CS 343 Fall 2024 – Assignment 4
Instructor: Peter Buhr
Due Date: Wednesday, November 6, 2024 at 22:00
Late Date: Friday, November 8, 2024 at 22:00

October 25, 2024

This assignment introduces complex locks in μC++ and continues examining synchronization and mutual exclusion. Use it to become familiar with these new facilities, and ensure you use these concepts in your assignment solution. **(Tasks may *not* have public members except for constructors and/or destructors.)**

1. Figure 1 is a C++ program comparing buffering using internal-data versus external-data format.

   (a) Compare the three versions of the program with respect to performance by doing the following:

   - Compile the program with u++ and run the program with preprocessor variables −DARRAY, −DSTRING and −DSTRSTREAM.
   - Time the executions using the time command:

     ```
     $ /usr/bin/time -f "%Uu %Ss %Er %Mkb" ./a.out 100000000 20
     3.21u 0.02s 0:03.32
     ```

     (Output from time differs depending on the shell, so use the system time command.) Compare the *user* time (3.21u) only, which is the CPU time consumed solely by the execution of user code (versus system and real time).
   - Start with the command-line argument 100000000 20 and adjust the times amount (if necessary) to get program execution into the range 1 to 100 seconds for the 3 versions of the program. (Timing results below 1 second are inaccurate.) Otherwise, increase/decrease the times amount as necessary and scale the difference in the answer.
   - Run the experiments again after recompiling the programs with compiler optimization turned on (i.e., compiler flag −O2).
   - Include all 6 timing results to validate the experiments and the number of calls to malloc.

   (b) Explain the performance and allocation difference (larger/smaller/by how much) between the three versions of the program.

   (c) Explain the performance difference (larger/smaller/by how much) when compiler optimization is used.

   (d) Very briefly (1-2 sentences) speculate on the cause of the performance difference between ARRAY and STRING / STRSTREAM.

2. Figure 2, p. 3 shows the Tsay[1] modified Peterson solution for two-thread mutual-exclusion.

   (a) Does the algorithm work if lines 5 and 6 are interchanged? If the changed algorithm does not work, explain a pathological situation that results in failure of one rule of the critical-section game. Be sure to clearly specify which rule is broken and why.

   (b) Run the changed algorithm multiple times with compilation command:

   ```
   $ u++ -O3 -multi -nodebug Tsay.cc
   ```

   Based on your answer to question 2a, explain why the algorithm runs without failure.

[1] Yih-Kuen Tsay, "Deriving a scalable algorithm for mutual exclusion", Distributed Computing, 1998, Springer, pp. 393–407,

```
#include <iostream>
#include <string>
#include <sstream>
using namespace std;

int main( int argc, char * argv[] ) {
    intmax_t times = 1'000'000, size = 20;                  // defaults
    const char * nosummary = getenv( "NOSUMMARY" );         // print summary ?
    struct cmd_error {};                                    // command-line errors

    try {
        switch ( argc ) {
          case 3: size = convert( argv[2] ); if ( size <= 0 ) { throw cmd_error(); }
          case 2: times = convert( argv[1] ); if ( times <= 0 ) { throw cmd_error(); }
          case 1: break;                                    // use defaults
          default: throw cmd_error();
        } // switch
    } catch( ... ) {
        cerr << "Usage: " << argv[0] << " [ times (> 0) [ size (> 0) ] ]" << endl;
        exit( EXIT_FAILURE );
    } // try

    #if defined( ARRAY )
    struct S { long long int i, j, k, l; };
    S buf[size];                                            // internal-data buffer
    #elif defined( STRING )
    string strbuf;                                          // external-data buffer
    #elif defined( STRSTREAM )
    stringstream ssbuf;                                     // external-data buffer
    #else
        #error unknown buffering style
    #endif

    enum { C = LLONG_MAX };                                 // print multiple characters
    int trick = rand();                                     // prevent dead-code removal

    for ( int t = 0; t < times; t += 1 ) {
        #if defined( ARRAY )
        for ( int s = 0; s < size; s += 1 ) {
            buf[s] = (S){ C, C, C, C };
            if ( trick == 0 ) cout << buf[s].i;
        } // for
        #elif defined( STRING )
        for ( int s = 0; s < size; s += 1 ) {
            strbuf += to_string(C) + '\t' + to_string(C) + '\t' + to_string(C) + '\t' + to_string(C) + '\t';
            if ( trick == 0 ) cout << strbuf;
        } // for
        strbuf.clear();                                     // reset string
        #elif defined( STRSTREAM )
        for ( int s = 0; s < size; s += 1 ) {
            ssbuf << C << '\t' << C << '\t' << C << '\t' << C << '\t';
            if ( trick == 0 ) cout << ssbuf.str();
        } // for
        ssbuf.seekp( 0 );                                   // reset stream
        #else
            #error unknown buffering style
        #endif
    } // for

    if ( ! nosummary ) { malloc_stats(); }                  // print heap statistics
}
```

