CS 341: ALGORITHMS

Lecture 2: background and analysis
Readings: CLRS Chapters 2.1, 2.2, 3

Trevor Brown (co-taught with Anna Lubiw)
https://www.student.cs.uwaterloo.ca/~cs341

trevor.brown@uwaterloo.ca
When your interviewer asks for the time complexity of your algorithm but you have no idea what that means

Big O notation

True story
O-notation:

\[ f(n) \in O(g(n)) \] if there exist constants \( c > 0 \) and \( n_0 > 0 \) such that \( 0 \leq f(n) \leq cg(n) \) for all \( n \geq n_0 \).

Here the complexity of \( f \) is not higher than the complexity of \( g \).
Ω-notation:

\[ f(n) \in \Omega(g(n)) \text{ if there exist constants } c > 0 \text{ and } n_0 > 0 \text{ such that } 0 \leq cg(n) \leq f(n) \text{ for all } n \geq n_0. \]

Here the complexity of \( f \) is not lower than the complexity of \( g \).
\[ f(n) \in \Theta(g(n)) \text{ if there exist constants } c_1, c_2 > 0 \text{ and } n_0 > 0 \text{ such that } 0 \leq c_1 g(n) \leq f(n) \leq c_2 g(n) \text{ for all } n \geq n_0. \]

Here \( f \) and \( g \) have the same complexity.
\textbf{o-notation:}

\( f(n) \in o(g(n)) \) if for all constants \( c > 0 \), there exists a constant \( n_0 > 0 \) such that \( 0 \leq f(n) \leq cg(n) \) for all \( n \geq n_0 \).

Here \( f \) has lower complexity than \( g \).

\textbf{ω-notation:}

\( f(n) \in \omega(g(n)) \) if for all constants \( c > 0 \), there exists a constant \( n_0 > 0 \) such that \( 0 \leq cg(n) \leq f(n) \) for all \( n \geq n_0 \).

Here \( f \) has higher complexity than \( g \).
EXERCISE

• Which of the following are true?
• \( n^2 \in O(n^3) \)
• \( n^2 \in o(n^3) \)
• \( n^3 \in \omega(n^3) \)
• \( \log n \in o(n) \)
• \( n \log n \in \Omega(n) \)
• \( n \log n^2 \in \omega(n \log n) \)
• \( n \in \Theta(n \log n) \)
EXERCISE

• Which of the following are true?
  • $n^2 \in O(n^3)$  YES
  • $n^2 \in o(n^3)$  YES
  • $n^3 \in \omega(n^3)$  NO
  • $\log n \in o(n)$  YES
  • $n \log n \in \Omega(n)$  YES
  • $n \log n^2 \in \omega(n \log n)$  NO
  • $n \in \Theta(n \log n)$  NO
Intuitively, we have the following correspondences between order notation and growth rates:

- $f(n) \in O(g(n))$ means the growth rate of $f$ is $\leq$ the growth rate of $g$
- $f(n) \in o(g(n))$ means the growth rate of $f$ is $<$ the growth rate of $g$
- $f(n) \in \Omega(g(n))$ means the growth rate of $f$ is $\geq$ the growth rate of $g$
- $f(n) \in \omega(g(n))$ means the growth rate of $f$ is $>$ the growth rate of $g$
- $f(n) \in \Theta(g(n))$ means the growth rate of $f$ is $=$ the growth rate of $g$

- $4n \in O(n^2)$
- $4n \in o(n^2)$
- $7n^2 \in O(n^2)$
- $7n^2 \notin o(n^2)$
- $7n^2 \in \Omega(n)$
- $7n^2 \in \omega(n)$
- $4n \in \Omega(n)$
- $4n \notin \omega(n)$

This is included for your notes.
Relationships between Order Notations

\[ f(n) \in \Theta(g(n)) \iff g(n) \in \Theta(f(n)) \]
\[ f(n) \in O(g(n)) \iff g(n) \in \Omega(f(n)) \]
\[ f(n) \in o(g(n)) \iff g(n) \in \omega(f(n)) \]

\[ f(n) \in \Theta(g(n)) \iff f(n) \in O(g(n)) \text{ and } f(n) \in \Omega(g(n)) \]
\[ f(n) \in o(g(n)) \Rightarrow f(n) \in O(g(n)) \]
\[ f(n) \in \omega(g(n)) \Rightarrow f(n) \in \Omega(g(n)) \]
Prove that $f(n) \in \Theta(g(n))$ implies $g(n) \in \Theta(f(n))$.

