

Tutorial 10: March 25

1. **Karp-Rabin**

For Karp-Rabin pattern matching, consider the following hash function for strings over the alphabet $\{A, C, G, T\}$:

$$h(P) = (\# \text{ of occurrences of } A) + 2 \times (\# \text{ of occurrences of } C) + 3 \times (\# \text{ of occurrences of } G) + 4 \times (\# \text{ of occurrences of } T)$$

Given the pattern $P = \text{TAGCAT}$ and sequence $T = \text{TGCCGATGTAGCTAGCAT}$, use the table below to show all the character comparisons performed during Karp-Rabin pattern matching. Start a new pattern shift (in which character comparison occurs) in a new row. You may not need all the available space.

T	G	C	C	G	A	T	G	T	A	G	C	T	A	G	C	A	T

Table 1: Table for Karp-Rabin problem.

2. **Boyer-Moore, Revisited**

Let $P[0..5] = \text{payday}$ and let $T[0..12] = \text{daypayplayayaya}$.

- a) Compute the last-occurrence array $S[0..5]$ for the pattern P .
- b) Show the search for P in T using the Boyer-Moore algorithm using only the bad character heuristic. Also, put square brackets around characters that are known to be matched, even if the algorithm matches them again.

3. **Moore's first name is 'J', not abbreviated**

Consider using the Boyer-Moore algorithm with only the Bad Character heuristic to search for a pattern P of length m in a text T of length n , with $n > m$, where P does **not** appear in T .

- a) Give an example of a pattern P with length n and text T with length n that achieves the worst-case runtime for searching. Do not consider preprocessing time.
- b) Same question, but for the best-case runtime.