

Lecture 12: Householder QR factorizations

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Outline

- ① Householder Triangularization
- ② Householder QR Factorization Algorithm
- ③ Example: Householder Reflector
- ④ Example: QR Factorization via Householder

Householder QR factorizations - Introduction

- Recall that in this course we consider three common algorithms for QR factorization:
 - ① Gram-Schmidt orthogonalization,
 - ② Householder reflections,
 - ③ Givens rotations.
- Gram-Schmidt orthogonalization was discussed in Lecture 11.
- This lecture will introduce the idea of Householder reflections for building the QR factorization.
- A final approach of Givens rotations will be presented in the next lecture.

Householder QR factorizations - Introduction

- Note that the QR factorization we produce here is similar, but not identical, to the one we produced last time:
 - 1 R is $m \times n$ and Q is square, instead of the other way around, and
 - 2 Negative entries can occur on R 's “diagonal”.

Householder Triangularization

- Note that Gram-Schmidt orthogonalization is a “triangular orthogonalization” process. In matrix form, Gram-Schmidt can be written as right-multiplication by triangular matrices that make the columns of A orthonormal (see end of Lecture 8 of Trefethen & Bau)

$$A \underbrace{R_1 R_2 \cdots R_n}_{\hat{R}^{-1}} = \hat{Q}.$$

- Householder reflections** instead provide an “orthogonal triangularization” process. The matrix A is made to be triangular (R) by applying orthogonal matrices Q_j , i.e.,

$$\underbrace{Q_n \cdots Q_2 Q_1}_{Q^{-1}} A = R.$$

Hence, the premise of Householder reflections (aka triangularization) is to find the orthogonal matrices $Q_j \in \mathbb{R}^{m \times m}$. This method is similar to LU -factorization, as each Q_j will zero the lower entries of column j .

Householder Triangularization

$$\begin{bmatrix} \times & \times & \times \\ \times & \times & \times \\ \times & \times & \times \\ \times & \times & \times \end{bmatrix} \xrightarrow{Q_1} \begin{bmatrix} \times & \times & \times \\ 0 & \times & \times \\ 0 & \times & \times \\ 0 & \times & \times \end{bmatrix} \xrightarrow{Q_2} \begin{bmatrix} \times & \times & \times \\ 0 & \times & \times \\ 0 & 0 & \times \\ 0 & 0 & \times \end{bmatrix} \xrightarrow{Q_3} \begin{bmatrix} \times & \times & \times \\ 0 & \times & \times \\ 0 & 0 & \times \\ 0 & 0 & 0 \end{bmatrix}$$

A $Q_1 A$ $Q_2 Q_1 A$ $Q_3 Q_2 Q_1 A$

Householder Triangularization

We will build orthogonal matrices of the following form

$$Q_j = \begin{bmatrix} I & 0 \\ 0 & F \end{bmatrix} \quad \left. \begin{array}{l} \} j-1 \text{ rows, already done,} \\ \} m-(j-1) \text{ rows, still to be done.} \end{array} \right\}$$

where F is the **Householder reflector** matrix. F reflects a vector x across a (specific) hyperplane H to produce a vector **along the axis**.

Householder Triangularization

See Figure 1 for a visualization of applying the Householder reflector (note $e_1 = [1, 0, \dots, 0]^T$).

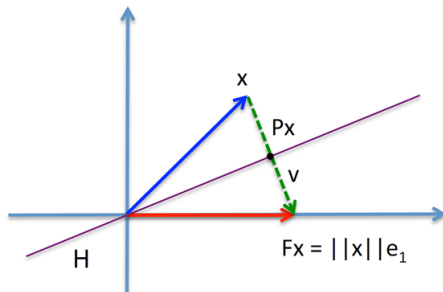


Figure: Applying the Householder reflector F to the vector x , which reflects x across the hyperplane H .

$$x + v^- = -||x||e_1 \Leftrightarrow v^- = -||x||e_1 - x.$$

Householder Triangularization

After reflection, the output vector has the same magnitude as x , and is parallel to e_1 . It depends on both of x and e_1 .

