Query Optimization

Introduction to Database Management

CS348 Fall 2022

Overview

- Many different ways of processing the same query
  - Scan? Sort? Hash? Use an index?
  - All have different performance characteristics and/or make different assumptions about data

last lecture

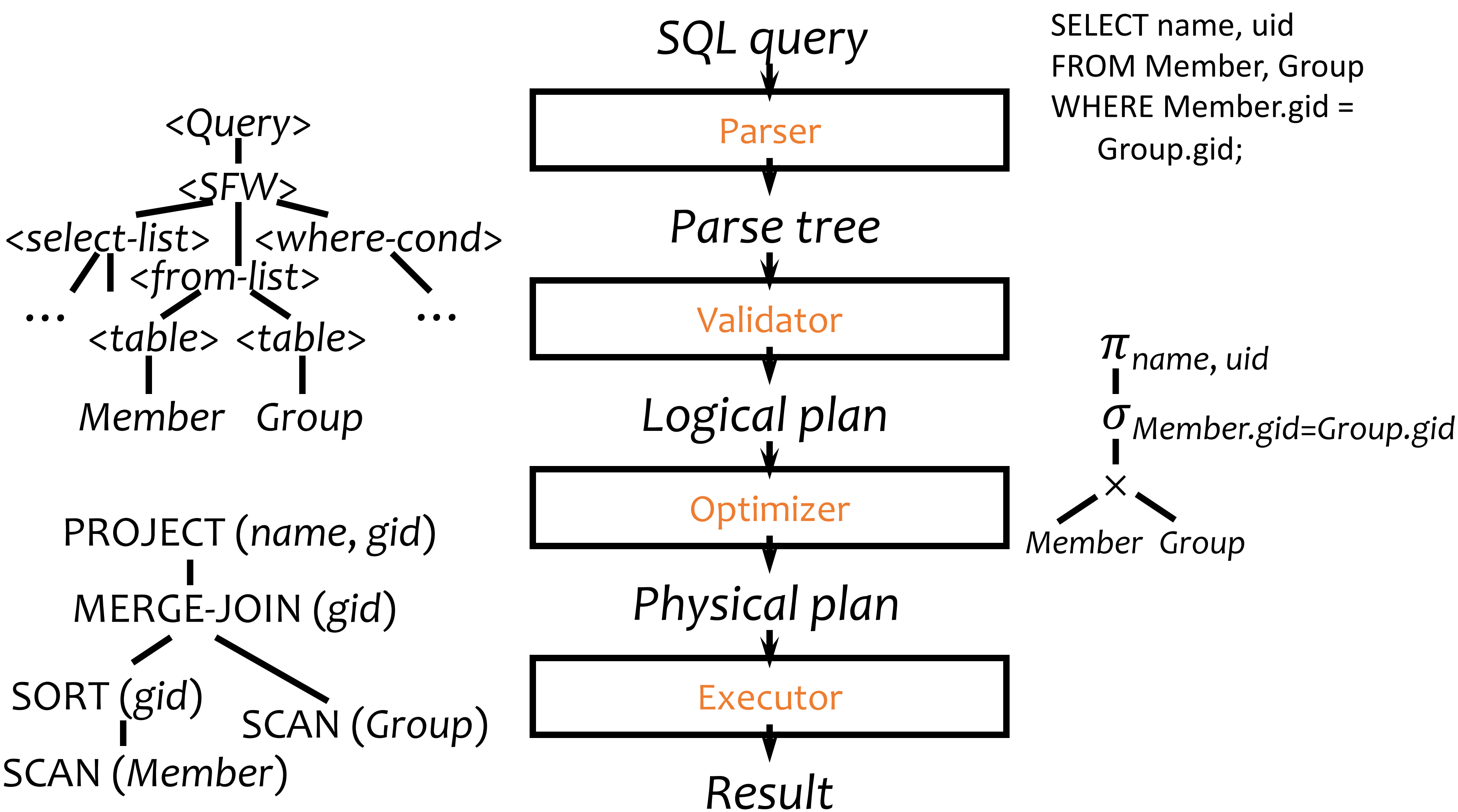
- Best choice depends on the situation
  - Implement all alternatives
  - Let the query optimizer choose at run-time (this lecture)

As some materials (sorting/hashing-based algorithms) are made optional in this term, some part of the edited video may not be smooth.

