CS341: ALGORITHMS (W21)

Lecture 1: course overview and Bentley’s problem

Readings: CLRS Chapter 1

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TABLE OF CONTENTS

- Course mechanics
- **Overview** of course material
- Worked example: Bentley’s problem
  - Multiple solutions, demonstrating **different algorithm design techniques**
COURSE MECHANICS
COURSE MECHANICS

- **Course website:** https://www.student.cs.uwaterloo.ca/~cs341/
  - Syllabus, calendar, policies, slides, assignments...
  - Read this and **mark** important dates.
- **Keep up with the lectures:** Material **builds** over time...
- **Piazza:** For questions and announcements.
ASSESSMENTS

- All sections have same assignments, midterm and final
  - Notify us long before the deadline of severe problems that will cause you to miss an assignment
- Midterm and final are to be take-home exams
- See website for grading scheme
TEXTBOOK

- Introduction to Algorithms, Third Edition
  Cormen, Leiserson, Rivest and Stein
  - Available **for free** via library website!
- You are expected to know
  - entire textbook sections, as listed on course website
  - **all the material presented in lectures**
    (unless we explicitly say you aren’t responsible for it)
ACADEMIC OFFENSES

- Beware plagiarism
  - **High level discussion** about solutions with individual students is **OK**
  - Don’t take written notes away from such discussions
  - Class-wide discussion of solutions is **not** OK (until the deadline)
COURSE OVERVIEW
Sketching out the road ahead
WHY IS CS341 IMPORTANT FOR YOU?

- Algorithms is the heart of CS
  - It appears often in later courses
  - It dominates technical interviews
    - Master this material... make your interviews easy!
  - Designing algorithms is creative work
  - Useful for some of the more interesting jobs out there
  - And, you want to graduate...

CS 341 is a required course for all CS
WHAT IS A COMPUTATIONAL PROBLEM?

- Informally: A description of input, and the desired output

WHAT IS AN ALGORITHM?

- Informally: A well-defined procedure (sequence of steps) to solve a computational problem
### Examples of Computational Problems

<table>
<thead>
<tr>
<th></th>
<th>Sorting</th>
<th>Matrix Multiplication</th>
<th>Traveling Salesman Problem</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Input</strong></td>
<td>An array of integers (in arbitrary order)</td>
<td>Two $n \times n$ matrices $A$, $B$</td>
<td>A set $S$ of cities, and distances between each pair of cities</td>
</tr>
<tr>
<td><strong>Desired output</strong></td>
<td>Same array of integers in <strong>increasing</strong> order</td>
<td>A matrix $C = A \times B$</td>
<td>Shortest possible path that visits each city, and returns to the origin city</td>
</tr>
</tbody>
</table>

Input:

- An array of integers (in arbitrary order)
- Two $n \times n$ matrices $A$, $B$
- A set $S$ of cities, and distances between each pair of cities

Desired output:

- Same array of integers in **increasing** order
- A matrix $C = A \times B$
- Shortest possible path that visits each city, and returns to the origin city
ANALYSIS OF ALGORITHMS

- Every software program uses resources
  - CPU instructions → we call this time
  - Memory (RAM) → we call this space
  - Others: I/O, network bandwidth/messages, locks... (not covered in this course)

Analysis is the study of **how many** resources an algorithm uses
  - Usually using big-O notation (to ignore constant factors)
TAXONOMY OF ALGORITHMS

- Serial vs Parallel
  - Serial: One instruction at a time
  - Parallel: Multiple instructions at once

- Deterministic vs Randomized
  - D: On multiple runs on same input, always do same thing
  - R: On multiple runs on same input, may do different things

  Example: flip a coin, and base your next action on the result

- Exact vs Approximate
  - Exact: exact solution to the problem
  - Approximate: produce something “close” to a solution

This course mainly covers: Serial, deterministic, exact
TRACTABILITY: DO ALL PROBLEMS HAVE FAST SOLUTIONS?

