than one statement, use a compound statement (block).

```c
if (n <= 0) {
    printf("n is zero\n"); // execute this
    printf("or less than zero\n"); // then this
}
```

Using a block with every if is strongly recommended even if there is only one statement. It is good style: it makes code easier to follow and less error prone.

```c
if (n <= 0) {
    printf("n is less than or equal to zero\n");
}
```

(In the notes, we occasionally omit them to save space.)
Statement A;
if (exp) {
    Code Block;
}
Statement Z;
As we have seen, the if statement can be combined with else statement(s) for multiple conditions.

```c
if (expression) {
    statement(s)
} else if (expression) {
    statement(s)
} else if (expression) {
    statement(s)
} else {
    statement(s)
}
```
Statement A;
if (exp1) {
    Code Block 1;
} else if (exp2) {
    Code Block 2;
} else {
    Code Block 3;
}
Statement Z;
If an if block contains a return statement, there may be no need for an else block.

```c
int sum(int k) {
    if (k <= 0) {
        return 0;
    } else {
        return k + sum(k - 1);
    }
}

// Alternate equivalent code
int sum(int k) {
    if (k <= 0) {
        return 0;
    }
    return k + sum(k - 1);
}
```
Braces are sometimes necessary to avoid a “dangling” else.

```c
if (y > 0)
    if (y != 7)
        printf("you lose");
else
    printf("you win!"); // when does this print?
```
The C switch control flow statement (see CP:AMA 5.3) has a similar structure to else if and cond, but very different behaviour.

A switch statement has “fall-through” behaviour where more than one branch can be executed.

In our experience, switch is very error-prone for beginner programmers.

Do not use switch in this course.
The C `goto` control flow statement (CP:AMA 6.4) is one of the most disparaged language features in the history of computer science because it can make “spaghetti code” that is hard to understand.

Modern opinions have tempered and most agree it is useful and appropriate in some circumstances.

To use gotos, *labels* (code locations) are required.

```c
if (k < 0) goto mylabel;
//...
mylabel:
//...
```

Do not use `goto` in this course.
Looping

With mutation, we can control flow with a method known as `looping`.

```c
while (expression) statement
```

while is similar to if: the statement is only executed if the expression is true.

The difference is, while repeatedly “loops back” and executes the statement until the expression is false.

Like with if, always use a block (`{}`) for a compound statement, even if there is only a single statement.
Statement A;
while (exp) {
    Code Block;
}
Statement Z;
example: while loop

```c
int i = 2;
while (i >= 0) {
    printf("%d\n", i);
    --i;
}
```

OUTPUT:
2
1
0

There is a supplemental video for this slide online.
Iteration vs. recursion

Using a loop to solve a problem is called **iteration**.

*Iteration* is an alternative to *recursion* and is much more common in imperative programming.

```c
// recursion
int sum(int k) {
    if (k <= 0) {
        return 0;
    }
    return k + sum(k - 1);
}

// iteration
int sum(int k) {
    int s = 0;
    while (k > 0) {
        s += k;
        --k;
    }
    return s;
}
```
When first learning to write loops, you may find that your code is very similar to using *accumulative recursion*.

```c
int accsum(int k, int acc) {
    if (k <= 0) {
        return acc;
    }
    return accsum(k - 1, k + acc);
}

int recursive_sum(int k) {
    return accsum(k, 0);
}
```

```c
int iterative_sum(int k) {
    int acc = 0;
    while (k > 0) {
        acc += k;
        --k;
    }
    return acc;
}
```

Looping is very “imperative”. Without mutation (side effects), the while loop condition would not change, causing an “endless loop”.

*CS 136 Winter 2022 04: C Model*
Loops can be “nested” within each other.

```c
int i = 5;
int j = 0;
while (i >= 0) {
    j = i;
    while (j >= 0) {
        printf("*");
        --j;
    }
    printf("\n");
    --i;
}
```

```
******
*****
****
***
**
*```
Tracing tools

The provided *tracing tools* can be used to help you understand your control flow and “see” what is happening in your program.

This can help you **debug** your code.

The tracing tools do **not** interfere with your I/O testing.

---

On your assignments, never `printf` any unnecessary output as it may affect your correctness results.

Always use our tracing tools to help debug your code.
int sum(int k) {
    trace_msg("sum called");
    int s = 0;
    trace_msg("loop starting");
    while (k > 0) {
        trace_int(k);
        s += k;
        trace_int(s);
        --k;
    }
    trace_msg("loop ended");
    return s;
}

int main(void) {
    trace_int(sum(3));
}
while errors

A simple mistake with while can cause an “endless loop” or “infinite loop”. The following examples are endless loops (for $i \geq 0$).