At step j , we start with x , and reflect onto the subspace spanned by $\{e_1, \dots, e_j\}$.

We find the Householder reflector matrix F , to perform the reflection, as follows.

Householder Triangularization

Suppose

$$x = \begin{bmatrix} \times \\ \times \\ \vdots \\ \times \end{bmatrix},$$

Find F such that

$$Fx = \begin{bmatrix} \|x\| \\ 0 \\ \vdots \\ 0 \end{bmatrix} = \|x\| e_1.$$

- The F reflects x across the hyperplane H **orthogonal to** $v = \|x\|e_1 - x$.
- That is because we want to produce a new vector of the same length as x , aligned with the axis e_1 (so all but the first entry are zeros).

Householder Triangularization

- The **orthogonal projection** P of x **onto** the hyperplane H (orthogonal to the vector v) is

$$P_x = x - \left(\left(\frac{v}{\|v\|} \right)^T x \right) \frac{v}{\|v\|} = x - v \left(\frac{v^T x}{v^T v} \right).$$

- Note that this orthogonal projection P is similar to the steps in Gram-Schmidt orthogonalization.
- **Idea:** Subtract out the component of x **along** v .

Householder Triangularization

- To **reflect** x across H (instead of projecting onto H) we must go twice as far in the same direction (see Figure 1)

$$Fx = x - 2v \left(\frac{v^T x}{v^T v} \right).$$

- Therefore, the **Householder reflector** F is given by

$$F = I - 2 \left(\frac{vv^T}{v^T v} \right)$$

where $v = \|x\|e_1 - x$.

Householder Triangularization

Remark that we could instead reflect to the point along the axis with a **negative** sign. That is, reflect to $-\|x\|e_1$ instead of $\|x\|e_1$, which gives

$$F_x = \begin{bmatrix} -\|x\| \\ 0 \\ \vdots \\ 0 \end{bmatrix} = -\|x\|e_1.$$

Either choice zeros out the desired entries of the active column. We just get a different v as shown in Figure 2.

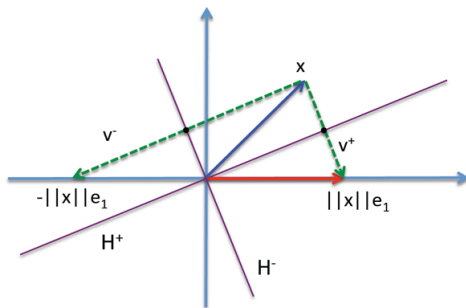


Figure: The two alternative Householder reflections.

Householder Triangularization

- Reflecting to either of $\|x\|e_i$ or $-\|x\|e_i$ will zero the remainder of the desired column.
- Which Householder reflector F should we choose?
- For numerical stability, we want the F that reflects x farther away from itself.
- Thus,
 - if $x_1 > 0$ we choose the negative one, $-\|x\|e_1$,
 - if $x_1 < 0$ we choose the positive one, $\|x\|e_1$.
- This gives $v = -\text{sign}(x_1)\|x\|e_1 - x$, or more simply (because only direction is important, and either choice gives the same F) $\text{sign}(x_1)\|x\|e_1 + x$.
- This choice of v avoids subtracting nearby quantities, which can introduce cancellation error.
- Therefore, choosing the F that reflects x farther away is more numerically stable.

Householder Triangularization

Alternative Derivation

- We will show an alternate derivation of Householder triangularization.
- Consider a reflection operator $F = I - 2\frac{vv^T}{v^T v}$ for an arbitrary vector v .
- We want to find v such that $Fx \in \text{span}\{e_1\}$ to zero the lower entries in column 1.
- Let $Fx \in \text{span}\{e_1\}$, in other words,

$$Fx = x - \frac{2vv^T x}{v^T v} = x - \left[\frac{2(v^T x)}{(v^T v)} \right] v \in \text{span}\{e_1\}.$$

Householder Triangularization

- Observe $v \in \text{span}\{e_1, x\}$ by construction, since $Fx \in \text{span}\{e_1\}$.
- Write $Fx = c_2 e_1$ for some scalar c_2 .
- Hence,

$$\begin{aligned}c_2 e_1 &= x - c_1 v \\ \Rightarrow v &= \hat{c}_1 x + \hat{c}_2 e_1,\end{aligned}$$

(for scalars \hat{c}_1 and \hat{c}_2) which means $v \in \text{span}\{e_1, x\}$.