- For some problems, such as the traveling salesman problem, we have only found **exponential time** algorithms.
  - These algorithms take **exponentially longer** to solve the problem as the number of cities increases!
  - Informally: adding one city **doubles** the runtime
  - This severely limits our ability to solve “real world” inputs…
- Is there a way around this limitation? Or should we stop trying?
- Open question (P vs NP): is it **possible** to solve such problems in polynomial time?
Topics to Cover

Fundamental (& Fast) Algorithms for Tractable Problems
- MergeSort
- Strassen’s MM
- BFS/DFS
- Dijkstra’s SSSP
- MST (Kruskal or Prim)
- Floyd Warshall APSP
- Topological Sort
- ...

Common Algorithm Design Paradigms
- Divide-and-Conquer
- Greedy
- Dynamic Programming
- Exhaustive search / brute force

Mathematical Tools to Analyze Algorithms
- Big-oh notation
- Recursion Tree
- Master method
- Substitution method
- Exchange Arguments
- Greedy-stays-ahead Arguments

Intractable Problems
- P vs NP
- Poly-time Reductions
- Undecidability
CS341: Before → After

1. Fundamental Algorithms
2. Fundamental Design Paradigms
3. Tractability/Intractability

Math Techniques for Algorithm Analysis
BENTLEY’S PROBLEM

A worked example to demonstrate algorithm design
Bentley’s Problem (introductory example)

Given an array of \( n \) integers, \( A[1], \ldots, A[n] \), find the maximum sum of consecutive entries of \( A \) (return 0 if all entries of \( A \) are negative).

**Example 1**

Array index:

| 1 | 7 | 4 | 0 | 2 | 1 | 3 | 1 |

Solution: 19
(take all of \( A[1..8] \))

**Example 2**

Index:

| -1 | -7 | -4 | -1 | -2 | -1 | -3 | -1 |

Solution: 0
(take no elements of \( A \))

**Example 3**

Index:

| 1 | -7 | 4 | 0 | 2 | -1 | 3 | -1 |

Solution: 8
(take \( A[3..7] \))
Bentley’s Problem: Solution 1

max := 0;
for i := 1 to n do
    for j := i to n do
        // compute A[i] + ... + A[j]
        sum := 0;
        for k := i to j do
            sum := sum + A[k];
        // compare to maximum sum observed so far
        if sum > max then max := sum;

Try all combinations of i, j
And for each combination, sum over k = i .. j

Time:
max := 0;
for i := 1 to n do
    // for each j, compute A[i] + ... + A[j]
    sum := 0;
    for j := i to n do
        // update sum by adding the next entry A[j]
        sum := sum + A[j];
        // compare to maximum sum observed so far
        if sum > max then max := sum;

Avoid summing over $k = i \ldots j$
Bentley’s Problem: Solution 3

Divide-and-Conquer can also be used here:
Divide an array into two equally-sized parts.
Our solution must either be entirely in the left part, or entirely in the right part, or it must be crossing the partition line.

Case 1: optimal sol’n is entirely in L

Case 2: optimal sol’n is entirely in R

Case 3: optimal sol’n crosses the partition
Therefore: Find the maximum subarray for left part \((maxL)\) and right part \((maxR)\) (done by recursive call).
Find the maximum subarray ”going over the middle partition line” \((maxM)\).
This can be done in linear time \(\Theta(n)\).
The solution is \(\max maxL, maxR, maxM\).

**Find:** maximum subarray **going over the middle** partition

Find \(i\) that maximizes the sum over \(i \ldots n/2\)

Find \(j\) that maximizes the sum over \(\left(\frac{n}{2} + 1\right) \ldots j\)

We can prove \(A[i \ldots j]\) **is** the maximum subarray going over the middle partition!
WHY $A[i \ldots j]$ IS MAXIMAL

- Suppose not for contradiction
- Then some $A[i' \ldots j']$ that crosses the partition has a larger sum

But both are impossible!