```c
while (i >= 0) // missing {}
    printf("%d\n", i);
    --i;

while (i >= 0); { // extra ;
    printf("%d\n", i);
    --i;
}

while (i = 100) { ... } // assignment typo

while (1) { ... } // constant true expression
    // (this may be on purpose...)
```
**do ... while**

The do control flow statement is very similar to while.

    do statement while (expression);

The difference is that statement is always executed *at least* once, and the expression is checked at the *end* of the loop.

    do {
        printf("try to guess my number!\n");
        guess = read_int();
    } while (guess != my_number && guess != READ_INT_FAIL);
Statement A;
do {
    Code Block;
    exp?
    true
} while (!exp);
Statement Z;
break

The break control flow statement is useful to exit from the middle of a loop. break immediately terminates the current (innermost) loop. break is often used with a (purposefully) infinite loop.

```c
while (1) {
    n = read_int();
    if (n == READ_INT_FAIL) {
        break;
    }
    //...
}
```

break only terminates loops. You cannot break out of an if.
The `continue` control flow statement skips over the rest of the statements in the current block (```{}```) and “continues” with the loop.

```c
// only concerned with fun numbers
while (1) {
    n = read_int();
    if (n == READ_INT_FAIL) {
        break;
    }
    if (!is_fun(n)) {
        continue;
    }
    //...
}
```
Statement A;
while (exp) {
    Code Block;
}
Statement Z;
continue;
break;
for loops

The final control flow statement we introduce is for, which is often referred to as a “for loop”.

for loops are a “condensed” version of a while loop.

The format of a while loop is often of the form:

```
setup statement
while (expression) {
  body statement(s)
  update statement
}
```

which can be re-written as a single for loop:

```
for (setup; expression; update) { body statement(s) }
```
for vs. while

Recall the for syntax.

   for (setup; expression; update) { body statement(s) }

This while example

   i = 100; // setup
   while (i >= 0) {
      // expression
      printf("%d\n", i);
      --i; // update
   }

is equivalent to

   for (i = 100; i >= 0; --i) {
      printf("%d\n", i);
   }
Statement A;
for (setup; exp; update) {
    Code Block;
}
Statement Z;
Statement A;
for (setup; exp; update) {
    Code Block;
}
Statement Z;

exp?

true

Code Block

update

false

continue;

break;
Most for loops follow one of these forms (or “idioms”).

    // Counting up from 0 to n - 1
    for (i = 0; i < n; ++i) {...}

    // Counting up from 1 to n
    for (i = 1; i <= n; ++i) {...}

    // Counting down from n - 1 to 0
    for (i = n - 1; i >= 0; --i) {...}

    // Counting down from n to 1
    for (i = n; i > 0; --i) {...}

It is a common mistake to be “off by one” (e.g., using < instead of <=). Sometimes re-writing as a while is helpful.
In C99, the \textit{setup} can be a \textbf{definition}.

This is very convenient for defining a variable that only has \textit{local (block) scope} within the for loop.

\begin{verbatim}
for (int i = 100; i >= 0; --i) {
    printf("\%d\n", i);
}
\end{verbatim}

The equivalent while loop would have an extra block.

\begin{verbatim}
{
    int i = 100;
    while (i >= 0) {
        printf("\%d\n", i);
        --i;
    }
}
\end{verbatim}
Any of the three components of a `for` statement can be omitted.

If the expression is omitted, it is always “true”.

```c
for (; i < 100; ++i) {...}  // i was setup previously
for (; i < 100;) {...}      // same as a while(i < 100)
for (;;) {...}              // endless loop
```

The *comma operator* (, ) allows for multiple sub-expressions in the *setup* and *update* statements of a *for* loop. Do not use it in this course. See CP:AMA 6.3 for more details.

