

CS 348 Lecture 6

SQL Part 3

Semih Salihoğlu

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UNIVERSITY OF
WATERLOO



SQL features covered so far

- Query
 - SELECT-FROM-WHERE statements
 - Set and bag operations
 - Subqueries
 - Aggregation and grouping
 - Ordering & LIMIT
 - Outerjoins (and NULL)
- Modification
 - INSERT/DELETE/UPDATE

👉 Today: Constraints, schema changes

Constraints

- Restricts what data is allowed in a database
 - In addition to the simple structure and type restrictions imposed by the table definitions
- Why use constraints?
 - Protect data integrity (catch errors)
 - Tell the DBMS about the data (so it can optimize better)
- Declared as **part of the schema** and enforced by the DBMS

Types of SQL constraints

- NOT NULL
- Key
- Referential integrity (foreign key)
- Tuple- and attribute-based CHECK's
- Another one: “General assertion” is also in the standard but not implemented in SQL systems

NOT NULL constraint examples

```
CREATE TABLE User  
(uid INT NOT NULL,  
name VARCHAR(30) NOT NULL,  
twitterid VARCHAR(15) NOT NULL,  
age INT,  
pop DECIMAL(3,2));
```

```
CREATE TABLE Group  
(gid CHAR(10) NOT NULL,  
name VARCHAR(100) NOT NULL);
```

```
CREATE TABLE Member  
(uid INT NOT NULL,  
gid CHAR(10) NOT NULL);
```

Key declaration examples

```
CREATE TABLE User
(uid INT NOT NULL PRIMARY KEY,
name VARCHAR(30) NOT NULL,
twitterid VARCHAR(15) NOT NULL UNIQUE,
age INT,
pop DECIMAL(3,2));
```

```
CREATE TABLE Group
(gid CHAR(10) NOT NULL PRIMARY KEY,
name VARCHAR(100) NOT NULL);
```

```
CREATE TABLE Member
(uid INT NOT NULL,
gid CHAR(10) NOT NULL,
PRIMARY KEY(uid,gid));
```

```
CREATE TABLE Member
(uid INT NOT NULL PRIMARY KEY,
gid CHAR(10) NOT NULL PRIMARY KEY,
```

Some systems allow PKs to be NULL (e.g., Sqlite), in which case often multiple nulls since null != null; others don't (e.g., Postgres).

At most one primary key per table

Any number of UNIQUE keys per table

Systems generally allow UNIQUE constraints to contain nulls (and multiple nulls because null != null)

This form is required for multi-attribute keys

Incorrect!