Figure 1: Internal versus External Buffering

```
static void CriticalSection() {
    static uBaseTask * volatile curr;                    // shared
    curr = &uThisTask();
    for ( unsigned int i = 0; i < 100; i += 1 ) {        // work
        if ( curr != &uThisTask() ) { abort( "Interference" ); } // check
    } // for
}

enum Intent { WantIn, DontWantIn };

_Task Tsay {
    static Intent * volatile Last;
    Intent & me, & you;

    void main() {
        for ( unsigned int i = 0; i < 10'000'000; i += 1 ) {
1           me = WantIn;                                 // entry protocol
2           Last = &me;                                  // race
            uFence();                                    // prevent hardware reordering (x86)
3           if ( you == WantIn )
4               while ( Last == &me ) {}
            CriticalSection();                           // critical section
5           me = DontWantIn;                             // exit protocol
6           Last = &me;
        }
    }
  public:
    Tsay( Intent & me, Intent & you ) : me( me ), you( you ) {}
};
Intent * volatile Tsay::Last;

int main() {
    uProcessor p;
    Intent me = DontWantIn, you = DontWantIn;            // shared
    Tsay t0( me, you ), t1( you, me );
}
```

Figure 2: Tsay 2-Thread Mutual Exclusion

3. (a) Consider the following situation involving a tour group of V tourists. The tourists arrive at the Louvre museum for a tour. However, a tour group can only be composed of G people at a time, otherwise the tourists cannot hear the guide. As well, there are 3 kinds of tours available at the Louvre: pictures, statues and gift shop. Therefore, each tour group must vote to select the kind of tour to take. Voting is a *ranked ballot*, where each tourist ranks the 3 tours with values 0, 1, 2, where 2 is the highest rank. Tallying the votes sums the ranks for each kind of tour and selects the highest ranking. If tie votes occur among rankings, prioritize the results by gift shop, pictures, and then statues, e.g.:

	P	S	G			P	S	G	
tourist1	0	1	2		tourist1	2	1	0	
tourist2	2	1	0		tourist2	1	2	0	
tally	2	2	2	all ties, select G	tally	3	3	0	two ties, select P

During voting, a tourist blocks until all G votes are cast, i.e., assume a secret ballot. Once a decision is made, the tourists in that group proceed on the specified tour. Tourists may take multiple tours, but because of voting, can take the same kind of tour.

The tour size G may not evenly divide the number of tourists, resulting in a *quorum* failure when the remaining tourists is less than G.

Implement a general vote-tallier as a class using *only*:

i. a single uOwnerLock and two uCondLocks to provide mutual exclusion and synchronization plus a signalling flag, and implement using ***barging avoidance***. Hint: uCondLock::signal returns true if a

```
#if defined( MC )                    // mutex/condition solution
#include "BargingCheckVote.h"
// includes for this kind of vote-tallier
class TallyVotes {
    // private declarations for this kind of vote-tallier
#elif defined( SEM )                 // semaphore solution
#include "BargingCheckVote.h"
// includes for this kind of vote-tallier
class TallyVotes {
    // private declarations for this kind of vote-tallier
#elif defined( BAR )                 // barrier solution
// includes for this kind of vote-tallier
_Cormonitor TallyVotes : public uBarrier {
    // private declarations for this kind of vote-tallier
#else
    #error unsupported voter type
#endif
    // common declarations
  public:                            // common interface
    _Exception Failed {};
    struct Ballot { unsigned int picture, statue, giftshop; };
    enum TourKind : char { Picture = 'p', Statue = 's', GiftShop = 'g' };
    struct Tour { TourKind tourkind; unsigned int groupno; };

    TallyVotes( unsigned int voters, unsigned int group, Printer & printer );
    Tour vote( unsigned int id, Ballot ballot );
    void done(
        #if defined( MC ) || defined( BAR )
        unsigned int id
        #endif
    );
};
```