```c
for (i = 1, j = 100; i < j; ++i, --j) {...}
```
A for loop is *not always* equivalent to a while loop.

The only difference is when a continue statement is used.

In a while loop, continue jumps back to the expression.

In a for loop, the “update” statement is executed before jumping back to the expression.
Memory

One bit of storage (in memory) has two possible states: 0 or 1.

A byte is 8 bits of storage. Each byte in memory is in one of 256 possible states.

In this course, we will usually be dealing with bytes and not individual bits.
Accessing memory

The smallest accessible unit of memory is a byte.

To access a byte of memory, its position in memory, which is known as the address of the byte, must be known.

For example, if you have 1 MB of memory (RAM), the address of the first byte is 0 and the address of the last byte is 1048575 (or \(2^{20} - 1\)).

Note: Memory addresses are usually represented in hex, so with 1 MB of memory, the address of the first byte is 0x0, and the address of the last byte is 0xFFFFF.
You can visualize computer memory as a collection of “labeled mailboxes” where each mailbox stores a byte.

<table>
<thead>
<tr>
<th>address</th>
<th>contents</th>
</tr>
</thead>
<tbody>
<tr>
<td>(1 MB of storage)</td>
<td>(one byte per address)</td>
</tr>
<tr>
<td>0x000000</td>
<td>00101001</td>
</tr>
<tr>
<td>0x000001</td>
<td>11001101</td>
</tr>
<tr>
<td>...</td>
<td>...</td>
</tr>
<tr>
<td>0xFFFFFE</td>
<td>00010111</td>
</tr>
<tr>
<td>0xFFFFFF</td>
<td>01110011</td>
</tr>
</tbody>
</table>

The *contents* in the above table are arbitrary values.
Defining variables

For a variable definition, C

- reserves (or “finds”) space in memory to store the variable
- “keeps track of” the address of that storage location
- stores the initial value of the variable at that location (address).

For example, with the definition

```
int n = 0;
```

C reserves space (an address) to store n, “keeps track of” the address n, and stores the value 0 at that address.
In our CS 135 substitution model, a variable is a “name for a value”. When a variable appears in an expression, a *substitution* occurs and the name is *replaced* by its value.

In our new model, a variable is a “name for a location” where a value is stored.

When a variable appears in an expression, C “fetches” the contents at its address to obtain the value stored there.
sizeof

When we define a variable, C reserves space in memory to store its value – but how much space is required?

It depends on the type of the variable.

It may also depend on the environment (the machine and compiler).
The **size operator** (`sizeof`) produces the number of bytes required to store a type (it can also be used on identifiers). `sizeof` looks like a function, but it is an operator.

```c
int n = 0;
trace_int(sizeof(int));
trace_int(sizeof(n));

sizeof(int) => 4
sizeof(n) => 4
```

In this course, the size of an integer is 4 bytes (32 bits).
In C, the size of an int depends on the machine (processor) and/or the operating system that it is running on.

Every processor has a natural “word size” (e.g., 32-bit, 64-bit). Historically, the size of an int was the word size, but most modern systems use a 32-bit int to improve compatibility.

In C99, the inttypes module (#include <inttypes.h>) defines many types (e.g., int32_t, int16_t) that specify exactly how many bits (bytes) to use.

In this course, only use int, and there are always 32 bits in an int.
**example: variable definition**

```c
int n = 0;
```

For this variable definition C reserves (or “finds”) 4 consecutive bytes of memory to store n (e.g., addresses 0x5000…0x5003) and then “keeps track of” the first (or “starting”) address.

<table>
<thead>
<tr>
<th>identifier</th>
<th>type</th>
<th># bytes</th>
<th>starting address</th>
</tr>
</thead>
<tbody>
<tr>
<td>n</td>
<td>int</td>
<td>4</td>
<td>0x5000</td>
</tr>
</tbody>
</table>

C updates the contents of the 4 bytes to store the initial value (0).

<table>
<thead>
<tr>
<th>address</th>
<th>0x5000</th>
<th>0x5001</th>
<th>0x5002</th>
<th>0x5003</th>
</tr>
</thead>
<tbody>
<tr>
<td>contents</td>
<td>00000000</td>
<td>00000000</td>
<td>00000000</td>
<td>00000000</td>
</tr>
</tbody>
</table>
Integer limits

Because C uses 4 bytes (32 bits) to store an int, there are only $2^{32}$ (4,294,967,296) possible values that can be represented.

The range of C int values is $-2^{31} \ldots (2^{31} - 1)$ or $-2,147,483,648 \ldots 2,147,483,647$.

In our CS 136 environment, the constants INT_MIN and INT_MAX are defined with those limit values.

unsigned int variables represent the values $0 \ldots (2^{32} - 1)$ but we do not use them in this course.
In the read_int function we provide, the value of the constant READ_INT_FAIL is actually INT_MIN, so the smallest value of int that can be successfully read by our read_int function is $-2,147,483,647$. 
Overflow

If we try to represent values outside of the int limits, overflow occurs.

Never assume what the value of an int will be after an overflow occurs.

The value of an integer that has overflowed is undefined.

By carefully specifying the order of operations, sometimes overflow can be avoided.

In CS 251 / CS 230 you learn more about overflow.
**example: overflow**