Householder Triangularization

- Now let $\hat{v} = x + \alpha e_1$ for some scalar α . (Note the length of v does not matter; only its direction matters.)
- We will write v for \hat{v} from now on.
- Then,

$$\begin{aligned} v^T x &= (x + \alpha e_1)^T x \\ &= x^T x + \alpha e_1^T x \\ &= x^T x + \alpha \underbrace{x_{(1)}}_{\text{scalar}} \end{aligned}$$

and

$$\begin{aligned} v^T v &= (x + \alpha e_1)^T (x + \alpha e_1) \\ &= x^T x + 2\alpha x_{(1)} + \alpha^2. \end{aligned}$$

Householder Triangularization

Plugging into Fx to determine α , we have

$$\begin{aligned}Fx &= x - 2 \left(\frac{v^T x}{v^T v} \right) (x + \alpha e_1), \\&= \left(1 - \frac{2v^T x}{v^T v} \right) x - \left(2\alpha \frac{v^T x}{v^T v} \right) e_1, \\&= \left(1 - \frac{2(x^T x + \alpha x_{(1)})}{x^T x + 2\alpha x_{(1)} + \alpha^2} \right) x - \left(2\alpha \frac{v^T x}{v^T v} \right) e_1, \\&= \left(\frac{x^T x + \cancel{2\alpha x_{(1)}} + \alpha^2 - 2x^T x - \cancel{2\alpha x_{(1)}}}{x^T x + 2\alpha x_{(1)} + \alpha^2} \right) x - \left(2\alpha \frac{v^T x}{v^T v} \right) e_1, \\&= \underbrace{\left(\frac{\alpha^2 - x^T x}{x^T x + 2\alpha x_{(1)} + \alpha^2} \right)}_{\text{must be 0}} x - \left(2\alpha \frac{v^T x}{v^T v} \right) e_1.\end{aligned}$$

Since $Fx \in \text{span}\{e_1\}$ the first term must be zero, so $\alpha^2 - x^T x = 0 \Rightarrow \alpha = \pm \|x\|$.

Householder Triangularization

Hence,

$$v = x \pm \|x\|e_1 \quad \text{and} \quad Fx = \mp\|x\|e_1,$$

as we saw last time.

Householder QR Factorization Algorithm

Algorithm 1 gives the QR factorization of A via Householder triangularization.

Algorithm 1 : Householder QR factorization algorithm

```
for  $k = 1, 2, \dots, n$   
     $x = A(k : m, k)$  ▷ Get current column  
     $v_k = \text{sign}(x_1) \|x\| e_1 + x$  ▷ Form the reflection vector  
     $v_k = \frac{v_k}{\|v_k\|}$  ▷ Normalize  
    for  $j = k, k + 1, \dots, n$  ▷ Apply  $F$  to active lower-right block  
         $A(k : m, j) = A(k : m, j) - 2v_k(v_k^T A(k : m, j))$   
    end for  
end for
```

