a) [1 mark]

Timesharing; **also accept** pre-emption

b) [4 marks; one for each correct point]

- 1. Thread block/sleep
- 2. Thread yield
- 3. Thread exit
- 4. preemption

c) [2 marks; one for each point]

- 1. semaphores have a counter instead of a boolean indicating they are held
- 2. semaphores have no ownership

d) [2 marks; one for each point --- if combined but correct, also accept]

thread_fork is passed a function pointer the new thread executes that function

e) [3 marks; one for each point]

- 1. The process has no living parent
- 2. When the parent process has called waitpid on the process, which has terminated
- 3. When the parent dies, zombie children can be fully deleted

ALSO ACCEPT for full marks

1. Any situation where a parent process is unable to call waitpid on the child

f) [2 marks; accept any reasonable answer]

an error code or 0

g) [2 marks; one for each point]

- 1. Security
- 2. Abstract design

h) [2 marks; one for each point]

- 1. Advantage: efficient OR easy to implement
- 2. Disadvantage: external fragmentation

a) [1 mark]

1 + 10 + 10 = 21

b) [6 marks; one for each correct answer]

- 1. NO
- 2. YES
- 3. NO
- 4. NO
- 5. NO
- 6. NO

c) [1 mark]

45

d) [1 mark]

No

Question 3 [12 marks total; deduct a mark for each mistake]

User	Kernel
Application frames (or user code, or)	trapframe
open()	mips_trap
	syscall
	sys_open
	trapframe
	mips_trap
	mainbus_interrupt [OPTIONAL, deduct no points if missing]
	timer_exception_handler
	thread_yeild
	thread_switch
	switchframe

a) [1 mark]

2^32/2^12 = 2^20

b) [1 mark]

2^48/2^12 = 2^36

d) [1 mark]

20

e) [1 mark]

36

f) [4 marks; one for each correct answer; accept hex answers]

1.	0x0000 0000	0
2.	0x0000 0ACE	0
3.	0x0110 EA5E	0x110E
4.	0x0000 00C5	0

g) [4 marks; one for each correct answer]

 $2^{13/2} = 2$ pages of memory (0 and 1)

- 1. VALID
- 2. VALID
- 3. INVALID
- 4. VALID

Question 5 [8 total marks; one for each point]

- 1. Acquire procTableLock
- 2. if ProcTable contains a process with name procName
 - a. Get pointer to process
 - b. Release lock
 - c. Return process PID
- 3. Otherwise
 - a. Release lock
 - b. Return ENOPROC

a) [7 marks; one for each point]

- bool try_acquire(lock *lk)
 - 1. acquire lk->spinlock
 - 2. If lock available
 - a. Take lock
 - b. Release lk->spinlock
 - c. Return true
 - 3. Release lk->spinlock
 - 4. Return false

b) [4 marks]

```
Yes (there is deadlock) [1 mark]
Corrections, identified with [], [1 mark each].
int total = 0;
   int account = 0;
   lock mutex = lock_create( "mutex" );
  void FuncA(int acc)
   {
     lock_acquire( mutex );
        for (i = 0 \text{ to } N)
        {
            total ++;
       }
     lock_acquire( mutex );
     account = acc;
     [lock_release(mutex)]
   }
  void FuncB(int acc, int val)
   {
     for (i = 0 \text{ to } N)
     {
           FuncA( acc );
          [lock_acquire(mutex)]
          total = total - val;
          [lock_release(mutex)]
     }
}
```

a) [2 marks; one for offset and one for segment number]

Segment Number = 2 bits Segment Offset = 30 bits

b) [2 marks; one for each point]

1. Add a register for each segment to indicate if the segment is read-only or not

2. If writing to a read-only segment; MMU must raise exception

c) [2 marks]

0

d) [2 marks; one for each point]

- 1. Clear relocation and limit registers
- 2. Load new process relocation and limit register values

e) [BONUS 2 marks; one for each point]

- 1. Clear STACK relocation and limit register
- 2. Load new THREAD STACK relocation and limit register values