**Proof:** Suppose $f(n) \in \Theta(g(n))$. Then there exist constants $c_1, c_2, n_0$ such that

$$0 \leq c_1 g(n) \leq f(n) \leq c_2 g(n)$$

if $n \geq n_0$. Thus

$$0 \leq (1/c_2)f(n) \leq g(n) \leq (1/c_1)f(n)$$

if $n \geq n_0$. Define $c'_1 = 1/c_2$, $c'_2 = 1/c_1$ and $n'_0 = n_0$. Then

$$0 \leq c'_1 f(n) \leq g(n) \leq c'_2 f(n)$$

if $n \geq n'_0$. 
1. Let \( f(n) = n^2 - 7n - 30 \). Prove from first principles that \( f(n) \in O(n^2) \).

2. Let \( f(n) = n^2 - 7n - 30 \). Prove from first principles that \( f(n) \in \Omega(n^2) \).

3. Suppose \( f(n) = n^2 + n \). Prove from first principles that \( f(n) \notin O(n) \).
EXAMPLE 1: $f(n) = n^2 - 7n - 30$

- Want to prove (WTP) **from first principles**: $f(n) \in O(n^2)$
  - More formally: there exist constants $c > 0, n_0 > 0$ such that for all $n \geq n_0$, we have $0 \leq f(n) \leq cn^2$
- Pick a value for $c$. How about $1$?
- Let’s visualize $c = 1$

Seems plausible that $c = 1$ will work

Let’s prove this algebraically
EXAMPLE 1: \( f(n) = n^2 - 7n - 30 \)

- WTP: there exist constants \( c > 0, n_0 > 0 \) such that for all \( n \geq n_0 \), we have \( 0 \leq f(n) \leq cn^2 \)

- Solution:
  - **When is** \( f(n) \leq cn^2 \)
    - \( n^2 - 7n - 30 \leq n^2 \) (for all \( n \geq 0 \))
  - **When is** \( 0 \leq f(n) \)
    - \( f(n) = n^2 - 7n - 30 = (n - 10)(n + 3) \)
    - When is \( (n - 10)(n + 3) \geq 0 \)? When \( n \geq 10 \).
    - (or when \( n \leq -3 \) ... but we want \( n_0 > 0 \))

So, the claim holds with \( c = 1, n_0 = 10 \)
EXAMPLE 2: \( f(n) = n^2 - 7n - 30 \)

- **WTP from first principles:** \( f(n) \in \Omega(n^2) \)
  - More formally: there exist constants \( c > 0, n_0 > 0 \) such that for all \( n \geq n_0 \), we have \( 0 \leq cn^2 \leq f(n) \)

- **Solution:**
  - Pick a value for \( c \).
  - How about 1?
  - Must show \( n^2 \leq n^2 - 7n - 30 \)
  - Impossible! \( c = 1 \) is too large.
  - Let’s try \( c = \frac{1}{2} \).
EXAMPLE 2: $f(n) = n^2 - 7n - 30$

- WTP: there exist constants $c > 0, n_0 > 0$ such that for all $n \geq n_0$, we have $0 \leq cn^2 \leq f(n)$
- Solution:
  - Let’s try $c = \frac{1}{2}$.
  - Goal: show $0 \leq \frac{1}{2}n^2 \leq n^2 - 7n - 30$
  - First part $0 \leq \frac{1}{2}n^2$ is easy: satisfied for all $n \geq 0$ (i.e., for any $n_0$).
  - Second part $\frac{1}{2}n^2 \leq n^2 - 7n - 30$ holds when $\frac{1}{2}n^2 - 7n - 30 \geq 0$
  - Roots are $7 \pm \sqrt{109}$, which are $< 18$

Result: $c = \frac{1}{2}, \ n_0 = 18$ works!
EXAMPLE 3: \( f(n) = n^2 + n \)

- **WTP from first principles** \( f(n) \notin O(n) \). Formally:
  
  \[
  \neg (f(n) \in O(n))
  \]
  
  \[
  \neg (\exists c > 0, n_0 > 0 \quad \forall n \geq n_0 : 0 \leq f(n) \leq cn)
  \]
  
  \[
  \forall c > 0, n_0 > 0 \quad \exists n \geq n_0 : f(n) < 0 \text{ or } f(n) > cn
  \]

- Consider any arbitrary \( c > 0, n_0 > 0 \)

- We find some \( n \geq n_0 \) such that \( n^2 + n < 0 \) or \( n^2 + n > cn \)
  