Householder QR Factorization Algorithm

- The notation used follows Matlab, i.e., $A(k : m, j) = j^{\text{th}}$ column of A from row k to row m .
- The algorithm converts A into R (upper “triangular”) using Householder reflections F .
- Note that one could **further** vectorize the inner loop (more efficient in Matlab) to the matrix operation

$$A(k : m, k : n) = A(k : m, k : n) - 2v_k(v_k^T A(k : m, k : n)).$$

Householder QR Factorization Algorithm

- Algorithm 1 does not construct Q , only the vectors v_k . Why is this not a problem in practice?
- We often do not need Q but just the products $Q^T b$ or Qx (e.g., for least squares we solve $Rx = Q^T b$). Since

$$\begin{aligned}Q^T &= Q_n Q_{n-1} \dots Q_2 Q_1, \\Q &= Q_1^T Q_2^T \dots Q_{n-1}^T Q_n^T,\end{aligned}$$

we can efficiently (in $O(mn)$ flops) compute $Q^T b$ or Qx using just the v_k 's to apply the appropriate reflections (see Algorithms 2 and 3). **Note the parentheses**, we compute $v(v^T b)$ rather than $(vv^T)b$ to avoid forming the matrix vv^T .

Algorithm 2 : Implicit $Q^T b$

```
for  $k = 1, 2, \dots, n$   
     $b(k : m) = b(k : m) -$   
     $2v_k (v_k^T b(k : m))$   
end for
```

Algorithm 3 : Implicit Qx

```
for  $k = n, n-1, \dots, 1$   
     $x(k : m) = x(k : m) -$   
     $2v_k (v_k^T x(k : m))$   
end for
```

Householder QR Factorization Algorithm

- Explicitly building the matrix Q may sometimes be necessary.
- So how **could** we use these implicit products to recover Q itself?
- We compute the product $QI = Q$ by applying Q to the columns of the identity matrix I (e_1, e_2, \dots), i.e., $q_1 = Qe_1, q_2 = Qe_2, \dots$.
- For the reduced QR factorization

$$A = \hat{Q}\hat{R} = \begin{bmatrix} & \\ & \\ & \end{bmatrix}_{m \times n} \begin{bmatrix} & \\ & \end{bmatrix}_{n \times n}.$$

- So \hat{Q} is given by just the first n columns of I

$$\hat{Q} = \begin{bmatrix} | & | & & | \\ Qe_1 & Qe_2 & \cdots & Qe_n \\ | & | & & | \end{bmatrix}.$$

Householder QR Factorization Algorithm

Complexity of Householder-Based QR Work is dominated by inner loop

$$A(k : m, j) = A(k : m, j) - 2v_k \left(v_k^T A(k : m, j) \right).$$

Tallying the cost gives:

- $v_k^T A(k : m, j) \approx 2(m - k + 1)$ flops (dot product),
- $2v_k (v_k^T A(k : m, j)) \approx (m - k + 1)$ flops (scalar multiply),
- $A(k : m, j) - 2v_k (v_k^T A(k : m, j)) \approx (m - k + 1)$ flops (subtraction).

So it approximately costs $4(m - k + 1)$ flops per inner step, which is done $(n - k + 1)$ times (j -loop). This totals to $4(m - k + 1)(n - k + 1)$ flops per outer iteration. The outer k -loop runs from 1 to n , which means the total flops can be approximated by

$$\sum_{k=1}^n 4(m - k + 1)(n - k + 1) \approx 2mn^2 - \frac{2}{3}m^3.$$

Note that this does **not** include forming Q .

Householder QR Factorization Algorithm

- For $m = n$ (square),
flops(Householder) $\approx \frac{4}{3}n^3 = 2 \times \text{flops(LU)}$.
- Recall from Lecture 11 that Gram-Schmidt orthogonalization cost $\approx 3 \times \text{flops(LU)}$.
- So Householder triangularization is faster than Gram-Schmidt orthogonalization.

Householder Reflector

- Given $x = \begin{bmatrix} 1 \\ 2 \\ 2 \end{bmatrix}$ find the Householder reflector F and the product Fx .
- Answer:** We see that $\|x\| = \sqrt{1^2 + 2^2 + 2^2} = 3$. Therefore

$$v = \pm\|x\|e_1 + x = \pm 3 \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} + \begin{bmatrix} 1 \\ 2 \\ 2 \end{bmatrix} = \begin{bmatrix} 4 \\ 2 \\ 2 \end{bmatrix} \text{ or } \begin{bmatrix} -2 \\ 2 \\ 2 \end{bmatrix}.$$

- We choose $v = \text{sign}(x_1)\|x\|e_1 + x = 3 \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} + \begin{bmatrix} 1 \\ 2 \\ 2 \end{bmatrix} = \begin{bmatrix} 4 \\ 2 \\ 2 \end{bmatrix}$, for its numerical stability.

