

University of Waterloo
CS240E, Winter 2025
Written Assignment 5 Post-Mortem

This document goes over common errors and general performance on the assignment. We create it using feedback from the graders, and it is meant to be used as a resource to understand common areas that we can improve in.

Question 1 [3+3=6 marks]

- Part (a) was done well.
- In part (b), some solutions argued along the lines of the Boyer-Moore pseudocode and last-occurrence array. It is possible to do this correctly, but some solutions incorrectly concluded that the shift amount is always m . This can't be true for every text T , so be sure to “sanity-check” your answers in cases like these.

Question 2 [4 marks]

- This question was done well.

Question 3 [2+2+6=10 marks]

- Parts (a) and (b) were done well.
- For part (c), some solutions stated that the Huffman trie for S would be a full binary tree without justifying why rigorously.

Question 4 -

Question 5 [2+2+3+2+2+5+2=18 marks]

- Parts (a) and (b)(i) were done well.
- In part (b)(ii), some solutions didn't account for decoding the first $m - 1$ characters of T_2 before searching through the indices of T_1 for P .
- In part (b)(iii), many solutions made the subtle error of forgetting that one character from T_2 is added to $\mathcal{W}$. Handling this case required using the assumption that $|P| > 1$.
- Part (b)(iv) was done well, with many solutions offering humorous remarks about the size of their constants.

- In part (b)(v), some solutions did not fully explain why the number of tuples was bounded by a constant, or why $P_2, \dots, P_{k-1}$ are uniquely determined by P_1 .
- In part (b)(vi), a few solutions forgot to use parts (i) and (ii) to handle the case where the pattern overlaps T_1 .

Question 6 [2+2+5=9 marks]

- In part (a), some solutions had an off-by-one error in the suffix array, perhaps due to misremembering the definition. This was not penalized, but be sure to get the indices right on the final!
- Part (b) was done well.
- In part (c), many solutions tried to maintain a set of suffixes of T that are candidates for a prefix match with P . While it is possible to achieve this efficiently, most of these solutions had a runtime of $\Omega(mN)$.