Figure 3: Tally Vote Interfaces

task is unblocked and false otherwise. **Warning: solutions with only one uCondLock are highly unlikely to work.**

ii. uSemaphores, used as binary not counting, to provide mutual exclusion and synchronization, and implement using ***barging prevention***.

iii. a single uBarrier to provide mutual exclusion and synchronization. Note, a uBarrier has implicit mutual exclusion so it is only necessary to manage the synchronization. As well, only the basic aspects of the uBarrier are needed to solve this problem.

No unbounded busy-waiting is allowed in any solution, and barging tasks can spoil an election and must be avoided/prevented.

Figure 3 shows the different forms for each μC++ vote-tallier implementation (you may add only a public destructor and private members), where the preprocessor is used to conditionally compile a specific interface **(see the Makefile)**. This form of header file removes duplicate code.

At creation, a vote-tallier is passed the number of voters, size of a voting group, and a printer for printing state transitions. There is only one vote-tallying object created for all of the voters, who share a reference to it. Each voter task calls the vote method with their id and a ranked vote, indicating their desire for a picture, statue, or gift-shop tour. The vote routine does not return until group votes are cast; after which, the majority result of the voting (Picture, Statue or GiftShop) is returned to each voter, along with a number to identify the tour group (where tours are numbered 1 to N). The groups are formed based on voter arrival; e.g., for a group of 3, if voters 2, 5, 8 cast their votes first, they form the first group, etc. Hence, all voting is serialized. When a tourist finishes taking tours and leaves the Louvre Museum, it *always* calls done (even if it has a quorum failure).

TallyVotes detects a quorum failure when the number of **remaining voters in the Louvre** is less than the group size. At this point, any new calls to vote immediately raise exception Failed, and any waiting voters

```
#include "BargingCheckVote.h"
class TallyVotes {
    ...                                     // regular declarations
    BCHECK_DECL;
  public:
    ...                                     // regular declarations
    Tour vote( unsigned int id __attribute__(( unused )), Ballot ballot ) {
        // acquire mutual exclusion
        VOTER_ENTER( tour-group-size );
        ...                                 // voter code
        VOTER_LEAVE( tour-group-size );
        // release mutual exclusion
        return ...
    }
};
```

Figure 4: Barging Check Macros: MC and SEM

must be unblocked so they can raise exception Failed. When a voter calls done, it must cooperate if there is a quorum failure by helping to unblock waiting voters. For example, with the mutex/condition and barrier lock, a voter calling done in the failure case may have to block (pretend to be a barger) to force waiting voters to unblock.

Note, even when V *is* a multiple of G *and* tourists take multiple tours, a quorum failure can occur. For example, one tour is faster than another or a tourist leaves a tour early and comes back to vote on another tour, so the quick tourist finishes all their tours and terminates. The slower tourists then encounter a situation where there are insufficient tourists to form a quorum for later tours.

Figure 4 shows the macro placement that ***must*** be present only in the MC and SEM tally-votes implementation to test for barging, and defining preprocessor variable BARGINGCHECK triggers barging testing **(see the Makefile)**. If barging is detected, a message is printed and the program continues, possibly printing more barging messages. To inspect the program with gdb when barging is detected, set BARGINGCHECK=0 to abort the program.

Figure 5 shows the interface for a voting task (you may add only a public destructor and private members). The task main of a voting task first

- yields a random number of times, between 0 and 19 inclusive, so all tasks do not start simultaneously

and then performs the following tours times:

- print start message
- yield a random number of times, between 0 and 4 inclusive
- vote
- yield a random number of times, between 0 and 4 inclusive
- print going on tour message
- after all tours, eventually report done and print terminate message

Casting a vote is accomplished by calling member cast. Yielding is accomplished by calling yield(times) to give up a task's CPU time-slice a number of times.