```c
int bil = 1000000000;
int four_bil = bil + bil + bil + bil;
int nine_bil = 9 * bil;

trace_int(bil);
trace_int(four_bil);
trace_int(nine_bil);

bil => 1000000000
four_bil => -294967296
nine_bil => 410065408
```

Remember, do not try to “deduce” what the value of an int will be after overflow – its behaviour is **undefined**.
Racket can handle arbitrarily large numbers, such as
\((\text{expt } 2 \ 1000)\).

Why did we not have to worry about overflow in Racket?

Racket does not use a fixed number of bytes to store numbers. Racket represents numbers with a structure that can use an arbitrary number of bytes (imagine a list of bytes).

There are C modules available that provide similar features (a popular one is available at gmplib.org).
The char type

Now that we have a better understanding of what an int in C is, we introduce some additional types.

The char type is also used to store integers, but C only allocates one byte of storage for a char (an int uses 4 bytes).

There are only $2^8$ (256) possible values for a char and the range of values is $(-128 \ldots 127)$ in our Seashell environment.

Because of this limited range, chars are rarely used for calculations. As the name implies, they are often used to store characters.
ASCII

Early in computing, there was a need to represent text (characters) in memory.

The American Standard Code for Information Interchange (ASCII) was developed to assign a numeric code to each character.

Upper case A is 65, while lower case a is 97. A space is 32.

ASCII was developed when teletype machines were popular, so the characters 0 ... 31 are teletype “control characters” (e.g., 7 is a “bell” noise).

The only control character we use in this course is the line feed (10), which is the newline \n character.
*/

32 space 48 0 64 @ 80 P 96 \ 112 p
33 ! 49 1 65 A 81 Q 97 a 113 q
34 " 50 2 66 B 82 R 98 b 114 r
35 # 51 3 67 C 83 S 99 c 115 s
36 $ 52 4 68 D 84 T 100 d 116 t
37 % 53 5 69 E 85 U 101 e 117 u
38 & 54 6 70 F 86 V 102 f 118 v
39 ' 55 7 71 G 87 W 103 g 119 w
40 ( 56 8 72 H 88 X 104 h 120 x
41 ) 57 9 73 I 89 Y 105 i 121 y
42 * 58 : 74 J 90 Z 106 j 122 z
43 + 59 ; 75 K 91 [ 107 k 123 {
44 , 60 < 76 L 92 \ 108 l 124 |
45 - 61 = 77 M 93 ] 109 m 125 }
46 . 62 > 78 N 94 ^ 110 n 126 ~
47 / 63 ? 79 O 95 _ 111 o
*/
ASCII worked well in English-speaking countries in the early days of computing, but in today’s international and multicultural environments it is outdated.

The **Unicode** character set supports more than 100,000 characters from all over the world.

A popular method of **encoding** Unicode is the UTF-8 standard, where displayable ASCII codes use only one byte, but non-ASCII Unicode characters use more bytes.

We will not be using Unicode in this course.
C characters

In C, **single** quotes (') are used to indicate an ASCII character.

For example, 'a' is equivalent to 97 and 'z' is 122. C “translates” 'a' into 97.

In C, there is **no difference** between the following two variables:

```c
char letter_a = 'a';
char ninety_seven = 97;
```

Always use **single** quotes with characters:

"a" is **not** the same as 'a'.

example: C characters

The printf format specifier to display a character is "\%c".

```c
char letter_a = 'a';
char ninety_seven = 97;

printf("letter_a as a character: \%c\n", letter_a);
printf("ninety_seven as a char: \%c\n", ninety_seven);

printf("letter_a in decimal: \%d\n", letter_a);
printf("ninety_seven in decimal: \%d\n", ninety_seven);
```

letter_a as a character: a
ninety_seven as a char: a

letter_a in decimal: 97
ninety_seven in decimal: 97
Character arithmetic

Because C interprets characters as integers, characters can be used in expressions to avoid having “magic numbers” in your code.

