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- Start as early as possible, and contact the instructor if you get stuck.
  - See the course outline for details about the course's marking policy and rules on collaboration.
  - Submit your completed solutions to **Crowdmark**.

1. Regular Expressions

[8] Let  $\Sigma = \{0, 1\}$ . Give a **rigorous** proof for the equality of languages

$$L((10^*)^*0) = L(0 + 1(0 + 1)^*0).$$

## 2. Closure Properties of Regular Languages

[8]

Let  $B$  and  $C$  be languages over  $\Sigma = \{0, 1\}$ . Define the binary operation on languages

$$B \stackrel{1}{\leftarrow} C = \{w \in B \mid \text{for some } y \in C, n_1(w) = n_1(y)\}.$$

Recall that  $n_1(w)$  denotes the number of occurrences of the symbol 1 in the string  $w$ . For example, if  $B = \{010, 101, 111\}$  and  $C = \{01, 011, 1111\}$ , then  $B \stackrel{1}{\leftarrow} C = \{010, 101\}$ .

Prove that the class of regular languages is closed under the  $\stackrel{1}{\leftarrow}$  operation.

## 3. Non-regular languages

Prove that each of the following languages is not regular.

- [4] (a)  $L_a = \{0^i 1^j \mid \gcd(i, j) = 1\}$  over  $\Sigma = \{0, 1\}$

[4]

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(b)  $L_b = \{a^m \mid m \neq n^2 \text{ for any } n \in \mathbb{N}\}$  over  $\Sigma = \{a\}$

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4. A Non-Regular Language In Which All Long Words Can Be Pumped

Let  $\Sigma = \{a, b, c\}$ .

- [4] (a) Prove that  $L = \{ab^j c^j \mid j \geq 0\}$  is not regular.

[4]

(b) Prove that  $F = \{a^i b^j c^k \mid i, j, k \geq 0 \text{ and if } i = 1 \text{ then } j = k\}$  is not regular.

[4]

- (c) Exhibit with proof a choice of a positive integer  $n$ , such that, for any  $z \in F$  with  $|z| \geq n$ , we may write  $z = uvw$  where
- $|uv| \leq n$ ,
  - $|v| \geq 1$  and
  - $uv^i w \in F$ , for all  $i \geq 0$ .

[2]

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(d) Explain briefly why the results of parts 4b and 4c do **not** contradict the Pumping Lemma for regular languages.



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5. Testing a Candidate Criterion for Regular Languages

Let  $\Sigma = \{0, 1\}$  be the alphabet for all languages in this problem.

(a) Prove that  $L_a = \{w \mid n_0(w) = n_1(w)\}$  is not regular.

[4]

[4]

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(b) Prove that  $L_b = \{w \mid n_0(w) \neq n_1(w)\}$  is not regular.

[4]

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(c) Suppose that  $L$  is a language over  $\Sigma$  and that there is a fixed integer  $k \geq 0$  such that, for every  $x \in \Sigma^*$ ,  $xz \in L$ , for some string  $z$  with  $|z| \leq k$ . Does it follow that  $L$  is regular? Prove your answer.