Timothy Chan's Convex Hull Algorithm (Chan 1995)
O(n lg n):combines Graham and Jarvis

Arbitrarily split P into subsets P_1, …, $P_{n/h}$ each of size h.
for k= I to n/h do
    U_k= upper bound of P_h (by Graham scan)
q_1= leftmost point
for 1= 1 to h do
    {
    if q_i == rightmost point return <q_1, … , q_i>
    q_{i+1} = rightmost point
    for k= 1 to n/h do {
        p_k= "right tangent" between q_i and U_k
        if p_k above q_i q_{i+1} then q_{i+1} = $p_{k\}}$
    }

Method works if estimate for h is too large … but runtime is n lg h, so we want a decent estimate for lg h. Start with lg h = 2 (ie h=4) and run the algorithm doubling estimate of lg h each time.