University of Waterloo CS240 Winter 2024 Assignment 1

Due Date: Tuesday, January 23 at 5:00pm

Please read https://student.cs.uwaterloo.ca/~cs240/w24/assignments.phtml#guidelines for guidelines on submission. Each question must be submitted individually to MarkUs as a PDF with the corresponding file names: a1q1.pdf, a1q2.pdf, ..., a1q6.pdf. It is a good idea to submit questions as you go so you aren't trying to create several PDF files at the last minute.

Late Policy: Assignments are due at 5:00pm, with the grace period until 11:59pm.

Notes:

- Logarithms are in base 2, if not mentioned otherwise.
- A positive function is a function that takes positive real values.

Question 1 [3+3+3+3+3=15 marks]

Provide a complete proof of the following statements from first principles (i.e., using the original definitions of order notation).

- a) $10n^5 \log n + 42n^4 15n^2 \in O(n^5 \log n)$
- **b)** $2n^4 10n^2 + 5n \in \Omega(n^4)$
- c) $1 + (n^2 \mod 10) \in \Theta(1)$
- **d)** $3^n \in o(3.1^n)$
- e) $n^2 \in \omega(n^{1.9})$

Question 2 [2+4=6 marks]

Prove the following statements based on the definitions of the order notations. All functions take positive values.

- a) If $f(n) \in \omega(1)$, then $f(n)h(n) \in \omega(h(n))$.
- **b)** If $f(n) \in \Theta(g(n)), h(n) \in \Theta(r(n))$, and $g(n) \in o(r(n))$ then $f(n) \in o(h(n))$.

Question 3 [3+3+3=9 marks]

For each pair of the following functions, fill in the correct asymptotic notation among Θ , o, and ω in the statement $f(n) \in \sqcup(g(n))$. Prove the relationship using any relationship or technique described in class.

a) $f(n) = 5n + 7n^2 \log n$ versus $g(n) = n^2 \sqrt{n}$

b)
$$f(n) = n!$$
 versus $g(n) = n^{100}$

c) $f(n) = 2^n$ versus $g(n) = 2^{\cos n}$

Question 4 [3+3+3+3=12 marks]

Analyze the following pieces of pseudocode and give a tight (Θ) bound on the running time as a function of n. Show your work. A formal proof is not required, but you should justify your answer (in all cases, n is assumed to be a positive integer).

```
a) s := 0
   for i=1 to n * n * n do
      for j=1 to i * i * i * i do
         s := s + 1
b) p = 1
   s = 0
   for i = 1 to n do
      p = p * 2
      for j = 1 to p do
         s = s + 1
c) s = 0
   i = 1
   while i <= n do
       j = n * n * n * n
       while j > 0 do
           s = s + 1
           j = j - i
       i = i * 2
```

d) For this question, you can derive O bound, but it must be as tight as possible.

```
foo(A[0...n-1]) # assume n>=1 is a power of two
    if n <= 1
        return
    print("x")
    foo(A[0...n/4-1]
        foo(A[n/4...n/2-1]</pre>
```

Question 5 [4+4=8 marks]

True or false? For each of the following assertions, indicate whether it is true or false. If true, prove it; if false, give a counter-example and briefly justify it.

a) f(n) is $\Theta(f(n/2))$, where f takes positive values.

b)
$$1 + n + n^2 + \dots + n^{k-1} \in \Theta(n^k).$$

Question 6 [1+4=5 marks]

For this problem, assume that we only deal with functions that are defined on non-negative integers and take positive real values (i.e. the domain is non-negative integers and the range is positive reals). Dr. I. M. Smart has recently invented a new class of functions, denoted O'(f). A function g(n) is in O'(f) if there is a constant c > 0 such that $g(n) \le cf(n)$ for all $n \ge 0$. Prove the following statements about relationship between O(f) and O'(f).

- a) If $g(n) \in O'(f(n))$, then $g(n) \in O(f(n))$
- **b)** If $g(n) \in O(f(n))$, then $g(n) \in O'(f(n))$