Householder Reflector

$$\begin{aligned}\frac{vv^T}{v^Tv} &= \begin{pmatrix} \frac{1}{\begin{bmatrix} 4 & 2 & 2 \end{bmatrix} \begin{bmatrix} 4 \\ 2 \\ 2 \end{bmatrix}} \\ \begin{bmatrix} 4 & 2 & 2 \end{bmatrix} \end{pmatrix} \begin{bmatrix} 4 \\ 2 \\ 2 \end{bmatrix} \begin{bmatrix} 4 & 2 & 2 \end{bmatrix} \\ &= \left(\frac{1}{24}\right) \begin{bmatrix} 16 & 8 & 8 \\ 8 & 4 & 4 \\ 8 & 4 & 4 \end{bmatrix} \\ &= \left(\frac{1}{6}\right) \begin{bmatrix} 4 & 2 & 2 \\ 2 & 1 & 1 \\ 2 & 1 & 1 \end{bmatrix}\end{aligned}$$

Householder Reflector

$$\begin{aligned} F &= I - 2 \left(\frac{vv^T}{v^T v} \right) \\ &= \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} - \frac{1}{3} \begin{bmatrix} 4 & 2 & 2 \\ 2 & 1 & 1 \\ 2 & 1 & 1 \end{bmatrix} \\ &= \frac{1}{3} \begin{bmatrix} -1 & -2 & -2 \\ -2 & 2 & -1 \\ -2 & -1 & 2 \end{bmatrix}, \text{ so that} \\ Fx &= \frac{1}{3} \begin{bmatrix} -1 & -2 & -2 \\ -2 & 2 & -1 \\ -2 & -1 & 2 \end{bmatrix} \begin{bmatrix} 1 \\ 2 \\ 2 \end{bmatrix} \\ &= \begin{bmatrix} -3 \\ 0 \\ 0 \end{bmatrix}. \end{aligned}$$

Householder Reflector

- Since we are computing exactly, we are not forced to consider numerical stability.
- We re-compute everything with the other possible choice,

namely $v = \begin{bmatrix} -2 \\ 2 \\ 2 \end{bmatrix}$.

Householder Reflector

$$\begin{aligned}\frac{vv^T}{v^Tv} &= \left(\frac{1}{[-2 \ 2 \ 2] \begin{bmatrix} -2 \\ 2 \\ 2 \end{bmatrix}} \right) \begin{bmatrix} -2 \\ 2 \\ 2 \end{bmatrix} [-2 \ 2 \ 2] \\ &= \left(\frac{1}{12} \right) \begin{bmatrix} 4 & -4 & -4 \\ -4 & 4 & 4 \\ -4 & 4 & 4 \end{bmatrix} \\ &= \left(\frac{1}{3} \right) \begin{bmatrix} 1 & -1 & -1 \\ -1 & 1 & 1 \\ -1 & 1 & 1 \end{bmatrix}\end{aligned}$$

Householder Reflector

$$\begin{aligned} F &= I - 2 \left(\frac{vv^T}{v^T v} \right) \\ &= \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} - \frac{2}{3} \begin{bmatrix} 1 & -1 & -1 \\ -1 & 1 & 1 \\ -1 & 1 & 1 \end{bmatrix} \\ &= \frac{1}{3} \begin{bmatrix} 1 & 2 & 2 \\ 2 & 1 & -2 \\ 2 & -2 & 1 \end{bmatrix}, \text{ so that} \\ Fx &= \frac{1}{3} \begin{bmatrix} 1 & 2 & 2 \\ 2 & 1 & -2 \\ 2 & -2 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 2 \\ 2 \end{bmatrix} \\ &= \begin{bmatrix} 3 \\ 0 \\ 0 \end{bmatrix} . \end{aligned}$$