Outline

- System view of query processing
  - Logical plan and physical plan
- Cost calculation of the physical plan
  - Cardinality estimation
- Search space and search strategy
  - Transformation rules

A query’s trip through the DBMS

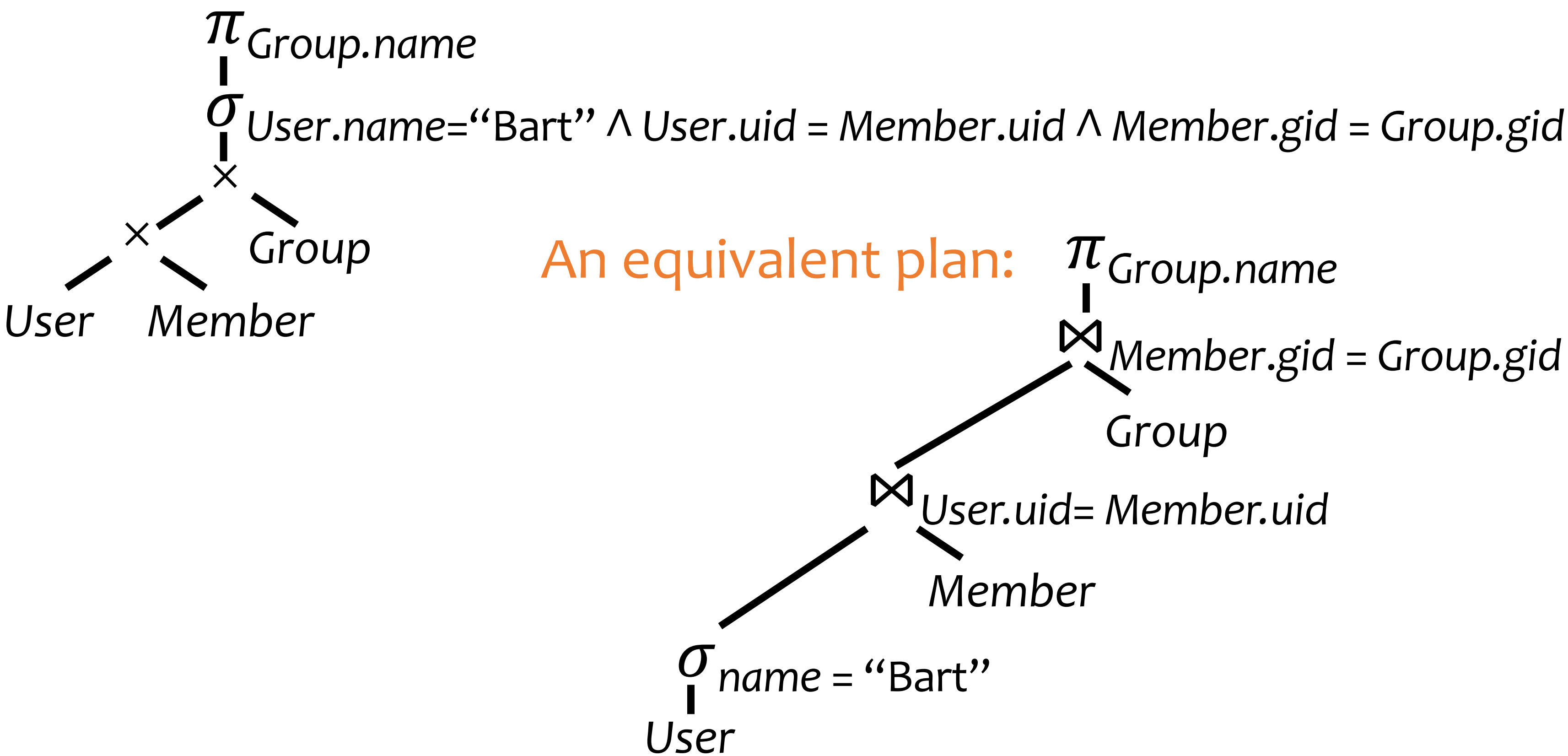


Parsing and validation

- Parser: SQL → parse tree
  - Detect and reject syntax errors
- Validator: parse tree → logical plan
  - Detect and reject semantic errors
    - Nonexistent tables/views/columns?
    - Insufficient access privileges?
    - Type mismatches?
      - Examples: AVG(name), name + pop, User UNION Member
  - Also
    - Expand *
    - Expand view definitions
  - Information required for semantic checking is found in system catalog (which contains all schema information)

Logical plan

- Nodes are logical operators (often relational algebra operators)
- There are many equivalent logical plans



