University of Waterloo CS240 - Fall 2022 Assignment 1

Due Date: Wednesday September 21 at 5:00pm

Please follow the guidelines for submission on the course webpage.

https://student.cs.uwaterloo.ca/~cs240/f22/assignments.phtml#guidelines

All logs are base 2. There are 63 marks available; the assignment will be marked out of 60.

Problem 1 [3+3+3+3+3=15 marks]

Provide a complete proof of the following statements from first principles (i.e., using the original definitions of order notation).

- a) $12n^3 + 11n^2 + 10 \in O(n^3)$ d) $1000n \in o(n \log n)$
- **b)** $12n^3 + 11n^2 + 10 \in \Omega(n^3)$ **e)** $n^n \in \omega(n^{20})$
- c) $12n^3 + 11n^2 + 10 \in \Theta(n^3)$

Problem 2 [4+4=8 marks]

For each pair of the following functions, fill in the correct asymptotic notation among Θ , o, and ω in the statement $f(n) \in \sqcup(g(n))$. Provide a brief justification of your answers. In your justification you may use any relationship or technique that is described in class.

a)
$$f(n) = \sqrt{n}$$
 versus $g(n) = (\log n)^2$

b) $f(n) = n^3(5 + 2\cos 2n)$ versus $g(n) = 3n^2 + 4n^3 + 5n$

Problem 3 [6+6=12 marks]

Prove or disprove each of the following statements. To prove a statement, you should provide a formal proof that is based on the definitions of the order notations. To disprove a statement, you can either provide a counter example and explain it or provide a formal proof. All functions are positive functions.

a)
$$f(n) \notin o(g(n))$$
 and $f(n) \notin \omega(g(n)) \Rightarrow f(n) \in \Theta(g(n))$
b) $\min(f(n), g(n)) \in \Theta\left(\frac{f(n)g(n)}{f(n)+g(n)}\right)$

Problem 4 [6 marks]

Suppose n is a power of two and θ is a parameter in the range $2 \le \theta \le 3$. Derive an exact closed form for the sum

$$f(n) := \sum_{i=0}^{\log_2 n} 4^i \left(\frac{n}{2^i}\right)^{\theta}$$

in terms of n and θ . *Hint:* Re-write the formula as a geometric series, and treat $\theta = 2$ as a special case.

Problem 5 [2+2+4+4=12 marks]

Consider the following procedure.

```
pre: n is a positive integer
pre: v[1..n] is a binary vector of length n,
     i.e., each entry is either 0 or 1
foo(v,n)
     i := 1;
1.
2.
     while i<=n and v[i]=0 do
3
         i := i+1
4
     od;
     for j from 1 to i do
5.
6.
         print("Hello world!")
7.
     od;
```

- a) How many inputs are there are of size n?
- b) What is the worst case number of calls to print? Give an exact formula in terms of n and justify your answer by giving an example of a worst case input of size n.
- c) For $i \in \{1, 2, ..., n\}$, let S_i denote the subset of inputs of size n for which the number of calls to print is i. Describe what an element of S_i looks like, and derive an expression for $|S_i|$, the number of elements of S_i .
- d) What is the average case number of calls to print? Derive an exact closed form formula in terms of n.

Problem 6 [5 marks]

Prove that the following code fragment will always terminate.

```
s := 3*n // n is an integer
while (s>0)
    if (s is even)
```

```
s := floor(s/4)
else
s := 2*s
```

Problem 7 [5 marks]

Analyze the following piece of pseudo-code and give a Θ bound on the running time as a function of n. Show your work. A formal proof is not required, but you should justify your answer.

| 1. | $mystery \leftarrow 0$ |
|----|---------------------------------------|
| 2. | for $i \leftarrow 1$ to $3n$ do |
| 3. | $mystery \leftarrow mystery \times 4$ |
| 4. | for $j \leftarrow 1388$ to 2010 do |
| 5. | for $k \leftarrow 4i$ to $6i$ do |
| 6. | $mystery \leftarrow mystery + k$ |
| | |