Example: QR Factorization via Householder

Perform QR factorization using Householder reflections on the matrix

$$A = \begin{bmatrix} 1 & -4 \\ 2 & 3 \\ 2 & 2 \end{bmatrix}$$

Example: QR Factorization via Householder

Step 1: $x = \begin{bmatrix} 1 \\ 2 \\ 2 \end{bmatrix}$, $v_- = \begin{bmatrix} 4 \\ 2 \\ 2 \end{bmatrix}$, $v_+ = \begin{bmatrix} 2 \\ -2 \\ -2 \end{bmatrix}$

(sign of v is irrelevant, since we only use vv^T and $v^T v$.) Ordinarily we would use v_- , but let's use v_+ for variety.

Then $F_1 = \begin{bmatrix} \frac{1}{3} & \frac{2}{3} & \frac{2}{3} \\ \frac{2}{3} & \frac{1}{3} & \frac{-2}{3} \\ \frac{2}{3} & \frac{-2}{3} & \frac{1}{3} \end{bmatrix} = Q_1$ and (by multiplying)

$$Q_1 A = \begin{bmatrix} 3 & 2 \\ 0 & -3 \\ 0 & -4 \end{bmatrix}$$

Example: QR Factorization via Householder

Step 2: $x = \begin{bmatrix} -3 \\ -4 \end{bmatrix}$ and so

$$\begin{aligned} v &= x + \text{sign}(x_1) \|x\| e_1 \\ &= \begin{bmatrix} -3 \\ -4 \end{bmatrix} + (-1) \cdot 5 \cdot \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} -8 \\ -4 \end{bmatrix} \end{aligned}$$

which lets us calculate

$$\begin{aligned} F_2 &= I - \frac{2vv^T}{v^T v} \\ &= \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} - \frac{2}{80} \begin{bmatrix} 64 & 32 \\ 32 & 16 \end{bmatrix} \\ &= \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} - \frac{1}{5} \begin{bmatrix} 8 & 4 \\ 4 & 2 \end{bmatrix} \\ &= \begin{bmatrix} \frac{-3}{5} & \frac{-4}{5} \\ \frac{-4}{5} & \frac{3}{5} \end{bmatrix} \end{aligned}$$

Example: QR Factorization via Householder

$$\text{Therefore } Q_2 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & F_2 \\ 0 & \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & \frac{-3}{5} & \frac{-4}{5} \\ 0 & \frac{-4}{5} & \frac{3}{5} \end{bmatrix}$$

$$\text{So } Q_2(Q_1 A) = R = \begin{bmatrix} 3 & 2 \\ 0 & 5 \\ 0 & 0 \end{bmatrix} \quad (\text{by multiplying})$$

Example: QR Factorization via Householder

Therefore

$$\begin{aligned} A &= Q_1^{-1} Q_2^{-1} R \\ &= Q_1^T Q_2^T R && \text{By orthogonality} \\ &= Q_1 Q_2 R && \text{by symmetry} \\ &= QR \end{aligned}$$

Example: QR Factorization via Householder

Orthogonality does not imply symmetry, but those Q 's were constructed to be symmetric, $I - \frac{2vv^T}{v^Tv}$.

$$\begin{aligned} Q = Q_1 Q_2 &= \begin{bmatrix} \frac{1}{3} & \frac{2}{3} & \frac{2}{3} \\ \frac{2}{3} & \frac{1}{3} & \frac{-2}{3} \\ \frac{2}{3} & \frac{-2}{3} & \frac{1}{3} \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & \frac{-3}{5} & \frac{-4}{5} \\ 0 & \frac{-4}{5} & \frac{3}{5} \end{bmatrix} \\ &= \frac{1}{15} \begin{bmatrix} 5 & -14 & -2 \\ 10 & 5 & -10 \\ 10 & 2 & 11 \end{bmatrix} \end{aligned}$$