Key declaration examples

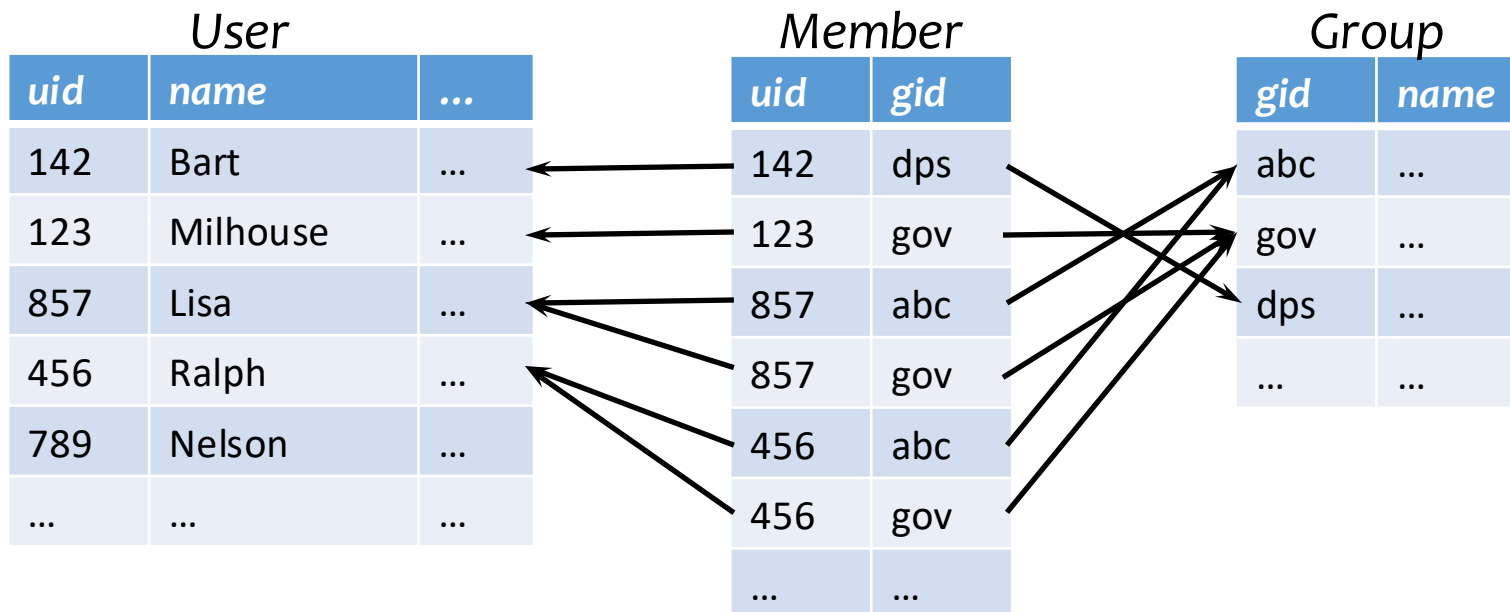
```
CREATE TABLE USCIT  
(uid INT NOT NULL PRIMARY KEY,  
name VARCHAR(30) NOT NULL,  
twitterid VARCHAR(15) NOT NULL,  
age INT,  
pop DECIMAL(3,2),  
UNIQUE (name,age));
```

Similarly, This form is required for multi-attribute unique constraints

Referential integrity example

- If a *uid* appears in *Member*, it must appear in *User*
 - *Member.uid* references *User.uid*
- If a *gid* appears in *Member*, it must appear in *Group*
 - *Member.gid* references *Group.gid*

☞ That is, no “dangling pointers”



Referential integrity in SQL

- Referenced column(s) must be **PRIMARY KEY**
- Referencing column(s) form a **FOREIGN KEY**
- Example

Some system allow them to be non-PK but must be **UNIQUE**

```
CREATE TABLE Member
(uid INT NOT NULL REFERENCES User(uid),
gid VARCHAR(10) NOT NULL,
PRIMARY KEY(uid,gid),
FOREIGN KEY (gid) REFERENCES Group(gid));
```

This form is required for multi-attribute foreign keys

```
CREATE TABLE MemberBenefits
(.....
FOREIGN KEY (uid,gid) REFERENCES Member(uid,gid));
```

Enforcing referential integrity

Example: *Member.uid* references *User.uid*

- Insert or update a *Member* row so it **refers to a non-existent uid**
 - **Reject**

User			Member	
uid	name	...	uid	gid
142	Bart	...	142	dps
123	Milhouse	...	123	gov
857	Lisa	...	857	abc
456	Ralph	...	857	gov
789	Nelson	...	456	abc
...	456	gov
			000	gov

Reject

Enforcing referential integrity

Example: *Member.uid* references *User.uid*

- Delete or update a *User* row whose *uid* is referenced by some *Member* row
 - Multiple Options (in SQL)

User			Member	
uid	name	...	uid	gid
142	Bart	...	142	dps
123	Milhouse	...	123	gov
857	Lisa	...	857	abc
456	Ralph	...	857	gov
789	Nelson	...	456	abc
...	456	gov
...