  - \( n^2 + n > cn \) iff \( n^2 + n - cn > 0 \) iff \( n(n + 1 - c) > 0 \)
  
  - For \( n \geq n_0 > 0 \), this holds iff \( n + 1 - c > 0 \), equivalently \( n > c - 1 \)
  
  - So, \( n = \max\{c, n_0\} \) will suffice
Comparing growth rates

you vs. the guy she tells you not to worry about

\[ O(n^2) \quad | \quad O(n \log n) \]
Some Common Growth Rates (in increasing order)

polynomial

- $\Theta(1)$
- $\Theta(\log n)$
- $\Theta(\sqrt{n})$
- $\Theta(n)$
- $\Theta(n^2)$
- $\Theta(n^c)$

exponential

- $\Theta(1.1^n)$
- $\Theta(2^n)$
- $\Theta(e^n)$
- $\Theta(n!)$
- $\Theta(n^n)$
Suppose that $f(n) > 0$ and $g(n) > 0$ for all $n \geq n_0$. Suppose that

$$L = \lim_{n \to \infty} \frac{f(n)}{g(n)}.$$

Then

$$f(n) \in \begin{cases} 
    o(g(n)) & \text{if } L = 0 \\
    \Theta(g(n)) & \text{if } 0 < L < \infty \\
    \omega(g(n)) & \text{if } L = \infty.
\end{cases}$$
Constant Function Rule

The limit of a constant function is the constant:

\[ \lim_{x \to a} C = C. \]

Sum Rule

This rule states that the limit of the sum of two functions is equal to the sum of their limits:

\[ \lim_{x \to a} [f(x) + g(x)] = \lim_{x \to a} f(x) + \lim_{x \to a} g(x). \]

All of the identities shown hold **only if the limits exist**.
**Product Rule**

This rule says that the limit of the product of two functions is the product of their limits (if they exist):

\[ \lim_{x \to a} [f(x)g(x)] = \lim_{x \to a} f(x) \cdot \lim_{x \to a} g(x) . \]

**Quotient Rule**

The limit of quotient of two functions is the quotient of their limits, provided that the limit in the denominator function is not zero:

\[ \lim_{x \to a} \frac{f(x)}{g(x)} = \frac{\lim_{x \to a} f(x)}{\lim_{x \to a} g(x)} , \text{ if } \lim_{x \to a} g(x) \neq 0. \]
**Limit Rules 3/3**

**Power Rule**

\[
\lim_{x \to a} [f(x)]^p = \left[ \lim_{x \to a} f(x) \right]^p,
\]

**Limit of an Exponential Function**

\[
\lim_{x \to a} b^{f(x)} = b^{\lim_{x \to a} f(x)}
\]

**Limit of a Logarithm of a Function**

\[
\lim_{x \to a} \log_b f(x) = \log_b \lim_{x \to a} f(x)
\]

(Where base \( b > 0 \))
L’HOSPITAL’S RULE

• Often we take the limit of \( \frac{f(n)}{g(n)} \) where both \( f(n) \) and \( g(n) \) tend to \( \infty \), or both \( f(n) \) and \( g(n) \) tend to 0

• Such limits require L’Hospital’s rule
  • This rule says the limit of \( \frac{f(n)}{g(n)} \) in this case is the same as the limit of the derivative
  
• In other words, \( \lim_{n \to \infty} \frac{f(n)}{g(n)} = \lim_{n \to \infty} \frac{\frac{dn}{dn} f(n)}{\frac{dn}{dn} g(n)} \)
Using The Limit Method: Exercise 1

• Compare growth rate of $n^2$ and $n^2 - 7n - 30$

• $\lim_{n \to \infty} \frac{n^2 - 7n - 30}{n^2}$

• $= \lim_{n \to \infty} (1 - \frac{7}{n} - \frac{30}{n^2})$

• $= 1$

• So $n^2 - 7n - 30 \in \Theta(n^2)$
• Compare growth rate of \((\ln n)^2\) and \(n^{1/2}\)

\[
\lim_{n \to \infty} \frac{(\ln n)^2}{n^{1/2}} = \lim_{n \to \infty} \frac{d}{dn}(\ln n)^2
\]

