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```

```
# Section 10.1
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#
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```
# Thm 10.1: Some example of the polynomials  $x^p - x$  over  $Z/(p)$ .
```

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#
```

```
> for i to 3 do
```

```
>   p := ithprime(i);
```

```
>   print(p, Factor( $x^p - x$ ) mod p);
```

```
> od:
```

```
2, x (x + 1)
```

```
3, x (x + 2) (x + 1)
```

```
5, (x + 3) x (x + 2) (x + 1) (x + 4)
```

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#####
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#
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```
# Thm 10.2: An example of the polynomial  $x^{(p^d)} - x$  over  $Z/(p)$ .
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```
#
```

```
# Consider  $p = 3$  and  $d = 4$ .
```

```
#
```

```
> Factor( $x^{(3^4)} - x$ ) mod 3;
```

```
4      3      2      2      4      3      2
(x  + 2 x  + x  + 2 x + 1) (x  + 2 x + 2) (x  + 2 x  + 2 x  + x + 2)
```

```
2      4      3      2      4      3      4      3      2
(x  + x + 2) (x  + 2 x  + x  + 1) (x  + 2 x  + 2) (x  + x  + x  + 2 x + 2)
```

```
4      3      2      4      3      2      4      3      4      2
(x  + 2 x  + x  + x + 2) (x  + x  + x  + 1) x (x  + x  + 2) (x  + 2 x  + 2)
```

```
4      2      4      2      4      2
(x  + 2 x + 2) (x  + 1) (x + 2) (x  + x  + x + 1) (x  + x  + 2)
```

```
4      2      4      3      2      4      3
(x  + x  + 2 x + 1) (x + 1) (x  + x  + 2 x  + 2 x + 2) (x  + 2 x  + x + 1)
```

$$(x^4 + x^3 + x^2 + x + 1)(x^4 + x + 2)(x^4 + x^3 + 2x + 1)$$

We know that $x^{(3^4)} - x$ is the product of all irreducible of
degree 1, 2 and 4. From inspecting the factorization above we
can see that there are 18 irreducibles of degree 4 over $Z/(3)[x]$.

#

But note that we can directly enumerate the number of distinct
irreducibles of degree d over $Z/(p)$, without factoring.

#

Question: How many distinct irreducibles of degree i are there over $Z/(3)$
for $i = 1, 2, 3, 4$?

#

Derivation:

#

Let k_i be the number of distinct irreducibles of degree i over $Z/(3)$.

Then we can set up a system of linear equations.

#

```
> sys := {degree(x^3-x,x) = k1*1,           # 1 divisible by 1
>         degree(x^(3^2)-x,x) = k1*1 + k2*2, # 2 divisible by 1, 2
>         degree(x^(3^3)-x,x) = k1*1 + k3*3, # 3 divisible by 1, 3
>         degree(x^(3^4)-x,x) = k1*1 + k2*2 + k4*4}: # 4 divisible by 1, 2, 4
```

```
> vars := {k1,k2,k3,k4};
```

```
vars := {k1, k2, k3, k4}
```

```
> sys;
```

```
{3 = k1, 9 = k1 + 2 k2, 27 = k1 + 3 k3, 81 = k1 + 2 k2 + 4 k4}
```

```
> solve(sys,vars);
```

```
{k1 = 3, k2 = 3, k3 = 8, k4 = 18}
```


#

As an aside, it is easy to derive the following function. In Maple,

```

# "option remember" stores the results of previous function calls so they
# don't need to be recomputed (or manually stored in a table like
# dynamic programming).
#
# Input: p - a prime
#        d - an integer in  $\mathbb{Z}_{\geq 1}$ 
#
# Output: the number of distinct irreducibles of degree d over  $\mathbb{Z}/(p)[x]$ 
#
> foo := proc(p,d)
>   option remember;
>   local i,c;

>   if d=1 then return p fi; # the degree of  $x^p - x$ 

>   # compute sum of degrees of all irreducibles in  $x^{(p^d)} - x$  of deg < d
>   c := 0;
>   for i to d-1 do
>     # degree i      # no. of irred. of degree i
>     if modp(d,i)=0 then c := c + i      * foo(p,i) fi
>   od;

>   return iquo(p^d - c,d);

> end:

> foo(3,4);

18

#
# Some more checks that foo is correct:
#
> degree(x^(5^4) - x,x) = foo(5,1)*1 + foo(5,2)*2 + foo(5,4)*4;
625 = 625

> degree(x^(17^26) - x,x) =foo(17,1)*1 + foo(17,2)*2 + foo(17,26)*26;
98100666009922840441972689847969 = 98100666009922830537394656942049

#####
#
# A key computational tool (for efficiency) that is used in the factoring
# algorithm is binary powering modulo another polynomial. This is

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# algorithm "RepeatedSquaring" in the script.
#
# One possible implementation is in the posted example "mypowmod.mpl".
#
# This operation is important enough that Maple has a built-in function
# for it, namely Powmod.
#
# Here is an example of an application of Powmod.
#
# Generate a rather large degree random polynomial with many factors
# modulo 3.
#
> f := mul(Randpoly(i,x)^i mod 3,i=1..60): f := Expand(f) mod 3:
#
# Let's not print out f! But do look at its degree.
#
> degree(f,x);

73810

#
# Compute the product of all irreducibles of degree 1 that divide f.
# (One copy of each).
#
> g1 := Gcd(x^3-x,f) mod 3;

3
g1 := x  + 2 x

#
# Check that f1 has only linear factors:
#
> Factor(g1) mod 3;

x (x + 2) (x + 1)

#
# To obtain all irreducibles of f of degree 1 (with multiplicity) we
# can compute the gcd of f with a high power of (x^3 - x). Since
# a linear factor can divide f at most deg f times, it suffices to compute
#
> g1_with_multiplicities := Gcd((x^3-x)^degree(f,x),f) mod 3:
> Factor(%) mod 3;

979      1119      704
x      (x + 2)      (x + 1)