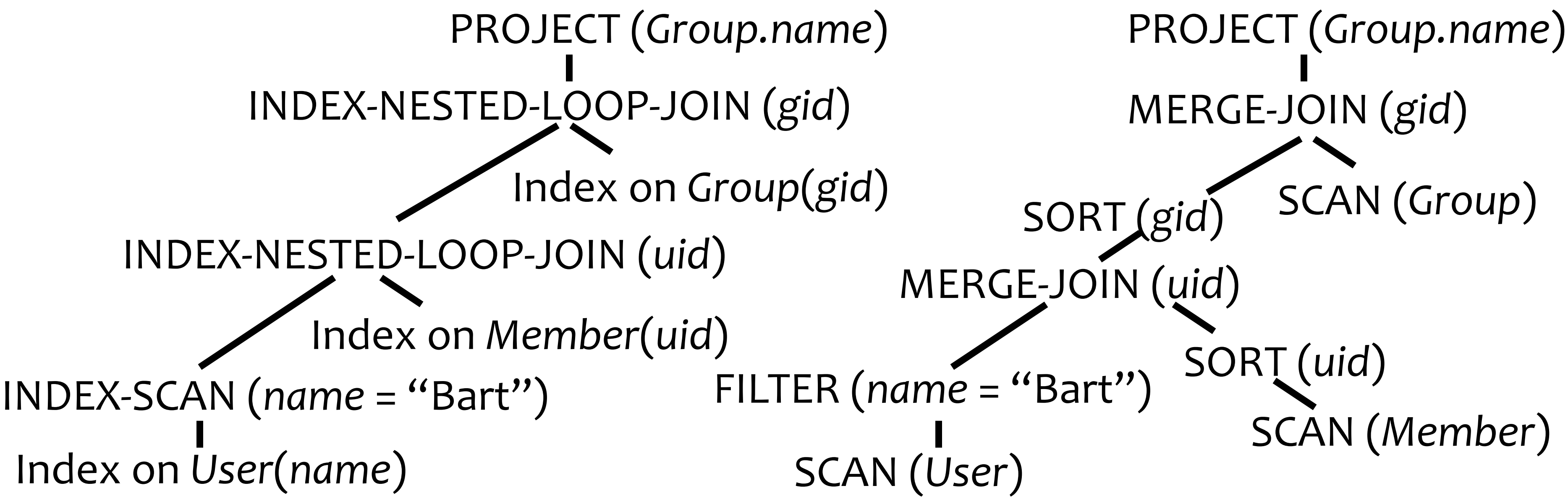
Physical (execution) plan

- A complex query may involve multiple tables and various query processing algorithms
  - E.g., table scan, index nested-loop join, sort-merge join, hash-based duplicate elimination… (Lecture 13)

- A physical plan for a query tells the DBMS query processor how to execute the query
  - A tree of physical plan operators
  - Each operator implements a query processing algorithm
  - Each operator accepts a number of input tables/streams and produces a single output table/stream

Examples of physical plans

SELECT Group.name
FROM User, Member, Group
WHERE User.name = 'Bart'
AND User.uid = Member.uid AND Member.gid = Group.gid;

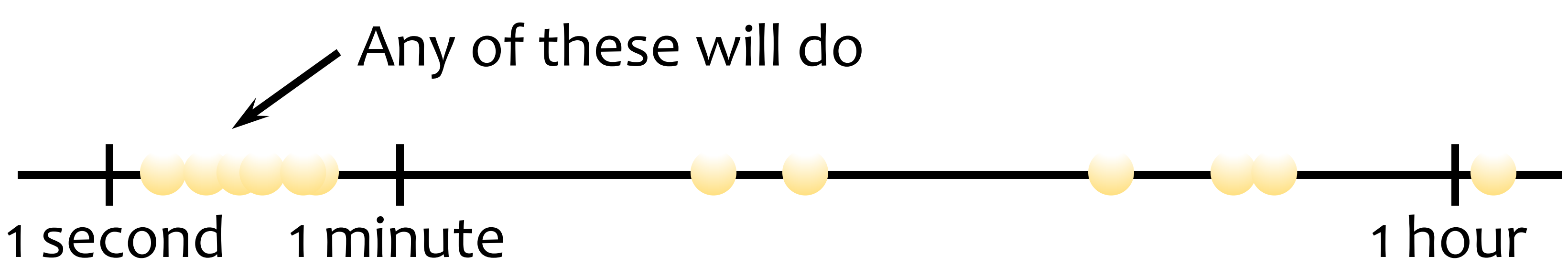


- Many physical plans for a single query
  - Equivalent results, but different costs and assumptions!

  ☞DBMS query optimizer picks the "best" possible physical plan

How to pick the “best” physical plan?

- One logical plan → “best” physical plan
- Questions
  - How to estimate costs
  - How to enumerate possible plans
  - How to pick the “best” one
- Often the goal is not getting the optimum plan, but instead avoiding the horrible ones



Cost estimation

Physical plan example:

PROJECT (*Group.name*)

INDEX-NESTED-LOOP-JOIN (*gid*)

Index on *Group*(*gid*)

Input to Join(*uid*):

What is its input size?