All output from the program is generated by calls to a printer, *excluding error messages*. Figure 6 shows the interface for the printer (you may add only a public destructor and private members). (For now, treat **_Monitor** as a class and **_Cormonitor** as a coroutine with public methods that implicitly provide mutual exclusion.) The printer attempts to reduce output by storing information for each voter until one of the stored elements is overwritten. When information is going to be overwritten, all the stored information is flushed and storing starts again. Output must look like that in Figure 8, p. 7.

Each column is assigned to a voter with the titles, "V_i", and Figure 7 shows the column entries indicating its current status. Information is buffered until a column is overwritten for a particular entry, which causes the buffered data to be flushed. If there is no new stored information for a column since the last buffer flush, an empty column is printed. After a task has terminated, no further output appears in that column. All output spacing can be accomplished using the standard 8-space tabbing (`'\t'`). Buffer any information necessary for printing in its internal representation; **do not build and store strings of text for output**.

```
_Task Voter {
    TallyVotes::Ballot cast() __attribute__(( warn_unused_result )) {  // cast 3-way vote
        // O(1) random selection of 3 items without replacement using divide and conquer.
        static const unsigned int voting[3][2][2] = { { {2,1}, {1,2} }, { {0,2}, {2,0} }, { {0,1}, {1,0} } };
        unsigned int picture = prng( 3 ), statue = prng( 2 );
        return (TallyVotes::Ballot){ picture, voting[picture][statue][0], voting[picture][statue][1] };
    }
  public:
    enum States : char { Start = 'S', Vote = 'V', Block = 'B', Unblock = 'U', Barging = 'b',
          Done = 'D', Complete = 'C', Going = 'G', Failed = 'X', Terminated = 'T' };
    Voter( unsigned int id, unsigned int tours, TallyVotes & voteTallier, Printer & printer );
};
```

Figure 5: Voter Interface

```
_Monitor / _Cormonitor Printer {        // chose one of the two kinds of type
  public:
    Printer( unsigned int voters );
    void print( unsigned int id, Voter::States state );
    void print( unsigned int id, Voter::States state, TallyVotes::Tour tour );
    void print( unsigned int id, Voter::States state, TallyVotes::Ballot vote );
    void print( unsigned int id, Voter::States state, unsigned int numBlocked );
    void print( unsigned int id, Voter::States state, unsigned int numBlocked, unsigned int group );
};
```

Figure 6: Printer Interface

Calls to perform printing may be performed from the vote-tallier and/or a voter task (you decide where to print).

For example, in the left-hand example of Figure 8, there are 3 voters, 1 voter in a group, and each voter attempts to vote once. At line 4, V0 has the value "S" in its buffer slot, V1 has value "S", and V2 is empty. When V1 attempts to print "V 0,2,1", which overwrites its current buffer value of "S", the buffer must be flushed generating line 4. V1's new value of "V 0,2,1" is then inserted into its buffer slot. When V1 attempts to print "C", which overwrites its current buffer value of "V 0,2,1", the buffer must be flushed generating line 5, and no other values are printed on the line because the print is consecutive (i.e., no intervening call from another object). Then V1 inserts value "C" and V0 inserts value "V 2,0,1" into the buffer. Assume V0 attempts to print "C", which overwrites its current buffer value of "V 2,0,1", the buffer must be flushed generating line 6, and so on. Note, a group size of 1 means a voter never has to block/unblock.

For example, in the right-hand example of Figure 8, there are 6 voters, 3 voters in a group, and each voter attempts to vote twice. Voters V3 and V4 are delayed (e.g., they went to Tom's for a coffee and donut). By looking at the F codes, V0, V1, V5 vote together (group 1), V0, V1 V2 vote together (group 2), and V2, V4, V5 vote together (group 3). Hence, V0, V1, V2, and V5 have voted twice and terminated. V3 needs to vote