*YO DAWG I HEARD YOU LIKE NATURAL LOGS*

*SO I PUT A NATURAL LOG INSIDE YOUR NATURAL LOG SO YOU CAN DERIVE WHILE YOU DERIVE*
USING THE LIMIT METHOD: EXERCISE 2

• Compare growth rate of \((\ln n)^2\) and \(n^{1/2}\)

\[
\lim_{n \to \infty} \frac{d}{dn} \left( \frac{d}{dn} \frac{(\ln n)^2}{n^{1/2}} \right)
\]

\[
= \lim_{n \to \infty} \frac{2 \ln n (1/n)}{\frac{1}{2} n^{-1/2}}
\]

\[
= \lim_{n \to \infty} \frac{4 \ln n}{n^{1/2}}
\]

\[
= \lim_{n \to \infty} \frac{4}{n^{1/2}}
\]

\[
= \lim_{n \to \infty} \frac{8}{n^{1/2}}
\]

\[
= 0
\]

\(\text{So, } (\ln n)^2 \in o(n^{1/2})\)
Additional Exercises

1. Compare the growth rate of the functions \((3 + (-1)^n)n\) and \(n\).

2. Compare the growth rates of the functions \(f(n) = n|\sin \pi n/2| + 1\) and \(g(n) = \sqrt{n}\).
SUMMATIONS AND SEQUENCES
Algebra of Order Notations

“Maximum” rules: Suppose that $f(n) > 0$ and $g(n) > 0$ for all $n \geq n_0$. Then:

\[ O(f(n) + g(n)) = O(\max\{f(n), g(n)\}) \]
\[ \Theta(f(n) + g(n)) = \Theta(\max\{f(n), g(n)\}) \]
\[ \Omega(f(n) + g(n)) = \Omega(\max\{f(n), g(n)\}) \]

“Summation” rules: Suppose $I$ is a finite set. Then

\[ O\left(\sum_{i \in I} f(i)\right) = \sum_{i \in I} O(f(i)) \]
\[ \Theta\left(\sum_{i \in I} f(i)\right) = \sum_{i \in I} \Theta(f(i)) \]
\[ \Omega\left(\sum_{i \in I} f(i)\right) = \sum_{i \in I} \Omega(f(i)) \]
Summation rules are commonly used in loop analysis.

Example:

\[
\sum_{i=1}^{n} O(i) = O \left( \sum_{i=1}^{n} i \right) = O \left( \frac{n(n+1)}{2} \right) = O(n^2).
\]
**Arithmetic sequence:**

\[
\sum_{i=0}^{n-1} (a + di) = na + \frac{dn(n-1)}{2} \in \Theta(n^2).
\]

**Geometric sequence:**

\[
\sum_{i=0}^{n-1} ar^i = \begin{cases} 
    ar^{n-1} \in \Theta(r^n) & \text{if } r > 1 \\
    na \in \Theta(n) & \text{if } r = 1 \\
    a \frac{1-r^n}{1-r} \in \Theta(1) & \text{if } 0 < r < 1.
\end{cases}
\]
**Arithmetic-geometric sequence:**

\[
\sum_{i=0}^{n-1} (a + di)r^i = \frac{a}{1 - r} - \frac{(a + (n - 1)d)r^n}{1 - r} + \frac{dr(1 - r^{n-1})}{(1 - r)^2}
\]

provided that \( r \neq 1 \).

**Harmonic sequence:**

\[
H_n = \sum_{i=1}^{n} \frac{1}{i} \in \Theta(\log n)
\]
Miscellaneous Formulae

\[ n! \in \Theta \left( n^{n+1/2}e^{-n} \right) \]

\[ \log n! \in \Theta(n \log n) \]

Another useful formula is

\[ \sum_{i=1}^{\infty} \frac{1}{i^2} = \frac{\pi^2}{6}, \]

which implies that

\[ \sum_{i=1}^{n} \frac{1}{i^2} \in \Theta(1). \]

A sum of powers of integers when \( c \geq 1 \):

\[ \sum_{i=1}^{n} i^c \in \Theta(n^{c+1}). \]
LOGARITHM RULES
Logarithm Formulae

1. \( \log_b xy = \log_b x + \log_b y \)
2. \( \log_b x/y = \log_b x - \log_b y \)
3. \( \log_b 1/x = -\log_b x \)
4. \( \log_b x^y = y \log_b x \)
5. \( \log_b a = \frac{1}{\log_a b} \)
6. \( \log_b a = \frac{\log_c a}{\log_c b} \)
7. \( a^{\log_b c} = c^{\log_b a} \)
We typically omit the base, and just write $\Theta(\log x)$ for this reason.
Running Time of a Program: $T_M(I)$ denotes the running time (in seconds) of a program $M$ on a problem instance $I$.

