SQL: Triggers, Views, Indexes

Introduction to Database Management CS348 Fall 2022

Announcements (Tue., Sep 27)

- Milestone o Project groups are formed by tonight!
 - Form a team on Learn
 - Report.pdf and link to GitHub repo
 - Not graded, but very important!

- Assignment #1 due by Sep 29 (Thur), 11:59pm
 - Part 1: general questions and r.a.
 - Submit via Crowdmark
 - Part 2: writing SQL on DB2 on school servers (try soon)
 - Submit via Marmoset

SQL

• Basic SQL (queries, modifications, and constraints)

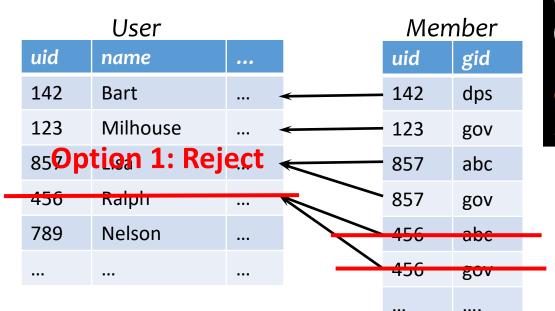
- Intermediate SQL
 - Triggers
 - Views
 - Indexes
- Advanced SQL
 - Programming
 - Recursive queries (Optional)

Lectures 5-6

Still remember "referential integrity"?

Example: Member.uid references User.uid

- Delete or update a User row whose uid is referenced by some Member row
 - Multiple Options (in SQL)



CREATE TABLE Member
(uid DECIMAL(3,0) NOT NULL
REFERENCES User(uid)
ON DELETE CASCADE,
....):

Option 2: Cascade (ripple changes to all referring rows)

Can we generalize it?

Delete/update a User row

Whether its uid is referenced by some Member row

Yes: reject/delete cascade/null

Condition
Action

Some user's popularity is updated

Whether the user is a member of "S group" and pop drops below 0.5

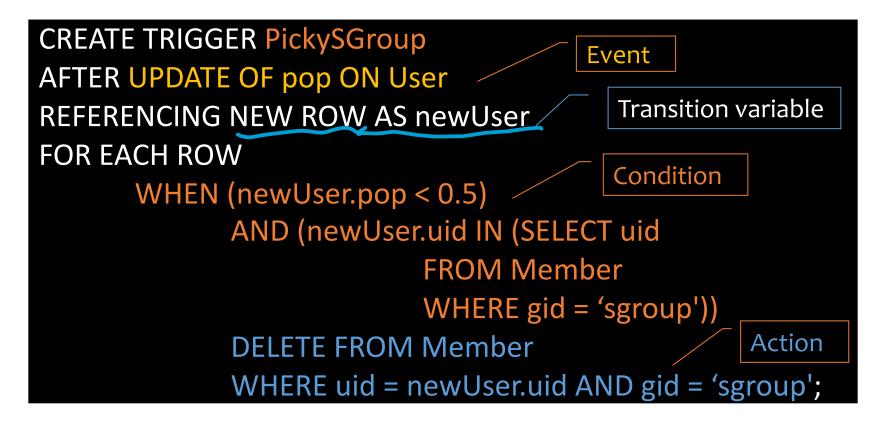
Yes: kick that user out of S group

Referential constraints

Data Monitoring

Triggers

- A trigger is an event-condition-action (ECA) rule
 - When event occurs, test condition; if condition is satisfied, execute action



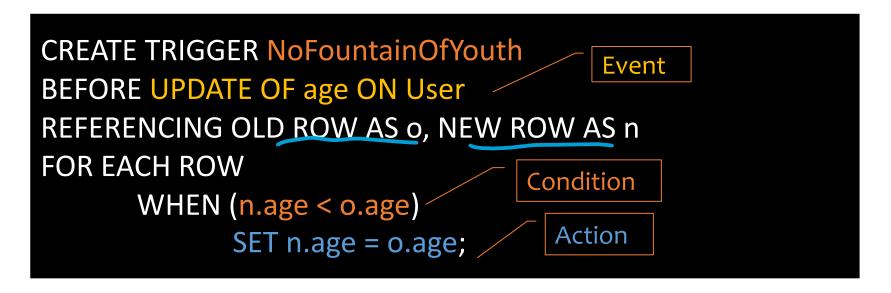
Trigger option 1 – possible events

- Possible events include:
 - INSERT ON table; DELETE ON table; UPDATE [OF column] ON table

```
CREATE TRIGGER PickySGroup
                                      Event
AFTER UPDATE OF pop ON User
REFERENCING NEW ROW AS new User
FOR EACH ROW
                                        Condition
      WHEN (newUser.pop < 0.5)
             AND (newUser.uid IN (SELECT uid
                           FROM Member
                          WHERE gid = 'sgroup'))
             DELETE FROM Member
                                                 Action
             WHERE uid = newUser.uid AND gid = 'sgroup';
```