```

```

#
# Instead of first computing  $(x^3 - x)^{\text{degree}(f,x)}$  and then taking the
# gcd, it is more efficient (polynomial time vs. exponential time!) to
# compute  $\text{Rem}((x^3-x)^{\text{degree}(f,x)}, f)$  using repeated squaring and then
# take the gcd:
#
> Gcd(Powmod((x^3-x), degree(f,x), f, x), f) mod 3:
> Factor(%) mod 3;

```

$$x^{979} (x + 2)^{1119} (x + 1)^{704}$$

```

#
# The above idea is used in the posted example "DDFact.mpl" which gives
# a function to compute the distinct degree factorization of the
# the squarefree part of f (even if f itself is not squarefree).
# Note: To improve the efficiency of that routine the computation
#  $\text{g}^{\text{ceil}(n/i)}$  should be replaced with the appropriate call to Powmod.

#####
#
# Section 10.2
#
# Let us first illustrate Thm 10.3.
#
# We know that  $R = \mathbb{Z}/(p)$  for  $p$  a prime is finite field with  $p$  element,
# in particular it is simply the set  $\{0, 1, \dots, p-1\}$  with addition and
# multiplication modulo  $p$ .
#
# The first part of Thm 10.3 states that if  $h$  is an irreducible of
# degree  $d$  then  $R[x]/\langle h \rangle$  is finite field with  $p^d$  elements.
#
# As an example, consider  $p = 3$  (so  $R = \mathbb{Z}/(3)$ ) and  $d = 2$ .
#
> h := x^2 + 1; # irred. of degree 2 over  $\mathbb{Z}/(3)$ 

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$$h := x^2 + 1$$

```

#
# Then  $RR = R[x]/\langle h \rangle$  is a finite field with  $p^d = 3^2 = 9$  elements.
# The elements of  $RR$  is all polynomials over  $R[x]$  of degree
# strictly less than 2 (i.e., the distinct polynomials of  $R[x]$  modulo  $h$ ).

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#
> RR := {seq(seq(i*x + j,i=0..2),j=0..2)};
      RR := {0, 1, 2, x, 2 x, x + 1, x + 2, 2 x + 1, 2 x + 2}

#
# Addition/subtraction in R is just addition/subtraction modulo 3.
# Multiplication is done modulo h, e.g., the product of (x+2) * (2*x+1) is
#
> Rem((x+2)*(2*x+1),h,x) mod 3;
      2 x

#
# The second part of Thm 10.3 should be familiar from our study
# of the Chinese remainder theorem, and algorithm such as multi-modular
# reduction.

#####
#
# Now consider Thm 10.4.
#
# First consider a prime  $Z/(p)$ . The theorem says that any nonzero
# element of  $Z/(p)$ , when raised to the power  $(p-1)/2$ , will be equal
# to 1 or -1, with exactly half equal to 1. Some experimental
# confirmation of the theorem.
#
> for i to 5 do
>   p := ithprime(i);
>   S := [seq(a,a=1..p-1)]; # nonzero elements of  $Z/(p)$ 
>   R := map(a->mods(a^((p-1)/2),p),S);
>   print(p,S,R);
> od:
      2, [1], [1]
      3, [1, 2], [1, -1]
      5, [1, 2, 3, 4], [1, -1, -1, 1]
      7, [1, 2, 3, 4, 5, 6], [1, 1, -1, 1, -1, -1]
     11, [1, 2, 3, 4, 5, 6, 7, 8, 9, 10], [1, -1, 1, 1, 1, -1, -1, -1, 1, -1]
#

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# Obviously, if we select an element of  $\{1,2,\dots,p-1\}$  uniformly
# at random, then with probably  $1/2$  it will be a quadratic residue
# and with probability  $1-1/2=1/2$  it will not be a non-quadratic residue.
#
#####
#
# Now let's give an illustration of the equal degree factorization
# on page 4 of the script. Let  $R = \mathbb{Z}/(p)$ ,  $p = 10000019$ .
#
> p := 10000019;
                                p := 10000019