```c
bool is_lowercase(char c) {
    return (c >= 'a') && (c <= 'z');
}

// to_lowercase(c) converts upper case letters to lowercase letters, everything else is unchanged
char to_lowercase(char c) {
    if ((c >= 'A') && (c <= 'Z')) {
        return c - 'A' + 'a';
    } else {
        return c;
    }
}
```
Reading characters from input

In Section 03, we used the `read_int` function to read integers from input.

We have also provided `read_char` for reading characters.

When reading `int` values, we ignored whitespace in the input.

When reading in characters, you **may** or **may not** want to ignore whitespace characters, depending on the application.

`read_char` has a parameter for specifying if whitespace is ignored.
Symbol type

There is no C equivalent of the Racket 'symbol type.

C symbols are constants (often ints) with meaningful identifiers ("names") but arbitrary (meaningless) values.

Use ALL_CAPS for symbol names.

    const int UP = 1;
    const int DOWN = 2;

    int direction = UP;

We have provided some tools for working with C "symbols" on your assignments.
In C, there are **enumerations** (enum, CP:AMA 16.5) which allow you to create your own enum types and help to facilitate defining constants with unique integer values.

Enumerations are an example of a C language feature that we do *not* introduce in this course.

After this course, we would expect you to be able to read about enums in a C reference and understand how to use them.

If you would like to learn more about C or use it professionally, we recommend reading through all of CP:AMA *after* this course is over.
Floating point types

The C `float` (floating point) type can represent real (non-integer) values.

```c
float pi = 3.14159;
float avogadro = 6.022e23;  // 6.022*10^23
```

Unfortunately, floats are susceptible to precision errors.

C’s float type is similar to **inexact numbers** in Racket (which appear with an `#i` prefix in the teaching languages):

```scheme
(sqrt 2) ; => #i1.4142135623730951
(sqr (sqrt 2)) ; => #i2.0000000000000004
```
example 1: inexact floats

```c
float penny = 0.01;
float money = 0;

for (int n = 0; n < 100; ++n) {
    money += penny;
}

printf("the value of one dollar is: \%.4f\n", money);

the value of one dollar is: 0.999999
```

The `printf` format specifier to display a float is "%.f".
example 2: inexact floats

float bil = 1000000000;
float bil_and_one = bil + 1;

printf("a float billion is: %f\n", bil);
printf("a float billion + 1 is: %f\n", bil_and_one);

a float billion is: 1000000000.000000
a float billion + 1 is: 1000000000.000000
In the previous two examples, we highlighted the precision errors that can occur with the float type.

C also has a double type that is still inexact but has significantly better precision.

Just as we use check-within with inexact numbers in Racket, we can use a similar technique for testing in floating point numbers C.

Assuming that the precision of a double is perfect or “good enough” can be a serious mistake and introduce errors.

Unless you are explicitly told to use a float or double, do not use them in this course.
Floats in memory

A double has more precision than a float because it uses more memory.

Just as we might represent a number in decimal as $6.022 \times 10^{23}$, a float uses a similar strategy.

A 32 bit float uses 24 bits for the mantissa and 8 bits for the exponent.

A 64 bit double uses $(53 + 11)$.

floats and their internal representation are discussed in CS 251 / 230 and in detail in CS 370 / 371.
Structures

Structures (compound data) in C are similar to structures in Racket.

```c
struct posn { // name of the structure
    int x;     // type and field names
    int y;
}; // don't forget this ;
```

Because C is statically typed, structure definitions require the type of each field.

Do not forget the last semicolon (;) in the structure definition.
The structure *type* includes the keyword “*struct*”. For example, the type is “*struct posn*”, not just “*posn*”. This can be seen in the definition of *p* below.

```c
struct posn p = {3, 4}; // note the use of {}

trace_int(p.x);
trace_int(p.y);

p.x => 3
p.y => 4
```

Instead of *selector functions*, C has a *structure operator* (.) which “selects” the requested field.

The syntax is *variablename.fieldname*
C99 supports an alternative way to initialize structures:

```c
struct posn p = { .y = 4, .x = 3};
```

This prevents you from having to remember the “order” of the fields in the initialization.

Any omitted fields are automatically zero, which can be useful if there are many fields:

```c
struct posn p = {.x = 3};    // .y = 0
```
Mutation with structures

The assignment operator can be used with structs to copy all of the fields from another struct. Individual fields can also be mutated.