Worst-case Running Time as a Function of Input Size: $T_M(n)$ denotes the \textit{maximum} running time of program $M$ on instances of size $n$:

$$T_M(n) = \max\{ T_M(I) : \text{Size}(I) = n \}.$$ 

Average-case Running Time as a Function of Input Size: $T_M^{\text{avg}}(n)$ denotes the \textit{average} running time of program $M$ over all instances of size $n$:

$$T_M^{\text{avg}}(n) = \frac{1}{|\{I : \text{Size}(I) = n \}|} \sum_{\{I : \text{Size}(I) = n \}} T_M(I).$$ 

But how do we know how much time $M$ will take on input $I$?

We don’t know how much time an \textit{individual step} in the program takes!
MODEL OF COMPUTATION

• Before we can analyze the running time of code, we need a precise **model of computation**

• We use the Word-RAM model
  
  • Each memory location is a word that can hold an integer
  
  • Accessing a word of memory takes constant time
  
  • Basic operations (arithmetic, shifting, logical operators) take constant time

• Is a word large enough to hold **any** integer?

  But, large enough to hold an **address** of an object in a data structure? **Yes!** If the data structure fits in RAM…

  Unlikely… split large integers into **arrays** of words. Watch cost of operations!
LOOP ANALYSIS
META-ALGORITHM FOR ANALYZING LOOPS

• Identify operations that require only constant time
• The complexity of a loop is the sum of the complexities of all iterations
• Analyze independent loops separately and add the results
• If loops are nested, it often helps to start at the innermost, and proceed outward… but,
  • sometimes you must express several nested loops together in a single equation (using nested summations),
  • and actually evaluate the nested summations… (can be hard)
TWO BIG-O ANALYSIS STRATEGIES

• Strategy 1
  • Prove a $O$-bound and a matching $\Omega$-bound separately to get a $\Theta$-bound.

• Strategy 2
  • Use $\Theta$-bounds throughout the analysis and thereby obtain a $\Theta$-bound for the complexity of the algorithm.

Often easier (but not always)
Algorithm: *LoopAnalysis1*(n : integer)

1. \( \text{sum} \leftarrow 0 \)
2. for \( i \leftarrow 1 \) to \( n \)
   
   \[
   \begin{align*}
   & \text{for } j \leftarrow 1 \text{ to } i \\
   & \quad \text{do } \\
   & \quad \quad \text{do } \\
   & \quad \quad \quad \text{sum} \leftarrow \text{sum} + (i - j)^2 \\
   & \quad \quad \text{sum} \leftarrow \lfloor \text{sum} / i \rfloor \\
   \end{align*}
   \]
3. return (\( \text{sum} \))
**Strategy 1:** big-O and big-Ω bounds

We focus on the two nested for loops (i.e., (2)). The total number of iterations is \( \sum_{i=1}^{n} i \), with \( \Theta(1) \) time per iteration.

**Upper bound:**
\[
\sum_{i=1}^{n} O(i) \leq \sum_{i=1}^{n} O(n) = O(n^2).
\]

**Lower bound:**
\[
\sum_{i=1}^{n} \Omega(i) \geq \sum_{i=n/2}^{n} \Omega(i) \geq \sum_{i=n/2}^{n} \Omega(n/2) = \Omega(n^2/4) = \Omega(n^2).
\]

Since the upper and lower bounds match, the complexity is \( \Theta(n^2) \).
Strategy 2: use $\Theta$-bounds throughout the analysis

Algorithm: $\text{LoopAnalysis1}(n : \text{integer})$

1. $sum \leftarrow 0$
2. \textbf{for} $i \leftarrow 1$ \textbf{to} $n$
   \hspace{1em} \textbf{for} $j \leftarrow 1$ \textbf{to} $i$
   \hspace{2em} \textbf{do}
   \hspace{3em} \textbf{do}
   \hspace{4em} $sum \leftarrow sum + (i - j)^2$
   \hspace{3em} $sum \leftarrow \lfloor sum / i \rfloor$
3. \textbf{return} $(sum)$

$\Theta$-bound analysis

\begin{align*}
\sum_{i=1}^{n} \Theta(i) &= \Theta \left( \sum_{i=1}^{n} i \right) = \Theta \left( \frac{n(n+1)}{2} \right) = \Theta(n^2).
\end{align*}

\begin{align*}
(1) & \quad \Theta(1) \\
(2) & \quad \text{Complexity of inner } \textbf{for} \text{ loop: } \Theta(i) \\
& \quad \text{Complexity of outer } \textbf{for} \text{ loop: } \sum_{i=1}^{n} \Theta(i) = \Theta(n^2) \\
(3) & \quad \Theta(1) \\
\text{total} & \quad \Theta(1) + \Theta(n^2) + \Theta(1) = \Theta(n^2)
\end{align*}
Consider this loop alone...

number of loop iterations?