> d := 2;
                                d := 2

> h1 := x^2+2090578*x+4297752: # irred. of degree 2
> h2 := x^2+4958404*x+3788058: # irred. of degree 2
> f := Expand(h1*h2) mod p;
                                4          3          2
                                f := x  + 7048982 x  + 8708093 x  + 5889928 x + 2913426

#
# Our goal is to factor f over  $R[x]$ . (Pretend we don't know h1 and h2.)
#
# The residue class ring
#
#  $R/\langle f \rangle \cong R/\langle h1 \rangle \times R/\langle h2 \rangle$  (*)
#
# contains  $(p)^2$  elements, of which  $(p-1)^2$  are relatively prime to f. (Why?)
#
# Four of the elements that are relatively prime to f are
#  $S = \{ (1,1), (1,-1), (-1,1), (-1,-1) \}$ : our goal is to uniformly and randomly
# select one of these four elements.
# On the left hand side of (*) the elements of S correspond to  $\{e1,e2,e3,e4\}$ :
#
> Gcdex(h1,h2,x,'s','t') mod p;
                                1

> e1 := Rem((1)*s*h1 + (1)*t*h2,f,x) mod p;
                                e1 := 1

> e2 := Rem((-1)*s*h1 + (1)*t*h2,f,x) mod p;

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          3          2
e2 := 2017226 x  + 4937057 x  + 544669 x + 3886491

> e3 := Rem((1)*s*h1 + (-1)*t*h2,f,x) mod p;
          3          2
e3 := 7982793 x  + 5062962 x  + 9455350 x + 6113528

> e4 := Rem((-1)*s*h1 + (-1)*t*h2,f,x) mod p;
e4 := 10000018

#
# Note that subtracting (1,1) from the elements in S gives the set
# { (0,0), (0,-2), (-2,0), (-2,-2) }: some of these (in this case
# two) have the nice property that they "split" the polynomial f:
#
> Gcd(e1-1,f) mod p; # should give us f (useless)
          4          3          2
x  + 7048982 x  + 8708093 x  + 5889928 x + 2913426

> Gcd(e2-1,f) mod p; # should give us h1 (great!)
          2
x  + 2090578 x + 4297752

> Gcd(e3-1,f) mod p; # should give us h2 (great!)
          2
x  + 4958404 x + 3788058

> Gcd(e4-1,f) mod p; # should give us 1 (useless)
          1

#
# Our goal is to select an element from {e1,e2,e3,e4} uniformly at random.
# To do so, we make use of Thm 10.4.
# First choose a random polynomial of degree < 4.
#
> u := Randpoly(4,x) mod p; # returns a random polynomial of degree 4
          4          3          2
u := 10323 x  + 1265632 x  + 4335155 x  + 3429475 x + 2575038

> u := u - lcoeff(u,x)*x^4; # now we have a random polynomial of degree < 4
          3          2
u := 1265632 x  + 4335155 x  + 3429475 x + 2575038

```



```

#
# We can check that u is relatively prime to f using Gcd.
#
> Gcd(f,u) mod p;
1

#
# [Excercise: What is the chance that our u is not relatively prime to f?]
# [Question: What if u is not realtively prime to f? Is this bad?]
#
# Assume now that u is relatively prime to f.
# Use Powmod to raise u to the power (p^d-1)/2.
#
> uu := Powmod(u,(p^2-1)/2,f,x) mod p;
uu := 10000018

#
# Then we must have uu in {e1,e2,e2,e4}.
# Let's try to take the gdd of uu - 1 with f:
#
> Gcd(uu-1,f) mod p;
1

#
# Unlucky! Try again.
#
> u := Randpoly(4,x) mod p; # returns a random polynomial of degree 4
u := 2262416 x4 + 2983381 x3 + 7057142 x2 + 5722259 x + 5518682

> u := u - lcoeff(u,x)*x^4; # now we have a random polynomial of degree < 4
u := 2983381 x3 + 7057142 x2 + 5722259 x + 5518682

> uu := Powmod(u,(p^2-1)/2,f,x) mod p;
uu := 1

> Gcd(uu-1,f) mod p;
x4 + 7048982 x3 + 8708093 x2 + 5889928 x + 2913426

```

```

#
# Unlucky! Try again.
#
> u := Randpoly(4,x) mod p; # returns a random polynomial of degree 4
      4      3      2
u := 4612324 x + 8986992 x + 3488871 x + 9954400 x + 5377058

> u := u - lcoeff(u,x)*x^4; # now we have a random polynomial of degree < 4
      3      2
u := 8986992 x + 3488871 x + 9954400 x + 5377058

> uu := Powmod(u,(p^2-1)/2,f,x) mod p;
      3      2
uu := 7982793 x + 5062962 x + 9455350 x + 6113528

> Gcd(uu-1,f) mod p;
      2
x + 4958404 x + 3788058

#
# Now we have split f.
> quit
memory used=117.6MB, alloc=32.6MB, time=40.29

```