State	Meaning
S	start
V p, s, g	vote with ballot containing 3 rankings
B n	block during voting, n voters waiting (including self)
U n	unblock after group reached, n voters still waiting (not including self)
b n gn	block barging task (avoidance only), n waiting for signalled tasks to unblock (including self), group number gn of last group that received a voting result
D	block in done (MC/BAR only)
C t	complete group and voting result is t (p/s/g)
G t gn	go on tour, t (p/s/g) in tour group number gn
X	failed to form a group (quorum failure)
T p, s, g	voter terminates (after call to done), total number of tours attended of each kind

Figure 7: Voter Status Entries

```
$ vote 3 1 1
V0      V1      V2
*******
S       S
        V 0,2,1
V 2,0,1 C s
C p     G s 1
G p 2   T 0,1,0 S
T 1,0,0         V 2,0,1
                C p
                G p 3
                T 1,0,0
*****************
All tours ended
```

```
$ vote 6 3 2
V0      V1      V2      V3      V4      V5
*******
S       S                               S
        V 1,0,2
        B 1                             V 2,0,1
V 2,0,1                                 B 2
C p     U 1
G p 1
S               S                       U 0
                V 2,0,1
                B 1                     G p 1
V 2,1,0                                 S
B 2     G p 1
        S
        V 1,2,0
        C p     U 1
U 0     G p 2                   S       b 1 2
        T 2,0,0 G p 2
G p 2           S       S       b 2 2   V 1,0,2
T 2,0,0         b 2 2           V 1,2,0 B 1
                V 1,0,2 b 2 2   B 2
                C g                     U 1
                G g 3           U 0     G g 3
                        V 0,2,1         T 1,0,1
                        B 1     G g 3
                T 1,0,1 U 0     S
                        X
                        T 0,0,0 X
                                T 0,0,1
*****************
All tours ended
```

Figure 8: Voters: Example Output

twice and V4 needs to vote again. However, there are now insufficient voters to form a group, so both V3 and V4 fail with X.

The executable program is named vote and has the following shell interface:

vote [voters | 'd' [group | 'd' [tours | 'd' [seed | 'd' [processors | 'd']]]]]

voters is the size of a tour (> 0), i.e., the number of voters (tasks) to be started. If d or no value for voters is specified, assume 6.

group is the size of a tour group (> 0). If d or no value for group is specified, assume 3.

tours is the number of tours (> 0) each voter takes of the museum. If d or no value for tours is specified, assume 1.

seed is the starting seed for the random-number generator (> 0). If seed is specified, call set_seed(seed). If d or no value for seed is specified, do nothing as PRNG sets the seed to an arbitrary value.

processors is the number of processors for parallelism (> 0). If d or no value for processors is specified, assume 1. Use this number in the following declaration placed in the program main immediately after checking command-line arguments but before creating any tasks:

```
uProcessor p[processors - 1] __attribute__(( unused )); // create more kernel thread
```

to adjust the amount of parallelism for computation. The program starts with one kernel thread so only processors − 1 additional kernel threads are added.

To obtain semi-repeatable results, all random numbers are generated using the μC++ task-member prng (see Appendix C in the μC++ reference manual). Note, because of the non-deterministic execution of concurrent programs, multiple runs with a common seed may not generate the same output. Nevertheless, shorts runs are often the same so the seed can be useful for testing.

Check all command arguments for correct form (integers) and range; print an appropriate usage message and terminate the program if a value is missing or invalid.

(b) Recompile the program to elide output for timing experiments by adding the following code and using it to bracketing all printer calls **(see the Makefile)**.

```
#ifdef NOOUTPUT
#define PRINT( stmt )
#else
#define PRINT( stmt ) stmt
#endif // NOOUTPUT
PRINT( printer.print( id, Voter::Vote, ballot ) ); // elide printer call
```

i. Compare the performance among the 3 kinds of locks:

- Time the executions using the time command with output and barging checking turned off:

  ```
  $ /usr/bin/time -f "%Uu %Ss %Er %Mkb" vote 100 10 10000 1003
  3.21u 0.02s 0:05.67r 32496kb
  ```

  Output from time differs depending on the shell, so use the system time command. Compare the *user* (3.21u) and *real* (0:05.67r) time among runs, which is the CPU time consumed solely by the execution of user code (versus system) and the total time from the start to the end of the program.
- If necessary, adjust the number of voters and then tours to get real time in range 1 to 100 seconds. (Timing results below 1 second are inaccurate.) Use the same number of tours for all experiments.
- Include all 3 timing results to validate your experiments.
- Repeat the experiment using 2 processors and include the 3 timing results to validate your experiments.

ii. Explain the performance difference (larger/smaller/by how much) among the locks.

iii. As the kernel threads increase, very briefly speculate on any performance difference.

