



**UNIVERSITY OF  
WATERLOO**

**Final Examination**

**Term: Winter    Year: 2023**

**CS343**

**Concurrent and Parallel Programming**

**Section 001**

**Instructor: Caroline Kierstead**

**Thursday, April 13, 2023**

**Start Time: 19:30                      End Time: 22:00**

**Duration of Exam: 2.5 hours**

**Number of Exam Pages (including cover sheet): 6**

**Total number of questions: 24 multiple-choice, 2 short-answer, 2 coding**

**Total marks available: 77**

**CLOSED BOOK, NO ADDITIONAL MATERIAL ALLOWED**

**Part A – Multiple Choice**

Elided for reuse.

**Part B – Short Answer**

1. (a) **2 marks** Explain why preventing *synchronization deadlocks* is not practical.
  - (b) **2 marks** Briefly explain the difference between *deadlock prevention* and *deadlock avoidance*.
  - (c) **1 mark** Name a technique for *mutual exclusion deadlock avoidance*.
  - (d) **1 mark** What is the only reasonably practical method for preventing *mutual exclusion deadlock*?
2. Assume a programming language that provides blocking mutual exclusion locks and condition locks:

```

MutexLock m   CondLock c;
m.acquire()   c.wait();
m.release();  c.wait( m ); // atomically block and release m, no reacquire of m on unblock
              c.signal();
              c.empty();

```

Note, MutexLock maybe acquired and released by different threads.

Create a Monitor class that is capable of performing internal scheduling, similar to the one shown below, where methods foo and bar are mutually exclusive with themselves and each other (just like **\_Mutex** public methods in  $\mu$ C++ monitors):

```

Monitor M {
    CondQueue bench; // blocking bench
    int value;
    public:
    void foo();
    int bar();
};

```

- (a) **3 marks** Write the locking declaration in M and the entry and exit protocol needed to provide mutual exclusion for method foo.
- (b) **3 marks** Write the entry and exit protocol needed to provide mutual exclusion for method bar.
- (c) **3 marks** Assume a thread enters foo and performs bench.wait(m), where bench is a CondLock and m is the MutexLock providing the monitor mutual-exclusion. Assume a thread now enters bar, performs bench.signal() and immediately returns. Assume signal() returns a boolean value indicating true if a thread is signalled and false otherwise. Write *barging-prevention* code in foo and bar that guarantees the signalled thread acquires the monitor next.

**Part C – Long Answer**

1. (a) The taxi cab company, Maple Leaf Cabs, has  $N$  taxi cabs scattered throughout the city. The dispatcher's job is to take requests from clients for a taxi and to dispatch a taxi to the client for a pickup. The dispatcher also takes requests from taxis for a client at the start of the day and after each client is delivered to their destination.

The interface for the dispatcher monitor is the following (you may not change this interface; you may only add code in the designated areas L1, L2 and L3). **(Do not copy the starting code into your answer booklet.)**

```

_Monitor Dispatcher {
    // communication variables
    int xclient, yclient;           // client (x,y) coordinates
    int taxild, clientld;         // identities of taxi/client pairing

    // L1: ANY VARIABLES NEEDED FOR EACH IMPLEMENTATION

public: // common interface
    int getTaxi( int id, int x, int y ) {           // called by client
        // L2: ANY SYNCHRONIZATION NEEDED FOR EACH IMPLEMENTATION
        return taxild;
    }
    int getClient( int id, int & x, int & y ) {     // called by taxi, x,y in/out parameters
        // L3: ANY SYNCHRONIZATION NEEDED FOR EACH IMPLEMENTATION
        return clientld;
    }
};

```

A client calls the `getTaxi` routine to ask for a taxi to pick them up at an address given by parameter coordinates  $(x,y)$ . `getTaxi` returns the id of the taxi picking up the client, `taxild`.

A taxi calls the `getClient` routine to indicate it is available and tell the dispatcher the taxi's current  $(x,y)$  location. The taxi's call to `getClient` returns the client's  $(x,y)$  position in arguments  $x$  and  $y$ , and returns the id of the client being collected, `clientld`.

**Do not write or create either the taxi or client tasks. You may assume for this question that the dispatcher never shuts down.**

Implement the Dispatcher monitor using:

- i. **10 marks** external scheduling,
- ii. **14 marks** internal scheduling,
- iii. **13 marks** implicit (automatic) signalling, using *only* the 3 macros defined below.

Assume the existence of the following preprocessor macros for implicit (automatic) signalling (I(a)iii):

```

#define AUTOMATIC_SIGNAL ...
#define WAITUNTIL( predicate ) ...
#define EXIT() ...

```

Macro `AUTOMATIC_SIGNAL` is placed only once in an automatic-signal monitor as a private member, and contains any private variables needed to implement the automatic-signal monitor. Macro `WAITUNTIL` is used to wait until the predicate evaluates to true. Macro `EXIT` must be called on return from **every** public routine of an automatic-signal monitor.

The  $\mu\text{C++}$  `uCondition` operations are available at the end of the exam.