Option 1: Reject

```
CREATE TABLE Member  
(uid INT NOT NULL  
REFERENCES User(uid)  
ON DELETE CASCADE,  
.....);
```

Option 2: Cascade
(ripple changes to all referring rows)

Enforcing referential integrity

Example: *Member.uid* references *User.uid*

- Delete or update a *User* row whose *uid* is referenced by some *Member* row
 - Multiple Options (in SQL)

User			Member	
uid	name	...	uid	gid
142	Bart	...	142	dps
123	Milhouse	...	123	gov
857	Lisa	...	857	abc
456	Ralph	...	857	gov
789	Nelson	...	NULL	abc
...	NULL	gov
...

```
CREATE TABLE Member  
(uid INT NOT NULL  
REFERENCES User(uid)  
ON DELETE SET NULL,  
... ..);
```

Option 3: Set NULL
(set all references to NULL)

Deferred constraint checking

- Example:

```
CREATE TABLE Dept
(name CHAR(20) NOT NULL PRIMARY KEY,
chair CHAR(30) NOT NULL
REFERENCES Prof(name));
```

```
CREATE TABLE Prof
(name CHAR(30) NOT NULL PRIMARY KEY,
dept CHAR(20) NOT NULL
REFERENCES Dept(name));
```

- The first INSERT will always violate a constraint!
- **Deferred constraint checking** is necessary
 - Check only at the end of a set of operations (transactions)
 - Allowed in SQL as an option
 - Use keyword **deferred**

Tuple- and attribute-based CHECK's

- Can be put on a single table:
 - Syntax: Check (P), where P is a boolean expression *that must be true for each tuple (i.e., checked per tuple)*
 - Either placed next to an attribute
 - Or at the end of table definition as a separate statement
- **Only checked when a tuple is inserted or modified**
 - Reject if P evaluates to FALSE
 - TRUE and UNKNOWN are fine
- SQL Standard: P is arbitrary and can contain sub-queries.
- In practice: SQL systems do not allow sub-queries

Tuple- and attribute-based CHECK's

- Useful to put domain constraints or correlate multiple attributes of the same tuple

- Examples:

```
CREATE TABLE User(...  
  age INTEGER CHECK(age > 0),  
  ...);
```

```
CREATE TABLE Products(...  
  pID INTEGER,  
  price INTEGER,  
  discountedPrice INTEGER,  
  CHECK(price <= discountedPrice));
```

Tuple- and attribute-based CHECK's

- Can specify complex constraints if sub-queries are supported (but again any system I know of does not)
- Reasoning about complex CHECK constraints can be hard:
- E.g:

Should you check when User is updated?
According to SQL Standard:
Checked when new tuples are added to Member but not when User is modified

```
CREATE TABLE Member  
(uid INTEGER NOT NULL,  
CHECK(uid IN (SELECT uid FROM User)),  
...);
```

- Similarly, if the sub-query contains complex joins interpreting the behaviour can be hard

General assertion (Optional)

- Also only in the SQL standard; Not supported in systems
- `CREATE ASSERTION assertion_name
CHECK assertion_condition;`
- *assertion_condition* is checked for each modification that could potentially violate it
- Example: *Member.uid* references *User.uid*

```
CREATE ASSERTION MemberUserRefIntegrity  
CHECK (NOT EXISTS  
      (SELECT * FROM Member  
       WHERE uid NOT IN  
         (SELECT uid FROM User)));
```

Can include
multiple
tables

Assertions are
statements
that must
always be true

Naming constraints

- It is possible to name constraints (similar to assertions)

```
CREATE TABLE User(...  
  age INT, constraint minAge check(age > 0),  
  ...);
```

Exercises

Consider this db instance:

<i>uid</i>	<i>gid</i>
857	dps
123	gov
857	abc
857	gov
456	abc
456	gov

<i>uid</i>	<i>gid</i>	<i>discount</i>
857	dps	10
123	gov	25
857	abc	5

- *MemberBenefits* table references the *Member* table
- (*uid, gid*) forms the primary key of *MemberBenefits* table
- Assume *discount* is of type INT (and *uid* is INT and *gid* is string with a max of 30 characters)
- Write a DDL to create the *MemberBenefits* table