INDEX-NESTED-LOOP-JOIN (*uid*)

Index on *Member*(*uid*)

INDEX-SCAN (*name* = “Bart”)

Index on *User*(*name*)

Lecture 13 (slide 20)

- We have: cost estimation for each operator
  - Example: INDEX-NESTED-LOOP-JOIN(*uid*) takes $O(B(R) + |R| \cdot (\text{index lookup}))$
- We need: size of intermediate results

Cardinality estimation

http://www.learningresources.com/product/estimation+station.do

Selections with equality predicates

- Q: $\sigma_{A=v} R$
- Suppose the following information is available
  - Size of R: $|R|$
  - Number of distinct A values in R: $|\pi_A R|$
- Assumptions
  - Values of A are uniformly distributed in R
  - Values of v in Q are uniformly distributed over all $R.A$ values
- $|Q| \approx {}^{|R|}/_{|\pi_A R|}$
  - Selectivity factor of $(A = v)$ is ${}^{1}/_{|\pi_A R|}$

Example

Physical plan example:

PROJECT (*Group.name*)

INDEX-NESTED-LOOP-JOIN (*gid*)

INDEX-NESTED-LOOP-JOIN (*uid*)

Index on *Group*(*gid*)

INDEX-SCAN (*name* = "Bart")

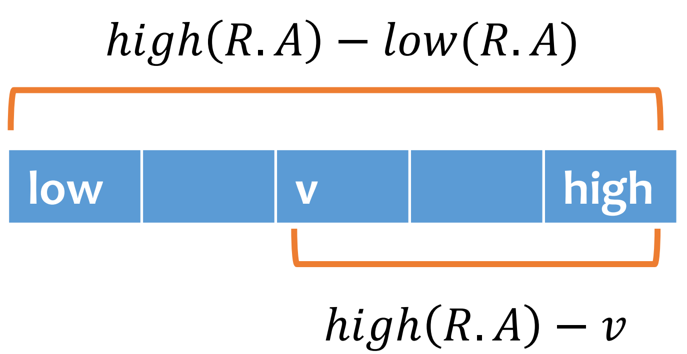
Index on *Member*(*uid*)

Index on *User*(*name*)

- |User|=1000, $|\pi_{name}(User)| = 50$ ➔ $|\sigma_{name="Bart"}(User)| =?$
- Assumptions:
  - Values of $name$ are uniformly distributed in $User$
  - Values of v in $\sigma_{name="Bart"}(User)$ are uniformly distributed over all $\text{User}.name$ values
- $|\sigma_{name="Bart"}(User)| = \frac{1000}{50} = 20$

Range predicates

- Q: $\sigma_{A>v}R$
- Not enough information!
  - Just pick, say, $|Q| \approx |R| \cdot {}^{1}/_{3}$
- With more information
  - Largest *R.A* value: $\mathrm{high}(R.A)$
  - Smallest *R.A* value: $\mathrm{low}(R.A)$
  - $|Q| \approx |R| \cdot \frac{\mathrm{high}(R.A)-v}{\mathrm{high}(R.A)-\mathrm{low}(R.A)}$
  - In practice: sometimes the second highest and lowest are used instead
    - The highest and the lowest are often used by inexperienced database designer to represent invalid values!



Example

- Database:
  - User(<u>uid</u>, name, age, pop), Member(<u>gid,uid</u>,date), Group(<u>gid</u>, gname)
  - |User|=1000 rows, |Group|=100 rows, |Member|=50000 rows
  - $|\pi_{name}(User)| = 50, \pi_{pop}(User) = \{1,2,3,4,5\}$
  - $|\pi_{uid}(Member)| = 900$

- Estimate size $|User \bowtie Member| =?$

Two-way equi-join

- Q: $R(A, B) \bowtie S(A, C)$
- Assumption: containment of value sets
  - Every tuple in the "smaller" relation (one with fewer distinct values for the join attribute) joins with some tuple in the other relation
    - That is, if $|\pi_A R| \leq |\pi_A S|$ then $\pi_A R \subseteq \pi_A S$
  - Certainly not true in general
  - But holds in the common case of foreign key joins
- $|Q| \approx \frac{|R| \cdot |S|}{\max(|\pi_A R|, |\pi_A S|)}$
  - Selectivity factor of $R.A = S.A$ is ${}^{1}/_{\max(|\pi_A R|, |\pi_A S|)}$

Example

- Database:
  - User(uid, name, age, pop), Member(gid,uid,date), Group(gid, gname)
  - |User|=1000 rows, |Group|=100 rows, |Member|=50000 rows
  - $|\pi_{name}(User)| = 50, \pi_{pop}(User) = \{1,2,3,4,5\}$
  - $|\pi_{uid}(Member)| = 500$

- Estimate size $|User \bowtie Member| = ?$
  - $|\,\pi_{uid}(User)| = 1000$
  - $|\pi_{uid}(Member)| = 500$
  - 1000*50000/max(500,1000)=50000