- (b) **25 marks** Write an administrator task to handle the dispatcher’s role for the Maple Leaf Taxi company. Figure 1 contains the starting code for the dispatcher-Administrator (you may add only a public destructor and private members). **(Do not copy the starting code into your exam booklet.)**

A client calls the `getTaxi` routine to ask for a taxi to pick them up at an address given by parameter coordinates  $(x,y)$ . A *future* taxi is returned immediate to the client, so the client can execute asynchronously (e.g., get ready to leave) before accessing the taxi. When the client accesses the future taxi, they may block because the taxi is not there; otherwise, the client gets into the taxi whose id is in the future.

A taxi calls the `getClient` routine to indicate it is available and tell the dispatcher the taxi’s current  $(x,y)$  location. The taxi is dispatched to a client if there is an outstanding request from a client; otherwise, the taxi blocks until a client request is made. The taxi’s call to `getClient` returns with the  $(x,y)$  arguments changed to the next client’s location.

The dispatcher creates a pool of 5 taxis, and dispatches the nearest taxi to the client’s address to minimize client waiting time, if a taxi is available. The taxi constructor is

```
Taxi( MapleLeafTaxi & employer, int id );
```

Routine `nearestTaxi` is used by the dispatcher to find the nearest available taxi address to the given client address. (Do not write `nearestTaxi`; just use the member interface in Figure 1.)

When the dispatcher’s `close` routine is called at the end of the day, it prints out the message “Closed for the day”, and stops accepting calls to `getTaxi`. (Assume no outstanding client requests for a taxi after this point.) Then, the dispatcher must deal with outstanding futures to clients that cannot be serviced (no cab will come to service their request), and it must wait for all the taxis to check in before telling them to go home so they can be deleted. Any client with an outstanding taxi-future has the exception `Closed` inserted into the future, so the client gets this exception raised when it accesses the future. Any waiting or arriving taxi has the exception `Closed` raised on its stack. The dispatcher must delete any allocated storage before terminating. Ensure the dispatcher task does as much administration works as possible; a monitor-style solution will receive little or no marks. Write the code for `MapleLeafTaxi::main` and any necessary declarations/initializations; do **NOT** write the client, taxi, or program main. Assume the program main creates and deletes all the necessary tasks, appropriately, and calls the dispatcher’s `close` routine.  $\mu$ C++ future server operations are:

- `delivery( T result )` - copy result to be returned to the client(s) into the future, unblocking clients waiting for the result.
- `delivery( uBaseEvent *cause )` - copy a server-generated exception into the future, and the exception cause is thrown at clients accessing the future.

The C++ list operations are:

$\mu$ C++ uCondition operations	C++ list operations
<b>bool</b> empty() // true if nobody blocked	<b>int</b> size() // list size
<b>void</b> wait() // wait on condition	<b>bool</b> empty() // size() == 0
<b>void</b> wait( int info ) // wait with info	T front() // first element
<b>bool</b> signal() // signal condition	T back() // last element
<b>bool</b> signalBlock() // signal condition	<b>void</b> push_front( const T & x ) // add x before first element
<b>int</b> front() // return front element info	<b>void</b> push_back( const T & x ) // add x after last element
	<b>void</b> pop_front() // remove first element
	<b>void</b> pop_back() // remove last element
	<b>void</b> clear() // erase all elements

```

_Task MapleLeafTaxiDispatcher {
public:
  _Event Closed {}; // indicate MapleLeafTaxi closed
  typedef Future_ISM<int> Ftaxi; // future taxi
private:
  struct LocnClient {
    int id, x, y; // client id and location coordinates
    Ftaxi ftaxi; // future returned to client
    LocnClient( int id, int x, int y ) : id( id ), x( x ), y( y ) {}
  };
  struct LocnTaxi {
    int id, x, y; // taxi id and location coordinates
    uCondition idle;
    LocnTaxi( int id, int x, int y ) : id( id ), x( x ), y( y ) {}
  };
  enum { NoOfTaxi = 5 };
  list<LocnClient *> clients; // client requests for taxis
  list<LocnTaxi *> taxis; // waiting taxis
  int xclient, yclient; // communication variables
  bool closed = false;
public:
  Ftaxi getTaxi( int id, int x, int y ) { // called by client
    LocnClient *client = new LocnClient( id, x, y ); // create work request
    clients.push_back( client ); // add to request list
    return client->ftaxi; // return future request
  }
  void getClient( int id, int & x, int & y ) { // called by taxi, x,y in/out parameters
    LocnTaxi taxi( id, x, y ); // use the stack!
    taxis.push_back( &taxi ); // add to waiting taxi list
    taxi.idle.wait(); // taxi always blocks
    if ( closed ) _Throw Closed();
    x = xclient; y = yclient; // taxi returns client info
  }
  void close() {} // called at closing time
private:
  list<LocnTaxi *>::iterator nearestTaxi( LocnClient * node, list<LocnTaxi *> & alist ) {
    // Find the element in parameter alist whose address is closest
    // to the address in parameter node. "alist" must be non-empty.
    // ASSUME THIS ROUTINE IS WRITTEN; DO NOT WRITE IT.
  }
  void main() {
    // YOU WRITE ONLY THIS ROUTINE!!!!
    // allocate taxi tasks
    // dispatch taxis to clients, until close called
    // print closed message
    // mark outstanding futures with Closed exception
    // tell each taxi to go home
    // delete taxi tasks
  }
};

```

Figure 1: Maple Leaf Taxi Dispatcher-Administrator