Exercises

Consider this db instance:

uid	gid
857	dps
123	gov
857	abc
857	gov
456	abc
456	gov

uid	gid	discount
857	dps	10
123	gov	25
857	abc	5

- *MemberBenefits* table references the *Member* table
- (*uid, gid*) forms the primary key of *MemberBenefits* table
- Assume *discount* is of type INT (and *uid* is INT and *gid* is string with a max of 30 characters)

```
CREATE TABLE MemberBenefits
(uid INT,
gid VARCHAR(30),
discount INT,
PRIMARY KEY (uid,gid),
FOREIGN KEY (uid,gid) REFERENCES Member(uid,gid));
```

Exercises

Consider this db instance:

<i>uid</i>	<i>gid</i>
857	dps
123	gov
857	abc
857	gov
456	abc
456	gov

- Assume all foreign key references are set to **ON DELETE SET NULL**
- (Assume the db allows this, just for this exercise)
- What happens when user 857 is deleted from the *User* table? (Recall *Member* table references *uid* of *User* table)

Exercise

User (uid int, name string, age int, pop float)
Group (gid string, name string)
Member (uid int, gid string)

- Assume the User table requires *pop* column values to be between 0 and 1. Complete the following DDL statement.

```
CREATE TABLE User
(uid INT PRIMARY KEY,
name VARCHAR(30),
age INT,
pop DECIMAL(3,2) ???);
```

Exercise.

User (uid int, name string, age int, pop float)
Group (gid string, name string)
Member (uid int, gid string)

- Assume the User table requires *pop* column values to be between 0 and 1 or NULL. Complete the following DDL statement.

```
CREATE TABLE User
(uid INT PRIMARY KEY,
 name VARCHAR(30),
 age INT,
 pop DECIMAL(3,2) CHECK(pop IS NULL OR (pop >= 0 AND pop < 1)));
```

Take home ex.

User (uid int, name string, age int, pop float)
Group (gid string, name string)
Member (uid int, gid string)

- Say every user with $\text{pop} \geq 0.9$ must belong to the Book Club ($\text{gid} = \text{'abc'}$). Create an assertion to check this constraint.

Schema modification

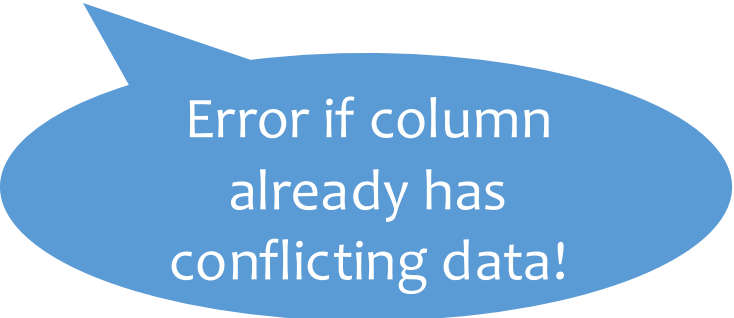
- How to add constraints once the schema is defined??
- Add or Modify attributes/domains
- Add or Remove constraints

Add or Modify attributes/domains

- *Alter table* table_name *Add column* column_name
- *Alter table* table_name *Rename column* old_name to new_name
- *Alter table* table_name *Drop column* column_name

Domain change:

- *Alter table* table_name *Alter column* column_name datatype



Error if column
already has
conflicting data!

Add or Remove constraints

- *Alter table* `table_name` *Add constraint*
`constraint_name constraint_condition`

```
ALTER TABLE Member  
ADD CONSTRAINT fk_user FOREIGN KEY(uid)  
REFERENCES User(uid)
```

- *Alter table* `table_name` *Drop constraint*
`constraint_name`

```
ALTER TABLE Member  
DROP CONSTRAINT fk_user
```

SQL

- Constraints
- Schema changes
- Triggers (Optional)

Note: The rest of these slides on triggers is optional material.

It is presented here to expose you to how much application logic you can push into SQL systems using a rule-based approach called triggers. You will not be tested on triggers.