Other estimations

- Using similar ideas, we can estimate the size of projection, duplicate elimination, union, difference, aggregation (with grouping)

- Lots of assumptions and very rough estimation
  - Accurate estimate is not needed
  - Maybe okay if we overestimate or underestimate consistently
  - May lead to very nasty optimizer "hints"

    SELECT * FROM User WHERE pop > 0.9;

    SELECT * FROM User WHERE pop > 0.9 AND pop > 0.9;

- Not covered: better estimation using histograms

Case Study

Physical plan example:

- PROJECT (*Group.name*)
  - INDEX-NESTED-LOOP-JOIN (*gid*)
    - INDEX-NESTED-LOOP-JOIN (*uid*)
      - INDEX-SCAN (*name* = “Bart”)
        - Index on *User*(*name*)
      - Index on *Member*(*uid*)
    - Index on *Group*(*gid*)

- System requirements:
  - Each disk/memory block can hold up to 10 rows (from any table);
  - All tables are stored compactly on disk (10 rows per block);
  - 8 memory blocks are available for query processing: M=8
- Database:
  - User(uid, age, pop), Member(gid,uid,date), Group(gid, gname)
  - |User|=1000 rows, |Group|=100 rows, |Member|=50000 rows
  - #of blocks: B(User)=1000/10=100; B(Group)=100/10=10; B(Member)=50000/10=5k

Case Study

Physical plan example:

PROJECT (*Group.name*)

INDEX-NESTED-LOOP-JOIN (*gid*)

Index on *Group*(*gid*)

INDEX-NESTED-LOOP-JOIN (*uid*)

Index on *Member*(*uid*)

INDEX-SCAN (*name* = “Bart”)

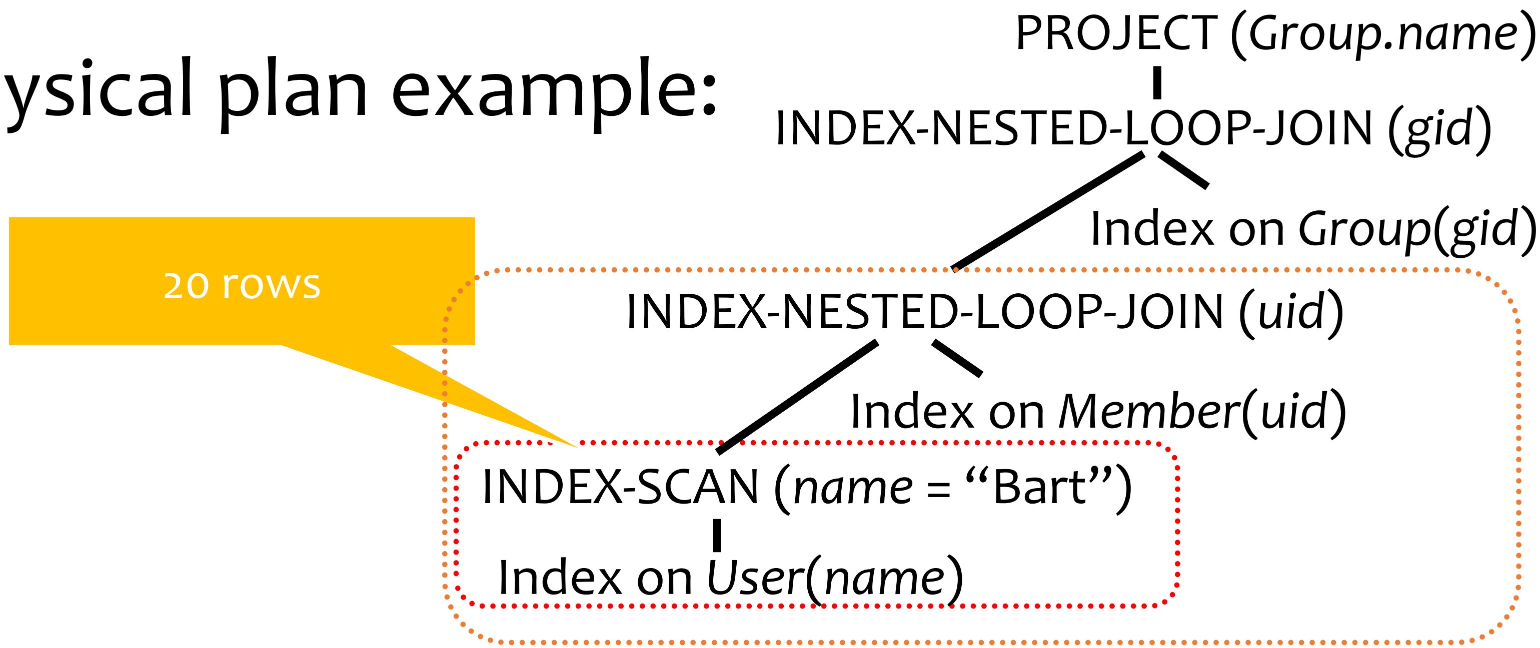
Index on *User*(*name*)