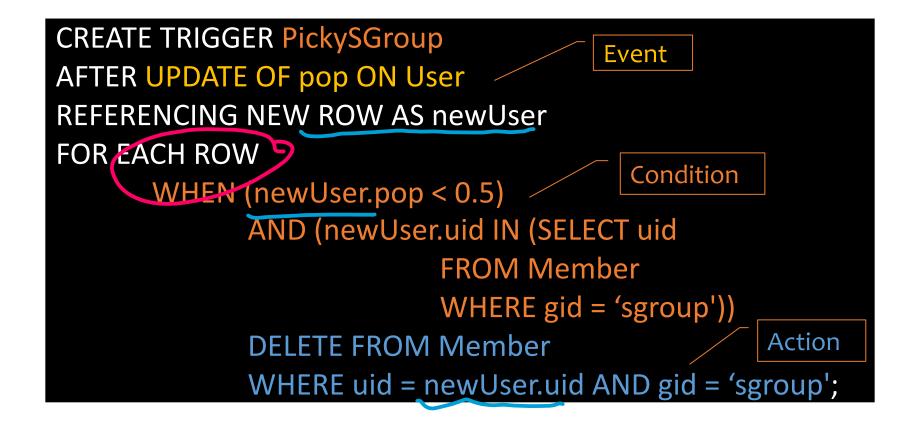
Trigger option 2 – timing

- Timing—action can be executed:
 - AFTER or BEFORE the triggering event
 - INSTEAD OF the triggering event on views (more later)



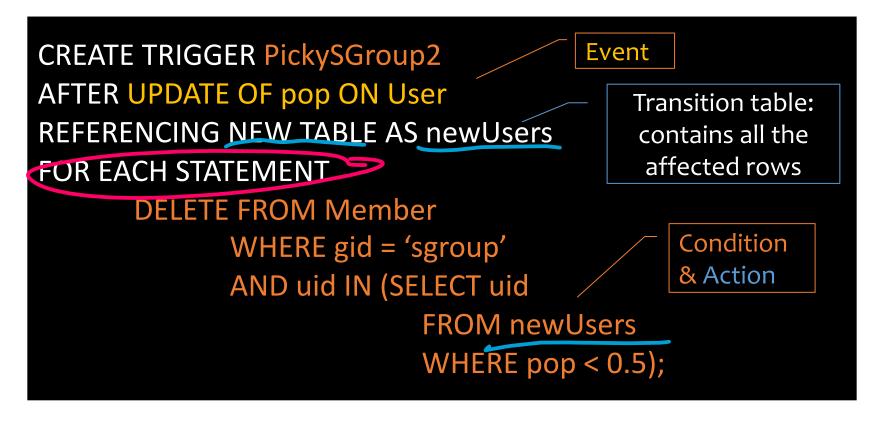
Trigger option 3 – granularity

- Granularity—trigger can be activated:
 - FOR EACH ROW modified



Trigger option 3 – granularity

- Granularity—trigger can be activated:
 - FOR EACH ROW modified
 - FOR EACH STATEMENT that performs modification



Trigger option 3 – granularity

- Granularity—trigger can be activated:
 - FOR EACH ROW modified
 - FOR EACH STATEMENT that performs modification

CREATE TRIGGER PickySGroup2 AFTER UPDATE OF pop ON User Transition table: REFERENCING NEW TABLE AS newUsers contains all the affected rows FOR EACH STATEMENT **DELETE FROM Member** Only can be used with **AFTER** WHERE gid = 'sgroup' triggers AND uid IN (SELECT uid FROM newUsers WHERE pop < 0.5);

Transition variables/tables

- OLD ROW: the modified row before the triggering event
- NEW ROW: the modified row after the triggering event
- OLD TABLE: a hypothetical read-only table containing all rows to be modified before the triggering event
- NEW TABLE: a hypothetical table containing all modified rows after the triggering event

Event	Row	Statement
Delete	old r; old t	old t
Insert	new r; new t	new t
Update	old/new r; old/new t	old/new t

AFTER Trigger

Event	Row	Statement
Update	old/new r	-
Insert	new r	-
Delete	old r	-

BEFORE Trigger

Statement- vs. row-level triggers

- Simple row-level triggers are easier to implement
 - Statement-level triggers: require significant amount of state to be maintained in OLD TABLE and NEW TABLE
- Exercise 1: However, can you think of a case when a row-level trigger may be less efficient?
- Exercise 2: Certain triggers are only possible at statement level. Can you think of an example?

System issues

- Recursive firing of triggers
 - Action of one trigger causes another trigger to fire
 - Can get into an infinite loop
- Interaction with constraints (tricky to get right!)
 - When to check if a triggering event violates constraints?
 - After a BEFORE trigger
 - Before an AFTER trigger
 - (based on db2, other DBMS may differ)
- Be best avoided when alternatives exist

SQL features covered so far

Basic SQL

- Intermediate SQL
 - Triggers
 - Views

Views

- A view is like a "virtual" table
 - Defined by a query, which describes how to compute the view contents on the fly
 - Stored by DBMS instead of view contents
 - Can be used in queries just like a regular table

```
CREATE VIEW SGroup AS

SELECT * FROM User

WHERE uid IN (SELECT uid

FROM Member

WHERE gid = 'sgroup');
```

SELECT AVG(pop) FROM SGroup;

SELECT MIN(pop) FROM SGroup;

SELECT ... FROM SGroup;

```
SELECT AVG(pop)

FROM (SELECT * FROM User

WHERE uid IN

(SELECT uid FROM Member

WHERE gid = 'jes'))

AS SGroup;
```

DROP VIEW SGroup;

Why use views?