This sum is bigger

So either $\sum L' > \sum L$ or $\sum R' > \sum R$
function solveDnC(A)
    let n = sizeof(A)

    // base case
    if n == 1 then return max(0, A[1])

    // recursive case
    maxL = solveDnC(A[1 .. n/2])
    maxR = solveDnC(A[n/2+1 .. n])

    // compute maxM
    tempSum = 0
    maxI = 0
    for i = n/2 .. 1
        tempSum = tempSum + A[i]
        if tempSum > maxI then maxI = tempSum

    tempSum = 0
    maxJ = 0
    for j = n/2+1 .. n
        tempSum = tempSum + A[j]
        if tempSum > maxJ then maxJ = tempSum

    maxM = maxI + maxJ
    return max( maxL, maxR, maxM )
function solveDnC(A)
    let n = sizeof(A)

    // base case
    if n == 1 then return max(0, A[1])

    // recursive case
    maxL = solveDnC(A[1 .. n/2])
    maxR = solveDnC(A[n/2+1 .. n])

    // compute maxM
    tempSum = 0
    maxI = 0
    for i = n/2 .. 1
        tempSum = tempSum + A[i]
        if tempSum > maxI then maxI = tempSum
    tempSum = 0
    maxJ = 0
    for j = n/2+1 .. n
        tempSum = tempSum + A[j]
        if tempSum > maxJ then maxJ = tempSum
    maxM = maxI + maxJ

    return max( maxL, maxR, maxM )

A
1  -7  4  0
2  -1  3  0

L
1  -7
4  0

R
2  -1
3  0

maxL = 4
maxR = 4
maxM = maxL + maxJ = 8
maxI = 4
maxJ = 4

Return max( 4, 4, 8 ) = 8
How do we analyze this running time? Need new mathematical techniques!

Recurrence relations, recursion tree methods, master theorem...

This result is really quite good... but can we do asymptotically better?
Define: \( \text{include}(j) = \text{maximum sum} \) of consecutive entries in array \( A[1..j] \)
if the sum must \textbf{include} \( A[j] \)

Define: \( \text{exclude}(j) = \text{maximum sum} \) of consecutive entries in array \( A[1..j] \)
if the sum must \textbf{exclude} \( A[j] \)

Observe: if we could solve for \( \text{include}(j), \text{exclude}(j) \) for all \( j \),
then the solution to our problem would be \( \max\{ \text{include}(n), \text{exclude}(n) \} \)
We can define **recurrence relations** to solve for **include** and **exclude**

- **Base case**: \(\text{include}(1) = A[1]\)
- **Base case**: \(\text{exclude}(1) = 0\)

\[
\text{include}(j) = \max\{ A[j], A[j] + \text{include}(j - 1) \}
\]

\[
\text{exclude}(j) = \max\{ \text{include}(j - 1), \text{exclude}(j - 1) \}
\]

"Max sum in A[1..1] if we must include A[1]"

If including \(A[j]\), there are two possibilities: either start a **new** sum of consecutive entries at \(A[j]\), or **extend** the best sum that ends at \(A[j - 1]\)

If excluding \(A[j]\), the best we can do in \(A[1..j]\) is simply the best we can do in \(A[1..j - 1]\)
Example: computing these recurrence relations with two arrays

- **Base case:** $\text{include}(1) = A[1]$
- $\text{include}(j) = \max\{ A[j], A[j] + \text{include}(j - 1) \}$

<table>
<thead>
<tr>
<th>Index</th>
<th>(A)</th>
<th>1</th>
<th>-7</th>
<th>4</th>
<th>0</th>
<th>2</th>
<th>-1</th>
<th>3</th>
<th>-9</th>
</tr>
</thead>
<tbody>
<tr>
<td>include</td>
<td></td>
<td>1</td>
<td>-6</td>
<td>4</td>
<td>4</td>
<td>6</td>
<td>5</td>
<td>8</td>
<td>-1</td>
</tr>
</tbody>
</table>

- **Base case:** $\text{exclude}(1) = 0$
- $\text{exclude}(j) = \max\{ \text{include}(j - 1), \text{exclude}(j - 1) \}$