Recall “referential integrity”

Example: *Member.uid* references *User.uid*

- Delete or update a *User* row whose *uid* is referenced by some *Member* row
 - Multiple Options (in SQL)

User			Member	
uid	name	...	uid	gid
142	Bart	...	142	dps
123	Milhouse	...	123	gov
857	Lisa	...	857	abc
456	Ralph	...	857	gov
789	Nelson	...	456	abc
...	456	gov
...

Diagram illustrating referential integrity. Arrows point from Member rows to User rows: 142 Member to Bart User, 123 Member to Milhouse User, 857 Member to Lisa User, 857 Member to Ralph User, 456 Member to Nelson User, and 456 Member to Ralph User. The User row for uid 456 (Ralph) and the Member rows for uid 456 (two rows) are crossed out with red lines. The text "Option 1: Reject" is overlaid on the User row for uid 456.

```
CREATE TABLE Member
(uid INT NOT NULL
REFERENCES User(uid)
ON DELETE CASCADE,
....);
```

Option 2: Cascade
(ripple changes to all referring rows)

Can we generalize it?

Referential constraints

Delete/update a
User row

Whether its uid is
referenced by some
Member row

If yes: reject/ delete
cascade/null

Event

Condition

Action

Data Monitoring

Some user's
popularity is updated

Whether the user is a
member of "Pop group"
and pop drops below 0.5

If yes: kick that user out
of Pop group!

Triggers

- A **trigger** is an event-condition-action (ECA) rule
 - When **event** occurs, test **condition**; if condition is satisfied, execute **action**

```
CREATE TRIGGER PickyPopGroup
AFTER UPDATE OF pop ON User
REFERENCING NEW ROW AS newUser
FOR EACH ROW
  WHEN (newUser.pop < 0.5)
    AND (newUser.uid IN (SELECT uid
                        FROM Member
                        WHERE gid = 'popgroup'))
  DELETE FROM Member
  WHERE uid = newUser.uid AND gid = 'popgroup'
```

The diagram illustrates the components of the SQL trigger code. Callouts point to specific parts of the code:

- Event:** Points to the `UPDATE OF pop` clause.
- Transition variable:** Points to the `newUser` alias in the `REFERENCING` clause.
- Condition:** Points to the `WHEN` clause.
- Action:** Points to the `DELETE FROM Member` clause.

Trigger option 1 – possible events

- Possible events include:
 - **INSERT ON** *table*; **DELETE ON** *table*; **UPDATE** [**OF** *column*]
ON *table*

```
CREATE TRIGGER PickyPopGroup
AFTER UPDATE OF pop ON User
REFERENCING NEW ROW AS newUser
FOR EACH ROW
    WHEN (newUser.pop < 0.5)
        AND (newUser.uid IN (SELECT uid
                              FROM Member
                              WHERE gid = 'popgroup'))
    DELETE FROM Member
    WHERE uid = newUser.uid AND gid = 'popgroup';
```

Event

Condition

Action

Trigger option 2 – timing

- Timing—action can be executed:
 - **AFTER** or **BEFORE** the triggering event
 - **INSTEAD OF** the triggering event on views (more later)

```
CREATE TRIGGER NoFountainOfYouth
BEFORE UPDATE OF age ON User
REFERENCING OLD ROW AS o, NEW ROW AS n
FOR EACH ROW
  WHEN (n.age < o.age)
    SET n.age = o.age;
```

The diagram illustrates the components of a SQL trigger. It shows a code snippet with three callout boxes: 'Event' pointing to 'UPDATE OF age ON User', 'Condition' pointing to '(n.age < o.age)', and 'Action' pointing to 'SET n.age = o.age;'. The words 'OLD ROW AS o' and 'NEW ROW AS n' are underlined in the original image.

