Spinlocks and WaitChannels are the low-level equivalents to Mutexes and Condition Variables. Spinlocks are designed to lock small sequences of code that do not include sleeps. WaitChannels maintain a list structure that contains a list of sleeping threads and coordinates sleeping and waking up with the CPU scheduler.

Below is a partial implementation of Semaphores that uses spinlocks and waitchannels.

Question 1. Use hand-over-hand locking to fill in the missing lines of Semaphore_Wait.

Question 2. Fill in the missing code in Semaphore_Post.

```
1 typedef struct Semaphore {
 \mathbf{2}
     int sem_count;
 3
     Spinlock *sem_lock;
     WaitChannel *sem_wchan;
 4
 5 } Semaphore;
 6
 7 Semaphore_Wait(Semaphore *sem) {
     Spinlock_Lock(&sem->sem_lock);
 8
 9
     while (sem->sem_count == 0) {
10
       // Fill me in
11
12
13
14
15
16
17
18
19
20
     }
21
     sem->sem_count--;
22
     Spinlock_Unlock(&sem->sem_lock);
23 }
24
25 Semaphore_Post(Semaphore *sem) {
     // Fill me in
26
27
28
29
30
31
32
33 }
```