Submission Guidelines

Follow these guidelines carefully. Review the Assignment Guidelines and C++ Coding Guidelines *before* starting each assignment. **Each text or test-document file, e.g., *.{txt,testdoc} file, must be ASCII text and not exceed 500 lines in length, using the command fold -w120 *.testdoc | wc -l.** Programs should be divided into separate compilation units, i.e., *.{h,cc,C,cpp} files, where applicable. Use the submit command to electronically copy the following files to the course account.

1. q1*.txt – contains the information required by question 1, p. 1.
2. q2*.txt – contains the information required by question 2, p. 1.
3. BargingCheckVote.h – barging checker (provided)
4. q3tallyVotes.h, q3*.{h,cc,C,cpp} – code for question question 3a, p. 3. **Program documentation must be present in your submitted code. No user, system or test documentation is to be submitted for this question.**
5. q3*.txt – contains the information required by question 3b.
6. Modify the following Makefile to compile the programs for question 3a, p. 3 by inserting the object-file names matching your source-file names.

```
VIMPL := MC
OUTPUT := OUTPUT
BCHECK := BARGINGCHECK

CXX = u++                                          # compiler
CXXFLAGS = -g -multi -Wall -Wextra -MMD -D"${VIMPL}" -D"${OUTPUT}" \
           -D"${BCHECK}" # compiler flags
ifeq ("${OUTPUT}","NOOUTPUT")
    CXXFLAGS += -O2 -nodebug
endif
MAKEFILE_NAME = ${firstword ${MAKEFILE_LIST}} # makefile name
```

```
OBJECTS = q3tallyVotes${VIMPL}.o # list of object files for question 3 prefixed with "q3"
EXEC = vote

DEPENDS = ${OBJECTS:.o=.d}                          # substitute ".o" with ".d"

##############################################################

.PHONY : all clean
.ONESHELL :

all : ${EXEC}                                       # build all executables

-include VoteImpl

ifeq (${shell if [ "${LOCKVIMPL}" = "${VIMPL}" -a "${OUTPUTTYPE}" = "${OUTPUT}"
        -a "${BCHECKIMPL}" = "${BCHECK}" ] ; then echo true ; fi },true)
${EXEC} : ${OBJECTS}
    ${CXX} ${CXXFLAGS} $^ -o $@
else                                                # implementation type has changed => rebuilt
.PHONY : ${EXEC}
${EXEC} :
    rm -f VoteImpl
    touch q3tallyVotes.h
    ${MAKE} ${EXEC} VIMPL="${VIMPL}" OUTPUT="${OUTPUT}" BCHECK="${BCHECK}"
endif

VoteImpl :
    echo "LOCKVIMPL=${VIMPL}\nOUTPUTTYPE=${OUTPUT}\nBCHECKIMPL=${BCHECK}" > VoteImpl
    sleep 1

##############################################################

${OBJECTS} : ${MAKEFILE_NAME}                       # OPTIONAL : changes to this file => recompile

-include ${DEPENDS}                                 # include *.d files containing program dependences

clean :                                             # remove files that can be regenerated
    rm -f *.d ${OBJECTS} ${EXEC} VoteImpl
```

This makefile is invoked as follows:

```
$ make vote VIMPL=MC BCHECK=BARGINGCHECK
$ vote ...
$ make vote VIMPL=SEM OUTPUT=OUTPUT
$ vote ...
$ make vote VIMPL=BAR OUTPUT=NOOUTPUT
$ vote ...
```

Put this Makefile in the directory with the programs, name the source files as specified above, and enter the appropriate make to compile a specific version of the programs. This Makefile must be submitted with the assignment to build the program, so it must be correct. Use the web tool Request Test Compilation to ensure you have submitted the appropriate files, your makefile is correct, and your code compiles in the testing environment.

Follow these guidelines. Your grade depends on it!