<table>
<thead>
<tr>
<th>Index</th>
<th>(A)</th>
<th>0</th>
<th>1</th>
<th>1</th>
<th>4</th>
<th>4</th>
<th>6</th>
<th>6</th>
<th>8</th>
</tr>
</thead>
<tbody>
<tr>
<td>exclude</td>
<td></td>
<td>0</td>
<td>1</td>
<td>1</td>
<td>4</td>
<td>4</td>
<td>6</td>
<td>6</td>
<td>8</td>
</tr>
</tbody>
</table>

Recall the definition:

- $\text{include}(1) = \text{“max solution in } A[1..1] \text{ that includes } A[1]...\text{”}$
- $\text{include}(2) = \text{“max solution in } A[1..2] \text{ that includes } A[2]...\text{”}$
- $\text{include}(3) = \text{“max solution in } A[1..3] \text{ that includes } A[3]...\text{”}$

- $\text{exclude}(1) = \text{“max solution in } A[1..1] \text{ that excludes } A[1]...\text{”}$
- $\text{exclude}(2) = \text{“max solution in } A[1..2] \text{ that excludes } A[2]...\text{”}$
- $\text{exclude}(3) = \text{“max solution in } A[1..3] \text{ that excludes } A[3]...\text{”}$

Full solution is $\max$ of these two: 8
Base case: $\text{exclude}(1) = 0$; $\text{include}(1) = A[1]$

Recursive case:

- $\text{exclude}(j) = \max\{ \text{include}(j-1), \text{exclude}(j-1) \}$
- $\text{include}(j) = \max\{ A[j], A[j] + \text{include}(j-1) \}$

Let's turn these recurrences into code…

```python
1 function solveDP(A)
2     define arrays exclude[1..n], include[1..n]
3
4     exclude[1] = 0
6     for j = 2..n
7         exclude[j] = max( include[j-1], exclude[j-1] )
8         include[j] = max( A[j], A[j] + include[j-1] )
9
10 return max(exclude[n], include[n])
```

Do we actually need these entire arrays? Only really care about the last entry of each…
function solveDP(A)
    define arrays exclude[1..n], include[1..n]

    exclude[1] = 0
    for j = 2..n
        exclude[j] = max( include[j-1], exclude[j-1] )
        include[j] = max( A[j], A[j] + include[j-1] )
    return max( exclude[n], include[n] )

Same running time, but only \(O(1)\) space (besides the input array)

function solveDP_spaceEfficient(A)
    exclude = 0
    include = A[1]
    for j = 2..n
        exclude = max( include, exclude )
        include = max( A[j], A[j] + include )
    return max( exclude, include )

At this time, include contains exactly “include[j-1]”

And similarly for exclude...

And these contain exactly “exclude[n]” and “include[n]”
BENTLEY’S PROBLEM: TIME CONSTRAINTS

- Consider solutions implemented in C
- Some values measured (on a Pentium II)
- Some estimated from other measurements
- $\epsilon$ represents time under 0.01s