20 rows

- |User|=1000, $|\pi_{name}(User)| = 50$ ➔ $|\sigma_{name="Bart"}(User)| = \frac{1000}{50} = 20$ records
- INDEX-SCAN on User
    - IO COST: index lookup (4 IOs, depending on the height of the tree)

Case Study

Physical plan example:



- |User|=1000, $|\pi_{name}(User)| = 50$ ➔ $|\sigma_{name="Bart"}(User)| = \frac{1000}{50} = 20$ records
- INDEX-SCAN on User
  - IO COST: index lookup (4 IOs, depending on the height of the index tree)
- JOIN: For each record with name = "Bart", probe the index on $Member(uid)$
  - IO cost: $B(R) + |R| \cdot (\text{index lookup})$
  - 20 rows are not clustered → at worst case, 20 blocks of data to be retrieved
  - 20 + 20 * (4 IOs for index lookup)

Case Study

Physical plan example:

PROJECT (*Group.name*)

INDEX-NESTED-LOOP-JOIN (*gid*)

Index on *Group*(*gid*)

INDEX-NESTED-LOOP-JOIN (*uid*)

Index on *Member*(*uid*)

INDEX-SCAN (*name* = "Bart")

Index on *User*(*name*)

2k rows ~
B(input) = 200 blocks

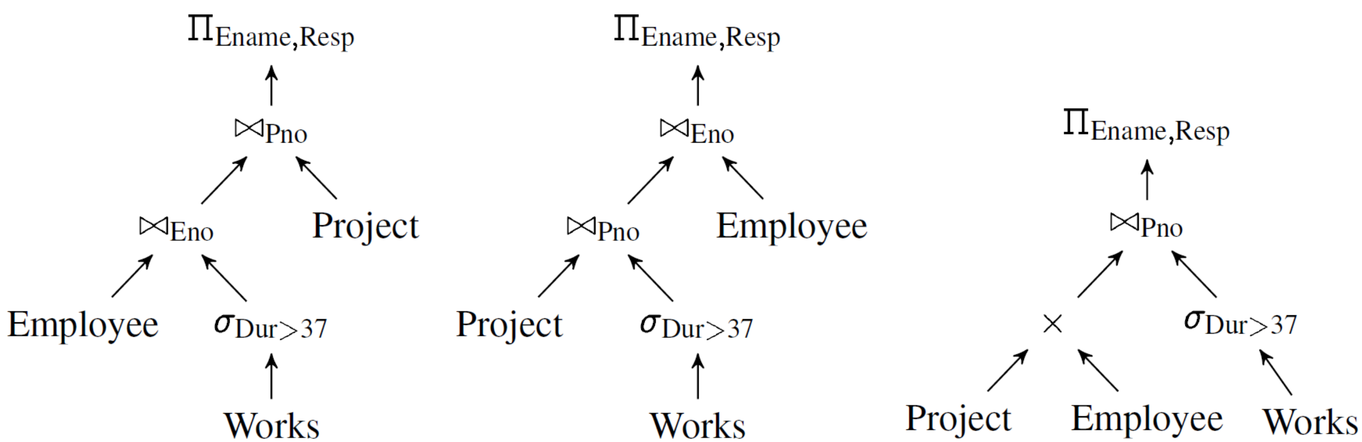
- Given $|\pi_{uid}(\sigma_{name="Bart"}User)| = 20, \ \ |\pi_{uid}(Member)| = 500$
- $|JOIN(uid)| \approx \frac{|R| \cdot |S|}{\max(|\pi_A R|, |\pi_A S|)} = \frac{20 \cdot 50k}{\max(20,500)} = \frac{1000k}{500} = 2k$
- Exercise: what is the IO cost for the next INDEX-NESTED-LOOP-JOIN(gid)?

Outline

- System view of query processing
  - Logical plan and physical plan
- Cost calculation of the physical plan
  - Cardinality estimation
- Search space and search strategy
  - Transformation rules
  - Heuristic approach