- To hide complexity from users
- To hide data from users
- Logical data independence
- To provide a uniform interface for different implementations or sources

Modifying views

Does it even make sense, since views are virtual?

 It does make sense if we want users to really see views as tables

 Goal: modify the base tables such that the modification would appear to have been accomplished on the view

A simple case

CREATE VIEW UserPop AS

SELECT uid, pop FROM User;

DELETE FROM UserPop WHERE uid = 123;

translates to:

DELETE FROM User WHERE uid = 123;

An impossible case

CREATE VIEW PopularUser AS

SELECT uid, pop FROM User

WHERE pop >= 0.8;

INSERT INTO PopularUser VALUES(987, 0.3);

 No matter what we do on User, the inserted row will not be in PopularUser

A case with too many possibilities

CREATE VIEW AveragePop(pop) AS

SELECT AVG(pop) FROM User;

Renamed column

UPDATE AveragePop SET pop = 0.5;

- Set everybody's pop to 0.5?
- Adjust everybody's pop by the same amount?
- Just lower one user's pop?

SQL92 updateable views

- More or less just single-table selection queries
 - No join
 - No aggregation
 - No subqueries
- Arguably somewhat restrictive
- Still might get it wrong in some cases
 - See the slide titled "An impossible case" (slide 20)
 - Adding WITH CHECK OPTION to the end of the view definition will make DBMS reject such modifications

INSTEAD OF triggers for views

CREATE VIEW AveragePop(pop) AS

SELECT AVG(pop) FROM User;

CREATE TRIGGER AdjustAveragePop
INSTEAD OF UPDATE ON AveragePop
REFERENCING OLD ROW AS o,
NEW ROW AS n
FOR EACH ROW
UPDATE User
SET pop = pop + (n.pop-o.pop);

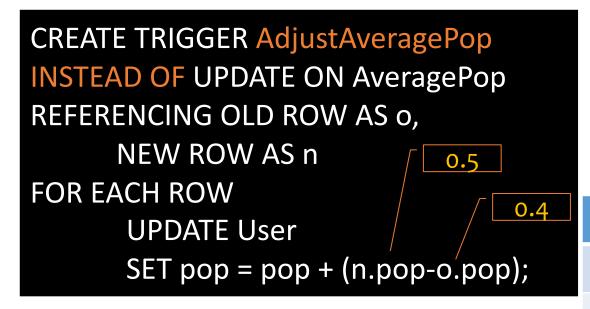
What does this trigger do?

UPDATE AveragePop SET pop = 0.5;

INSTEAD OF triggers for views

CREATE VIEW AveragePop(pop) AS

SELECT AVG(pop) FROM User;



• What does this trigger do?

UPDATE AveragePop SET pop = 0.5;

User

pop

0.4 + 0.1

SQL features covered so far

Basic SQL

- Intermediate SQL
 - Triggers
 - Views
 - Indexes

Motivating examples of using indexes

SELECT * FROM User WHERE name = 'Bart';

- Can we go "directly" to rows with name='Bart' instead of scanning the entire table?
 - → index on User.name

SELECT * FROM User, Member WHERE User.uid = Member.uid AND Member.gid = 'sgroup';

- Can we find relevant Member rows "directly"?
 - → index on Member.gid or (gid, uid)
- For each relevant Member row, can we "directly" look up User rows with matching uid
 - → index on User.uid

Indexes

- An index is an auxiliary persistent data structure
 - Search tree (e.g., B+-tree), lookup table (e.g., hash table), etc.
 - More on indexes later in this course!
- CREATE [UNIQUE] INDEX indexname ON tablename(columnname₁,...,columnname_n);
 - With UNIQUE, the DBMS will also enforce that $\{columnname_1, \dots, columnname_n\}$ is a key of tablename
- DROP INDEX indexname;
- Typically, the DBMS will automatically create indexes for PRIMARY KEY and UNIQUE constraint declarations

Indexes

- An index on R. A can speed up accesses of the form
 - R.A = value
 - R.A > value (sometimes; depending on the index type)
- An index on $(R.A_1, ..., R.A_n)$ can speed up
 - $R.A_1 = value_1 \land \cdots \land R.A_n = value_n$
 - $(R.A_1, ..., R.A_n) > (value_1, ..., value_n)$ (again depends)

Questions (lecture 12):

- Thow about an index on R. A plus another on R. B?
- More indexes = better performance?

SQL features covered so far

Basic & Intermediate SQL

- Query
- Modification
- Constraints
- Triggers
- Views
- Indexes

Next: Programming