\( j \) starts at \( i \) and is repeatedly divided by 2... this can happen only \( \Theta(\log i) \) times

So inner loop has runtime \( \Theta(\log i) \)

And the entire inner loop is executed for \( i = 1, 2, ..., n \)

So, we have \( T(n) \in \Theta(\sum_{i=1}^{n} \log i) \)

\[
T(n) \in O \left( \sum_{i=1}^{n} \log i \right) \subseteq O(n \log n)
\]

\[
T(n) \in \Omega \left( \sum_{i=1}^{n} \log i \right) \subseteq \Omega(n \log n)
\]

\[
T(n) \in \Omega \left( \sum_{i=1}^{n} \log \frac{n}{i} \right) \subseteq \Omega(n \log n)
\]
... ANOTHER EXERCISE IN LOOP ANALYSIS?

Olive Garden waiter: Sir, you've already had 5 baskets of breadsticks

Me:

We're done when I say we're done
max := 0;
for i := 1 to n do
  for j := i to n do
    sum := 0;
    for k := i to j do
      sum := sum + A[k];
    if sum > max then max := sum;
Strategy 1: big-O and big-Ω bounds

\[ T(n) \in \Theta(1) + \sum_{i=1}^{n} \sum_{j=i}^{n} \left( \Theta(1) + \sum_{k=i}^{j} \Theta(1) + \Theta(1) \right) \]

\[ = \sum_{i=1}^{n} \sum_{j=i}^{n} \Theta(j - i) \in \Theta \left( \sum_{i=1}^{n} \sum_{j=i}^{n} (j - i) \right) \]

\[ T(n) \in O \left( \sum_{i=1}^{n} \sum_{j=i}^{n} (j - i) \right) \leq O \left( \sum_{i=1}^{n} \sum_{j=i}^{n} n \right) \]

\[ \leq O \left( \sum_{i=1}^{n} \sum_{j=1}^{n} n \right) \]

\[ T(n) \in O(n^3) \]

This is the maximum number of iterations that could be performed in this loop.
Proving a big-Ω bound…

Recall:

\[ T(n) \in \Theta \left( \sum_{i=1}^{n} \sum_{j=i}^{n} (j - i) \right) \]

\[ T(n) \in \Omega \left( \sum_{i=1}^{n} \sum_{j=i}^{n} (j - i) \right) \]

\[ \geq \Omega \left( \sum_{i=1}^{n/2} \sum_{j=i}^{n} (j - i) \right) \]

\[ \geq \Omega \left( \sum_{i=1}^{n/2} \sum_{j=3n/4}^{n} (j - i) \right) \]

Intuition: \( j - i \) is \( \Omega(n) \) in some iterations. How many iterations? Lots?

To get a good \( \Omega \)-bound, we ask questions like:

When do our loops have many iterations?

When is our dominant term large?

Many iterations: when our \( j \) loop does \( \Omega(n) \) iterations! For example, when \( i \leq n/2 \)…

Large dominant term: when \( j \) is much larger than \( i \) (i.e., by a factor of \( n \))
Proving a big-Ω bound... continued

Recall:

\[ T(n) \in \Omega \left( \sum_{i=1}^{n/2} \sum_{j=3n/4}^{n} (j - i) \right) \]

\[ \geq \Omega \left( \sum_{i=1}^{n/2} \sum_{j=3n/4}^{n} \left( \frac{3n}{4} - \frac{n}{2} \right) \right) \]

\[ = \Omega \left( \sum_{i=1}^{n/2} \sum_{j=3n/4}^{n} \frac{n}{4} \right) \]

\[ \geq \Omega \left( \frac{n \cdot n \cdot n}{2 \cdot 4 \cdot 4} \right) = \Omega(n^3) \]

Smallest possible value of \( j - i \) for these bounds on \( i,j \)

We will perform at least this much work in every iteration!

This term does not depend on the loop indexes, so just multiply by the total number of loop iterations...

Since we have \( O(n^3) \) and \( \Omega(n^3) \), we have proved \( \Theta(n^3) \)