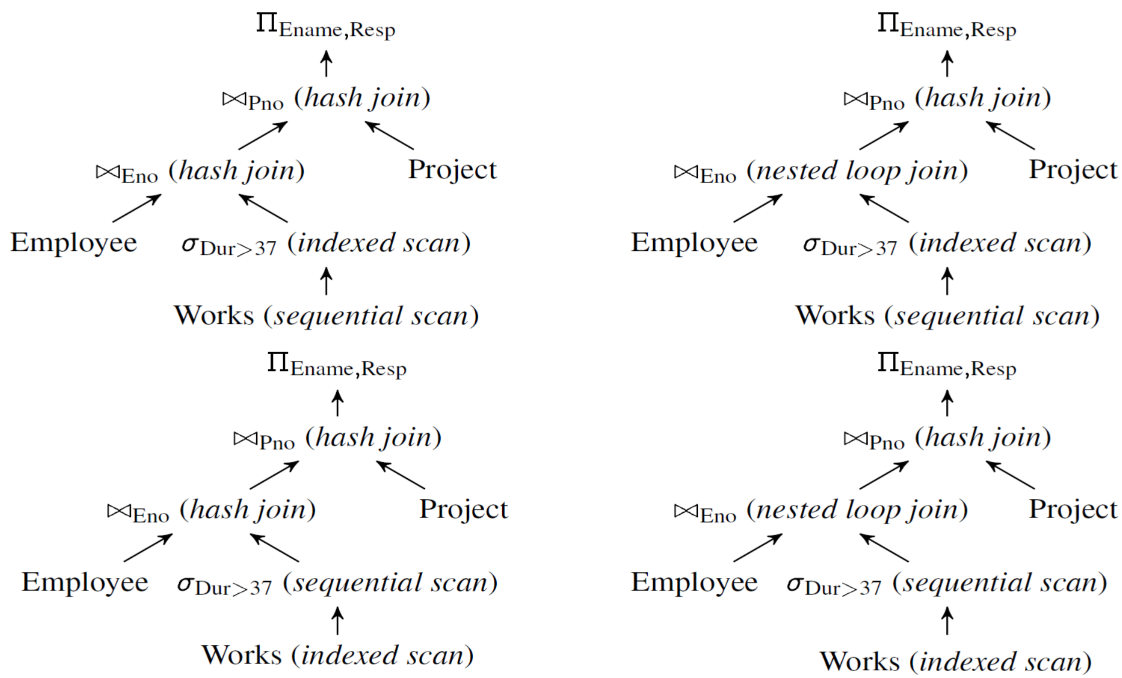
Search space is huge

- Characterized by "equivalent" logical query plans
  - **select** E.Ename, W. Resp
    from Employee E, Projects P, Works W
    where E.ENo = W.Eno **and** W.Pno=P.Pno **and** W.Dur > 37



This gets complicated very quickly

- Each logical plan can have multiple physical plans



- Do we need to exam all the logical plans?
  - No. We can use apply heuristic transformation rules to find a cheaper logical plan

Transformation rules (a sample)

- Convert σ_p-$\times$ to/from $\bowtie_p$: $\sigma_p(R \times S) = R \bowtie_p S$
  - Example: $\sigma_{User.uid=Member.uid}(User \times Member) = User \bowtie Member$

- Merge/split σ's: $\sigma_{p_1}\left(\sigma_{p_2}R\right) = \sigma_{p_1 \wedge p_2}R$
  - Example: $\sigma_{age>20}\left(\sigma_{pop=0.8}User\right) = \sigma_{age>20 \wedge pop=0.8}User$

- Merge/split π's: $\pi_{L_1}\left(\pi_{L_2}R\right) = \pi_{L_1}R$, where $L_1 \subseteq L_2$
  - Example: $\pi_{age}\left(\pi_{age,pop}User\right) = \pi_{age}User$

Transformation rules (a sample)

- Push down/pull up σ:
  $\sigma_{p \wedge p_r \wedge p_s}(R \bowtie_{p'} S) = (\sigma_{p_r} R) \bowtie_{p \wedge p'} (\sigma_{p_s} S)$, where
  - p_r is a predicate involving only R columns
  - p_s is a predicate involving only S columns
  - p and p' are predicates involving both R and S columns
  - Example:

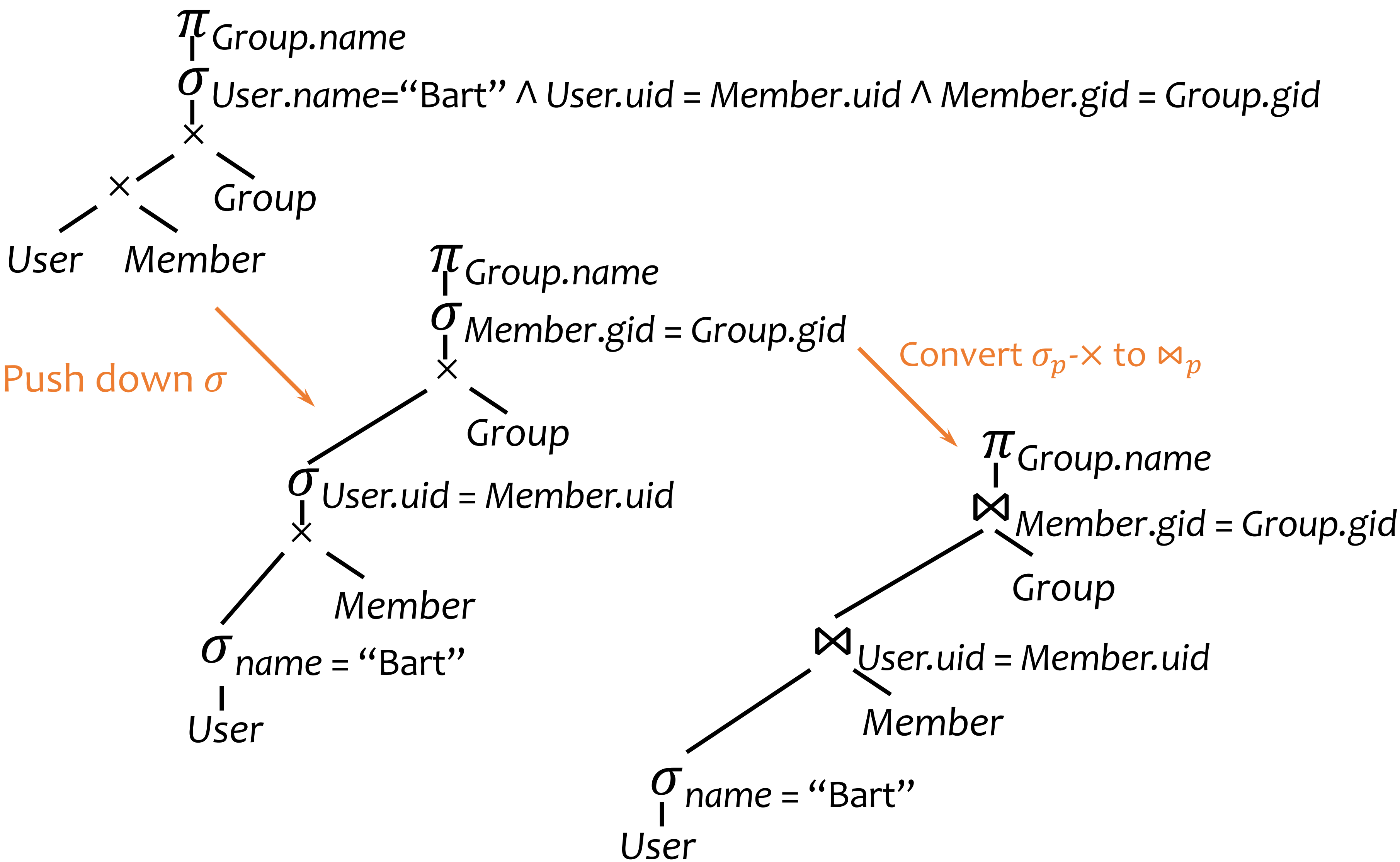
$$\sigma_{\text{U1.name=U2.name} \wedge U1.\text{pop}>0.8 \wedge U2.pop>0.8}(\rho_{U1} User \bowtie_{U1.uid \neq U2.uid} \rho_{U2} User)$$
$$= \sigma_{pop>0.8}(\rho_{U1} User) \bowtie_{U1.uid \neq U2.uid, U1.name=U2.name} (\sigma_{pop>0.8}(\rho_{U2} User))$$

Transformation rules (a sample)

- Push down π: $\pi_L(\sigma_p R) = \pi_L\left(\sigma_p(\pi_{LL'} R)\right)$, where
  - L' is the set of columns referenced by p that are not in L
  - Example:
    $\pi_{age}(\sigma_{pop>0.8} User) = \pi_{age}(\sigma_{pop>0.8}(\pi_{age,pop} User))$

- Many more (seemingly trivial) equivalences…
  - Can be systematically used to transform a plan to new ones

Relational query rewrite example



Heuristics-based query optimization

- Start with a logical plan
- Push selections/projections down as much as possible
  - Why? Reduce the size of intermediate results
  - Why not? May be expensive; maybe joins filter better
- Join smaller relations first, and avoid cross product
  - Why? Reduce the size of intermediate results
  - Why not? Size depends on join selectivity too
- Convert the transformed logical plan to a physical plan (by choosing appropriate physical operators)

Search strategy

- Heuristics-based optimization
  - Apply heuristics to rewrite "logical plans" into cheaper ones
- Cost-based optimization
  - Need statistics to estimate sizes of intermediate results to find the best "physical plan"

→ Course CS448 "Database Systems Implementation"

Summary

- System view of query processing
  - Logical plan and physical plan
- Heuristics-based optimization
  - Apply heuristics to rewrite “logical plans” into cheaper ones
- Cost-based optimization
  - Need statistics to estimate sizes of intermediate results to find the best “physical plan”