```c
struct posn p = {1, 2};
struct posn q = {3, 4};

p = q;
p.x = 23;
```

```
trace_int(p.x);
trace_int(p.y);
```

```
p.x => 23
p.y => 4
```
The braces ({} ) are part of the initialization syntax and can not simply be used in assignment. Instead, just mutate each field.

On rare occasions, you may want to define a new struct so you can mutate “all at once”.

```c
struct posn p = {1, 2};

p = {5, 6}; // INVALID
p.x = 5;    // VALID
p.y = 6;

// alternatively:
struct posn new_p = {5, 6};
p = new_p;
```
The *equality* operator (==) does not work with structures. You have to define your own equality function.

```c
bool posn_equal (struct posn a, struct posn b) {
    return (a.x == b.x) && (a.y == b.y);
}
```

Also, `printf` only works with elementary types. Print each field of a structure individually:

```c
dstruct posn p = {3, 4};
printf("The value of p is (%d, %d)\n", p.x, p.y);
```

The value of p is (3, 4)
Structures in the memory model

For a structure *definition*, no memory is reserved:

```c
struct posn {
    int x;
    int y;
};
```

Memory is only reserved when a `struct` *variable* is defined:

```c
struct posn p = {3, 4};
```
sizeof a struct

struct mystruct {
    int x; // 4 bytes
    char c; // 1 byte
    int y; // 4 bytes
};

The amount of space reserved for a struct is at least the sum of the sizeof each field, but it may be larger.

trace_int(sizeof(struct mystruct));

sizeof(struct mystruct) => 12

You must use the sizeof operator to determine the size of a structure.
The size may depend on the order of the fields:

```c
struct s1 {
    char c;
    int i;
    char d;
};
```

```c
struct s2 {
    char c;
    char d;
    int i;
};
```

```c
trace_int(sizeof(struct s1));
trace_int(sizeof(struct s2));
```

```c
sizeof(struct s1) => 12
sizeof(struct s2) => 8
```

C may reserve more space for a structure to improve efficiency and enforce alignment within the structure.
Sections of memory

In this course we model five *sections* (or “regions”) of memory:

<table>
<thead>
<tr>
<th>Code</th>
</tr>
</thead>
<tbody>
<tr>
<td>Read-Only Data</td>
</tr>
<tr>
<td>Global Data</td>
</tr>
<tr>
<td>Heap</td>
</tr>
<tr>
<td>Stack</td>
</tr>
</tbody>
</table>

Other courses may use alternative names.

The **heap** section is introduced in Section 10.
Sections are combined into memory **segments**, which are recognized by the hardware (processor).

When you try to access memory outside of a segment, a **segmentation fault** occurs (more on this in CS 350).
**Temporary results**

When evaluating C expressions, the intermediate results must be temporarily stored.

\[ a = f(3) + g(4) - 5; \]

In the above expression, C must temporarily store the value returned from \( f(3) \) “somewhere” before calling \( g \).

In this course, we are not concerned with this “temporary” storage.

Temporary storage is discussed in CS 241.
The code section

When you program, you write *source code* in a text editor using ASCII characters that are “human readable”.

To “run” a C program, the *source code* must first be converted into *machine code* that is “machine readable”.

This machine code is then placed into the *code section* of memory where it can be executed.

Converting source code into machine code is known as *compiling*. It is briefly discussed in Section 13 and covered extensively in CS 241.
The read-only & global data sections

Earlier we described how C “reserves space” in memory for a variable definition. For example:

```c
int n = 0;
```

The location of memory depends on whether the variable is global or local.

First, we discuss global variables.

All global variables are placed in either the read-only data section (constants) or the global data section (mutable variables).
Global variables are available throughout the entire execution of the program, and the space for the global variables is reserved before the program begins execution.

- First, the code from the entire program is scanned and all global variables are identified.
- Next, space for each global variable is reserved.
- Finally, the memory is properly initialized.
- This happens before the main function is called.

The read-only and global memory sections are created and initialized at compile time.
example: read-only & global data

// global

int gmv = 13;
int n = 0;

// read-only

const int c = 42;

Memory Diagram:

```
<table>
<thead>
<tr>
<th>Global</th>
</tr>
</thead>
<tbody>
<tr>
<td>gmv</td>
</tr>
<tr>
<td>n</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Read-Only</th>
</tr>
</thead>
<tbody>
<tr>
<td>c</td>
</tr>
</tbody>
</table>
```
Function Calls (revisited)

Recall the control flow for function calls:

- When a function is called, the program location “jumps” \textit{from} the current location \textit{to} the start of the function

- \texttt{return} changes the program location to go \textit{back} to the most \texttt{recent} calling function (where it “jumped from”)

- C needs to track where it “jumped from” so it knows where to \texttt{return} to

We model function calls with a \texttt{stack}, and store the information in the \texttt{stack section} of memory.
Stacks

A stack in computer science is similar to a physical stack where items are “stacked” on top of each other.

For example, a stack of papers or a stack of plates.

Only the top item on a stack is “visible” or “accessible”.

Items are pushed onto the top of the stack and popped off of the stack.
The call stack

Whenever a function is called, we can imagine that it is pushed onto a stack, and it is now on the top of the stack.

If another function is called, it is then pushed so it is now on top.

The call stack illustrates the “history” or “sequence” of function calls that led us to the current function.

When a function returns, it is popped off of the stack, and the control flow returns to the function now on top.
example: call stack

```c
void blue(void) {
    return;
}

void green(void) {
    // DRAW DIAGRAM
    return;
}

void red(void) {
    green();
    blue();
    return;
}

int main(void) {
    red();
}
```

Call Stack
(when green is called)

```
green
   red
      main
```

When green returns:
* green is popped
* red is now on top
* control flow returns to red
void blue(void) {
    // DRAW DIAGRAM
    return;
}

void green(void) {
    return;
}

void red(void) {
    green();
    blue();
    return;
}

int main(void) {
    red();
}

Call Stack
(when blue is called)

blue
red
main

green does not appear because it was previously popped before blue was called
The return address

When C encounters a return, it needs to know: “where was the program location before this function was called?”

In other words, it needs to “remember” the program location to “jump back to” when returning.

This location is known as the return address.

In this course, we use the name of the calling function and a line number (or an arrow) to represent the return address.

The operating system calls the main function, so that is shown as "OS".
When `foo` is called, the program location is on line 8 of the function `main` so we would record the `return address` as:

`main: 8`
Stack frames

The “entries” pushed onto the call stack are known as *stack frames*.

Each function call creates a *stack frame* (or a “frame of reference”).

Each *stack frame* contains:

- the **argument values**

- all **local variables** (both mutable variables and constants) that appear within the function *block* (including any sub-blocks)

- the **return address** (the program location within the *calling function* to *return to*)
example: stack frames

int pow4(int j) {
    printf("inside pow4\n");
    int k = j * j;
    // DRAW DIAGRAM
    return k * k;
}

int main(void) {
    printf("inside main\n");
    int i = 1;
    printf("%d\n", pow4(i + i));
}
As with Racket, **before** a function can be called, all of the arguments must be values.

**C makes a copy** of each argument value and **places the copy** in the stack frame.

This is known as the “pass by value” convention.
Whereas space for a *global* variable is reserved *before* the program begins execution, space for a *local* variable is only reserved *when the function is called*.

The space is reserved within the newly created stack frame.

When the function returns, the variable (and the entire frame) is popped and effectively “disappears”.

In C, local variables are known as *automatic* variables because they are “automatically” created when needed. There is an `auto` keyword in C but it is rarely used.
Calling a function

We can now model all of the **control flow** when a function is called:

- a *stack frame* is created ("pushed" onto the Stack)
- the current program location is placed in the stack frame as the *return address*
- a *copy* of each of the arguments is placed in the stack frame
- the program location is changed to the start of the new function
- the initial values of local variables are set when their definition is encountered
return

When a function returns:

- the current program location is changed back to the return address (which is retrieved from the stack frame)
- the stack frame is removed ("popped" from the Stack memory area)

The return value (for non-void functions) is stored in a temporary memory area we are not discussing in this course. This is discussed further in CS 241.
Return vs. return address

Beginners often confuse the return address and return statements.

Remember, the return address is a location within the calling function.

It has nothing to do with the location of any return statement(s) in the called function or if one does not exist (e.g., a void function).

There is always one (and only one) return address in a stack frame.
Recursion in C

Now that we understand how stack frames are used, we can see how recursion works in C.

In C, each recursive call is simply a new stack frame with a separate frame of reference.

The only unusual aspect of recursion is that the return address is a location within the same function.

In this example, we also see control flow with the if statement.
int sum_first(int n) {
    if (n == 0) {
        return 0;
    } else {
        return n + sum_first(n - 1);
    }
}

int main(void) {
    int a = sum_first(2);
    //...
}

sum_first:
    n: 0
    return address: sum_first:5

sum_first:
    n: 1
    return address: sum_first:5

sum_first:
    n: 2
    return address: main:10

main:
    a: ???
    return address: OS
Stack section

The call stack is stored in the stack section, the fourth section of our memory model. We refer to this section as “the stack”.

In practice, the “bottom” of the stack (i.e., where the main stack frame is placed) is placed at the highest available memory address. Each additional stack frame is then placed at increasingly lower addresses. The stack “grows” toward lower addresses.

If the stack grows too large, it can “collide” with other sections of memory. This is called “stack overflow” and can occur with very deep (or infinite) recursion.
Uninitialized memory

In most situations, mutable variables *should* be initialized, but C allows variable definitions without any initialization.