<table>
<thead>
<tr>
<th>Time to solve a problem of size:</th>
<th>Sol.4 $\Theta(n)$</th>
<th>Sol.3 $\Theta(n \log n)$</th>
<th>Sol.2 $\Theta(n^2)$</th>
<th>Sol.1 $\Theta(n^3)$</th>
</tr>
</thead>
<tbody>
<tr>
<td>10</td>
<td>$\epsilon$</td>
<td>$\epsilon$</td>
<td>$\epsilon$</td>
<td>$\epsilon$</td>
</tr>
<tr>
<td>50</td>
<td>$\epsilon$</td>
<td>$\epsilon$</td>
<td>$\epsilon$</td>
<td>$\epsilon$</td>
</tr>
<tr>
<td>100</td>
<td>$\epsilon$</td>
<td>$\epsilon$</td>
<td>$\epsilon$</td>
<td>$\epsilon$</td>
</tr>
<tr>
<td>1000</td>
<td>$\epsilon$</td>
<td>0.01s</td>
<td>2.1s</td>
<td>4.5s</td>
</tr>
<tr>
<td>10000</td>
<td>$\epsilon$</td>
<td>0.12s</td>
<td>3.5m</td>
<td>75m</td>
</tr>
<tr>
<td>100000</td>
<td>0.04s</td>
<td>0.12s</td>
<td>3.5m</td>
<td>75m</td>
</tr>
<tr>
<td>1 mil.</td>
<td>0.42s</td>
<td>1.4s</td>
<td>5.8h</td>
<td>142yrs.</td>
</tr>
<tr>
<td>10 mil.</td>
<td>4.2s</td>
<td>16.1s</td>
<td>24.3d</td>
<td>140000yrs.</td>
</tr>
<tr>
<td>Max size problem solved in</td>
<td>1s</td>
<td>2.3 mil.</td>
<td>740000</td>
<td>6900</td>
</tr>
<tr>
<td></td>
<td>1m</td>
<td>140 mil.</td>
<td>34 mil.</td>
<td>53000</td>
</tr>
<tr>
<td></td>
<td>1d</td>
<td>200 bil.</td>
<td>35 bil.</td>
<td>2 mil.</td>
</tr>
<tr>
<td>time if $n$ increases:</td>
<td>x 2</td>
<td>x 2</td>
<td>x 2+</td>
<td>x 4</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
HOW ABOUT A MORE MODERN SYSTEM? 😊
<table>
<thead>
<tr>
<th>N</th>
<th>Sol.4</th>
<th>Sol.3</th>
<th>Sol.2</th>
<th>Sol.1</th>
</tr>
</thead>
<tbody>
<tr>
<td>100</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>1,000</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0.12</td>
</tr>
<tr>
<td>10,000</td>
<td>0</td>
<td>0</td>
<td>0.036</td>
<td>2 min</td>
</tr>
<tr>
<td>100,000</td>
<td>0</td>
<td>0.002</td>
<td>3.582</td>
<td>33 hrs</td>
</tr>
<tr>
<td>1M</td>
<td>0.001</td>
<td>0.017</td>
<td>6 min</td>
<td>4 yrs</td>
</tr>
<tr>
<td>10M</td>
<td>0.012</td>
<td>0.195</td>
<td>12 hrs</td>
<td>3700 yrs</td>
</tr>
<tr>
<td>100M</td>
<td>0.112</td>
<td>2.168</td>
<td>50 days</td>
<td>3.7M yrs</td>
</tr>
<tr>
<td>1 billion</td>
<td>1.124</td>
<td>24.57</td>
<td>1.5 years</td>
<td>&gt; age of life</td>
</tr>
<tr>
<td>10 billion</td>
<td>19.15</td>
<td>5 min</td>
<td>150 years</td>
<td>&gt; age of universe</td>
</tr>
</tbody>
</table>

### AMD Threadripper 3970x (2020)

- Sol.4
- Sol.3
- Sol.2
- Sol.1

### Pentium II (circa 1997)

<table>
<thead>
<tr>
<th>Sol.4</th>
<th>Sol.3</th>
<th>Sol.2</th>
<th>Sol.1</th>
</tr>
</thead>
<tbody>
<tr>
<td>ε</td>
<td>ε</td>
<td>ε</td>
<td>ε</td>
</tr>
<tr>
<td>ε</td>
<td>0.01s</td>
<td>2.1s</td>
<td>75m</td>
</tr>
<tr>
<td>0.04s</td>
<td>0.12s</td>
<td>3.5m</td>
<td>52d</td>
</tr>
<tr>
<td>0.42s</td>
<td>1.4s</td>
<td>5.8h</td>
<td>142 yrs</td>
</tr>
<tr>
<td>4.2s</td>
<td>16.1s</td>
<td>24.3d</td>
<td>140000 yrs</td>
</tr>
</tbody>
</table>

- ε
- 0.02s
- 4.5s
- 0.001
- 0.12s
- 3.5m
- 52d
- 0.112
- 2.168
- 3.7M
- 1.124
- 24.57
- > age of life
- 19.15
- 5 min
- > age of universe
BONUS

- Trevor’s study-song of the day
- Tool - Descending
- [youtu.be/PcSoLwFisaw](https://youtu.be/PcSoLwFisaw)