Trigger option 3 – granularity

- Granularity—trigger can be activated:
 - **FOR EACH ROW** modified

```
CREATE TRIGGER PickyPopGroup
AFTER UPDATE OF pop ON User
REFERENCING NEW ROW AS newUser
FOR EACH ROW
  WHEN (newUser.pop < 0.5)
    AND (newUser.uid IN (SELECT uid
                        FROM Member
                        WHERE gid = 'popgroup'))
  DELETE FROM Member
  WHERE uid = newUser.uid AND gid = 'popgroup';
```

Event

Condition

Action

Trigger option 3 – granularity

- Granularity—trigger can be activated:
 - **FOR EACH ROW** modified
 - **FOR EACH STATEMENT** that performs modification

```
CREATE TRIGGER PickyPopGroup2
AFTER UPDATE OF pop ON User
REFERENCING NEW TABLE AS newUsers
FOR EACH STATEMENT
DELETE FROM Member
WHERE gid = 'popgroup'
AND uid IN (SELECT uid
FROM newUsers
WHERE pop < 0.5);
```

Event

Transition table:
contains all the
affected rows

Condition
& Action

Trigger option 3 – granularity

- Granularity—trigger can be activated:
 - **FOR EACH ROW** modified
 - **FOR EACH STATEMENT** that performs modification

```
CREATE TRIGGER PickyPopGroup2
AFTER UPDATE OF pop ON User
REFERENCING NEW TABLE AS newUsers
FOR EACH STATEMENT
    DELETE FROM Member
        WHERE gid = 'popgroup'
        AND uid IN (SELECT uid
                    FROM newUsers
                    WHERE pop < 0.5);
```

Transition table:
contains all the
affected rows

Can only be used
with **AFTER**
triggers

Transition variables/tables

- **OLD ROW**: the modified row before the triggering event
- **NEW ROW**: the modified row after the triggering event
- **OLD TABLE**: a read-only table containing all old rows modified by the triggering event
- **NEW TABLE**: a table containing all modified rows after the triggering event

Event	Row	Statement
Delete	old r; old t	old t
Insert	new r; new t	new t
Update	old/new r; old/new t	old/new t

AFTER Trigger

Event	Row	Statement
Update	old/new r	-
Insert	new r	-
Delete	old r	-

BEFORE Trigger

Statement- vs. row-level triggers

- Simple row-level triggers are easier to implement
 - Statement-level triggers: require significant amount of state to be maintained in OLD TABLE and NEW TABLE
- However, in some cases a row-level trigger may be less efficient
 - E.g., 4B rows and a trigger may affect 10% of the rows. Recording an action for 4 Million rows, one at a time, is not feasible due to resource constraints.
- Certain triggers are only possible at statement level
 - E.g., ??

Certain triggers are only possible at statement level

```
CREATE TRIGGER MaintainAvgPop
AFTER UPDATE OF pop ON User
REFERENCING NEW TABLE AS newUsers
          OLD TABLE AS oldUsers
FOR EACH STATEMENT
  WHEN (0.5 > (SELECT AVG(pop) from User))
  BEGIN
    DELETE FROM User WHERE uid IN (SELECT uid
    FROM newUsers)
    INSERT INTO User (SELECT * FROM oldUsers)
  END
```

The diagram illustrates the components of the SQL trigger:

- Event:** `UPDATE OF pop ON User`
- Transition tables:** `newUsers` (NEW TABLE) and `oldUsers` (OLD TABLE)
- Condition:** `0.5 > (SELECT AVG(pop) from User)`
- Action:** `DELETE FROM User WHERE uid IN (SELECT uid FROM newUsers)` and `INSERT INTO User (SELECT * FROM oldUsers)`

System issues

- Recursive firing of triggers
 - Action of one trigger causes another trigger to fire
 - Can get into an infinite loop
- Interaction with constraints (tricky to get right!)
 - When to check if a triggering event violates constraints?
 - After a BEFORE trigger
 - Before an AFTER trigger
 - (based on db2, other DBMS may differ)
- Best to avoid when alternatives exist

SQL features covered so far

Basic & Intermediate SQL

- Query
- Modification
- Constraints
- Triggers

👉 Next: Views, Indexes, Programming & recursion