```c
int i;
```

For all **global** variables, C automatically initializes the variable to be zero.

Regardless, it is good style to explicitly initialize a global variable to be zero, even if it is automatically initialized.

```c
int g = 0;
```
A local variable (on the stack) that is uninitialized has an arbitrary initial value.

```c
void mystery(void) {
    int k;
    printf("the value of k is: %d\n", k);
}
```

Seashell gives you a warning if you obtain the value of an uninitialized variable.

In the example above, the value of k will likely be a leftover value from a previous stack frame.
## Memory sections (so far)

<table>
<thead>
<tr>
<th>Low</th>
<th>High</th>
</tr>
</thead>
<tbody>
<tr>
<td>Code</td>
<td>Stack</td>
</tr>
<tr>
<td>Read-Only Data</td>
<td></td>
</tr>
<tr>
<td>Global Data</td>
<td></td>
</tr>
<tr>
<td>Heap</td>
<td></td>
</tr>
</tbody>
</table>
Memory snapshot

You may be asked to draw a memory diagram (including the call stack) at a particular moment in the code execution. For example, “draw the memory when line 19 is reached”.

- make sure you show any variables in the global and read-only sections, separate from the stack
- include all local variables in stack frames, including definitions that have not yet been reached (or are incomplete)
- local variables not yet fully initialized have a value of ???
- you do not have to show any temporary storage (e.g., intermediate results of an expression)
When a variable is defined inside of a loop, only one occurrence of the variable is placed in the stack frame. The same variable is re-used for each iteration.

Each time the definition is reached in the loop, the variable is re-initialized (it does not retain its value from the previous iteration).

```c
for (int j = 0; j < 3; ++j) {
    int k = 0;
    k = k + j;
    trace_int(k);
}
```

```
k => 0
k => 1
k => 2
```
Scope vs. memory

Just because a variable exists in memory, it does not mean that it is *in scope*.

**Scope** is part of the C syntax and determines when a variable is “visible” or “accessible”.

**Memory** is part of our C model (which closely matches how it is implemented in practice).
example: snapshot and scope

```c
int foo(void) {
    // SNAPSHOT HERE (at line 2)
    // 5 variables are in memory,
    // but none are in scope
    int a = 1;
    {
        int b = 2;
    }
    return a;
}

const int c = 3;
int d = 4;

int main(void) {
    int e = 5;
    foo();
}
```

**READ-ONLY DATA:**
c: 3

**GLOBAL DATA:**
d: 4

**STACK:**

```
foo:
    a: ???
    b: ???
    return address: main:17
```

```
main:
    e: 5
    return address: OS
```
Model

We now have the tools to model the behaviour of a C program.

At any moment of execution, a program is in a specific *state*, which is the combination of:

- the current *program location*, and
- the current contents of the *memory*.

To properly interpret a program’s behaviour, we must keep track of the program location and all of the memory contents.
Goals of this Section

At the end of this section, you should be able to:

- use the introduced control flow statements, including \( \text{return} \), \( \text{if} \), \( \text{while} \), \( \text{do} \), \( \text{for} \), \( \text{break} \), \( \text{continue} \)

- re-write a recursive function with iteration and \textit{vice versa}

- explain why C has limits on integers and why overflow occurs

- use the \texttt{char} type and explain how characters are represented in ASCII

- use structures in C
• explain how C execution is modelled with memory and control flow, as opposed to the substitution model of Racket

• describe the 4 areas of memory seen so far: code, read-only data, global data and the stack

• identify which section of memory an identifier belongs to

• explain a stack frame and its components (return address, parameters, local variables)

• explain how C makes copies of arguments for the stack frame

• model the execution of small programs by hand, and draw the memory snapshot (including stack frames) at specific execution points